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PLÁNOVÁNÍ A HRY - CV 3

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State – space Planning

- **Forward Search**
- **Backward Search**
- **Lifting**
- **STRIPS**

Forward Search

Forward-search(O, s_0, g)

$s \leftarrow s_0$

$\pi \leftarrow$ the empty plan

loop

if s satisfies g then return π

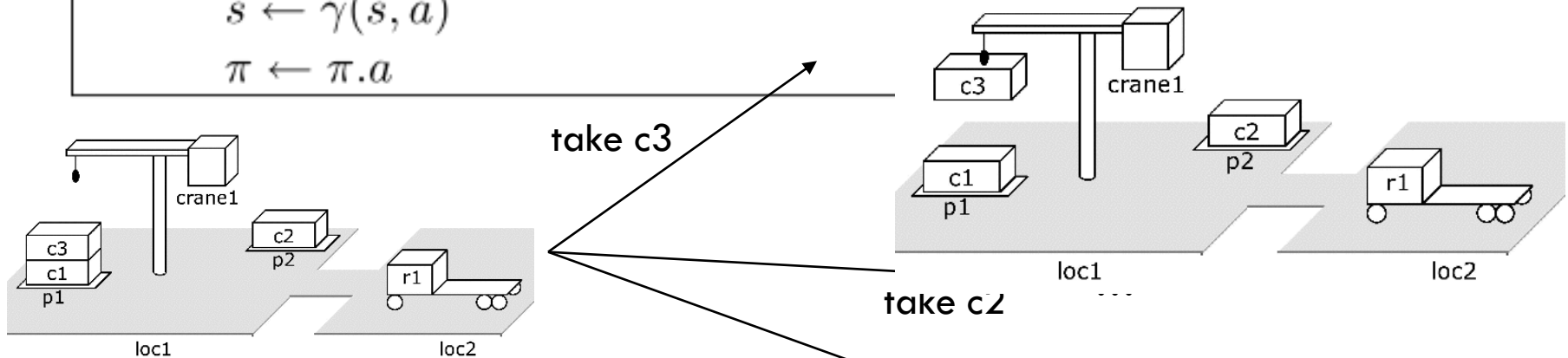
$E \leftarrow \{a \mid a \text{ is a ground instance an operator in } O,$
and $\text{precond}(a)$ is true in $s\}$

if $E = \emptyset$ then return failure

nondeterministically choose an action $a \in E$

$s \leftarrow \gamma(s, a)$

$\pi \leftarrow \pi.a$

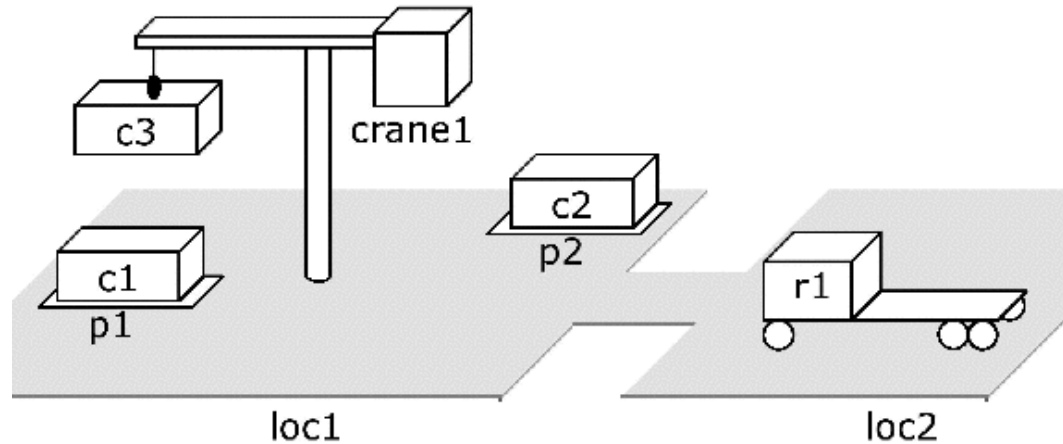


Forward Search Properties

- Forward-search is *sound*
 - ▣ for any plan returned by any of its nondeterministic traces, this plan is guaranteed to be a solution
- Forward-search also is *complete*
 - ▣ if a solution exists then at least one of Forward-search's nondeterministic traces will return a solution.

Task 1: DWR, find 1 finite and 1 infinite trace

□ s_0 :

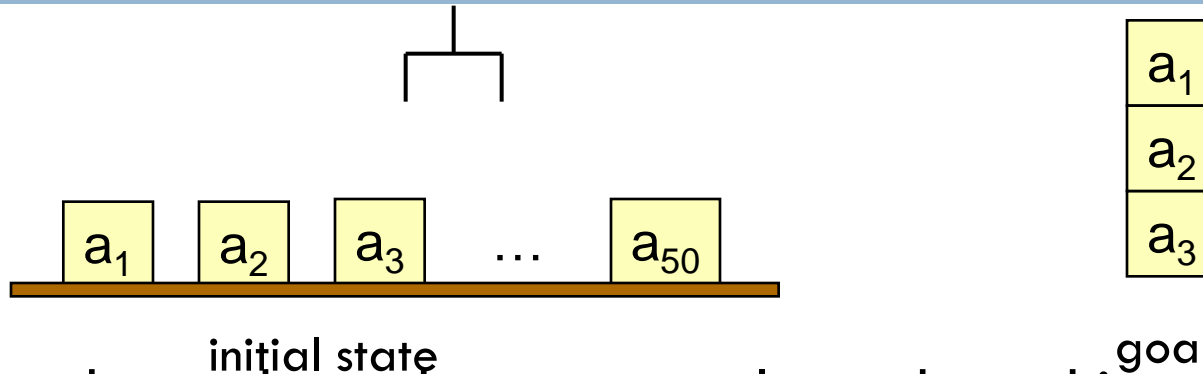


□ $g: \{at(r1, loc1), loaded(r1, c3)\}$

Task 2: Interchanging variables

- Objective: Interchange the values of variables v_1 and v_2 .
- $s_0 = \{\text{value}(v_1, 3), \text{value}(v_2, 5), \text{value}(v_3, 0)\}$
- $g = \{\text{value}(v_1, 5), \text{value}(v_2, 3)\}$
- $\text{assign}(v, w, x, y)$
 - ▣ precondition: $\text{value}(v, x), \text{value}(w, y)$
 - ▣ effects: $\neg \text{value}(v, x), \text{value}(v, y)$

Branching Factor of Forward Search



- Forward search can have a very large branching factor
 - ▣ E.g., many applicable actions that don't progress toward goal
- Why this is bad:
 - ▣ Deterministic implementations can waste time trying lots of irrelevant actions
- Need a good heuristic function and/or pruning procedure
- How to do pruning?

Backward Search

- For forward search, we started at the initial state and computed state transitions
 - ▣ new state = $\gamma(s,a)$
- For backward search, we start at the goal and compute inverse state transitions
 - ▣ new set of subgoals = $\gamma^{-1}(g,a)$
- To define $\gamma^{-1}(g,a)$, must first define *relevance*:
 - ▣ An action a is relevant for a goal g if
 - a makes at least one of g 's literals true
 - $g \cap \text{effects}(a) \neq \emptyset$
 - a does not make any of g 's literals false
 - $g^+ \cap \text{effects}^-(a) = \emptyset$ and $g^- \cap \text{effects}^+(a) = \emptyset$

Inverse State Transitions

- If a is relevant for g , then
 - $\gamma^{-1}(g,a) = (g - \text{effects}(a)) \cup \text{precond}(a)$
- Otherwise $\gamma^{-1}(g,a)$ is undefined
- Example: suppose that
 - $g = \{\text{on}(b1,b2), \text{on}(b2,b3)\}$
 - $a = \text{stack}(b1,b2)$
- What is $\gamma^{-1}(g,a)$?

Backward Search

Backward-search(O, s_0, g)

$\pi \leftarrow$ the empty plan

loop

if s_0 satisfies g then return π

$A \leftarrow \{a \mid a \text{ is a ground instance of an operator in } O$
and $\gamma^{-1}(g, a)$ is defined}

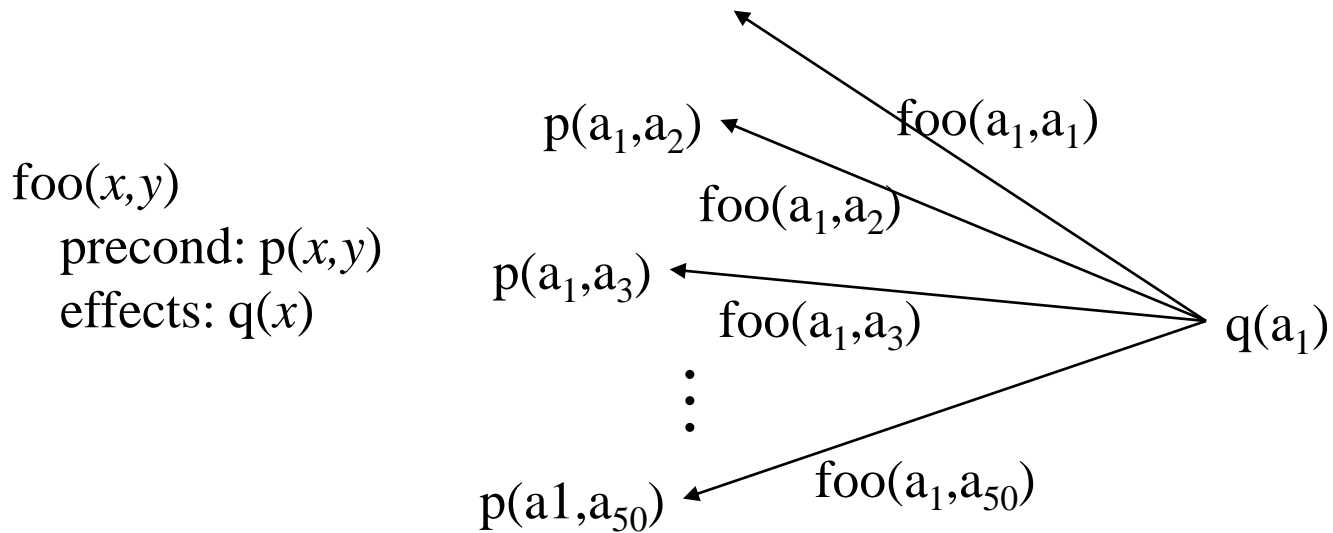
if $A = \emptyset$ then return failure

nondeterministically choose an action $a \in A$

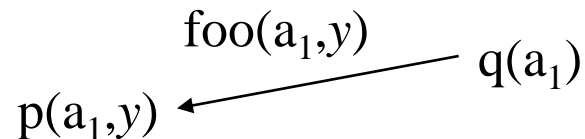
$\pi \leftarrow a.\pi$

$g \leftarrow \gamma^{-1}(g, a)$

Lifting



- Can reduce the branching factor of backward search if we *partially* instantiate the operators
 - ▣ this is called *lifting*



Lifted Backward Search

- More complicated than Backward-search
 - ▣ Have to keep track of what substitutions were performed
- But it has a much smaller branching factor

Lifted-backward-search(O, s_0, g)

$\pi \leftarrow$ the empty plan

loop

if s_0 satisfies g then return π

$A \leftarrow \{(o, \theta) \mid o \text{ is a standardization of an operator in } O,$
 $\theta \text{ is an mgu for an atom of } g \text{ and an atom of effects}^+(o),$
 $\text{and } \gamma^{-1}(\theta(g), \theta(o)) \text{ is defined}\}$

if $A = \emptyset$ then return failure

nondeterministically choose a pair $(o, \theta) \in A$

$\pi \leftarrow$ the concatenation of $\theta(o)$ and $\theta(\pi)$

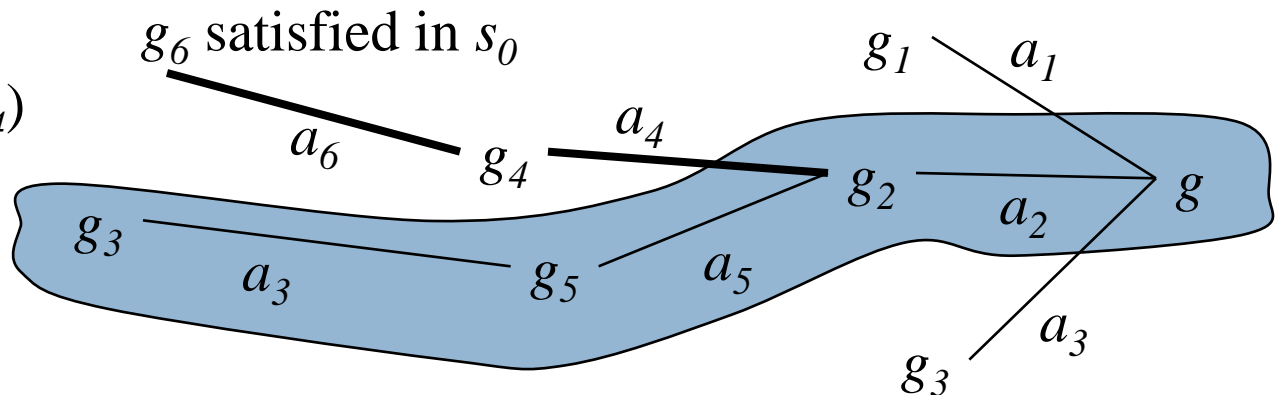
$g \leftarrow \gamma^{-1}(\theta(g), \theta(o))$

STRIPS

- $\pi \leftarrow$ the empty plan
- do a modified backward search from g
 - instead of $\gamma^{-1}(s,a)$, each new set of subgoals is just $\text{precond}(a)$
 - whenever you find an action that's executable in the current state, then go forward on the current search path as far as possible, executing actions and appending them to π
 - repeat until all goals are satisfied

$$\pi = \langle a_6, a_4 \rangle$$

$$s = \gamma(\gamma(s_0, a_6), a_4)$$



STRIPS

```
function groundStrips( $O, s, g$ )  
   $plan \leftarrow \langle \rangle$   
  loop  
    if  $s.satisfies(g)$  then return  $plan$   
     $applicables \leftarrow$   
      {ground instances from  $O$  relevant for  $g-s$ }  
    if  $applicables.isEmpty()$  then return failure  
     $action \leftarrow applicables.chooseOne()$   
     $subplan \leftarrow groundStrips(O, s, action.preconditions())$   
    if  $subplan = failure$  then return failure  
     $s \leftarrow \gamma(s, subplan \bullet \langle action \rangle)$   
     $plan \leftarrow plan \bullet subplan \bullet \langle action \rangle$ 
```

Blocks World ?

unstack(x,y)

Precond: $\text{on}(x,y)$, $\text{clear}(x)$, handempty

Effects: $\neg\text{on}(x,y)$, $\neg\text{clear}(x)$, $\neg\text{handempty}$,
 $\text{holding}(x)$, $\text{clear}(y)$

stack(x,y)

Precond: $\text{holding}(x)$, $\text{clear}(y)$

Effects: $\neg\text{holding}(x)$, $\neg\text{clear}(y)$,
 $\text{on}(x,y)$, $\text{clear}(x)$, handempty

pickup(x)

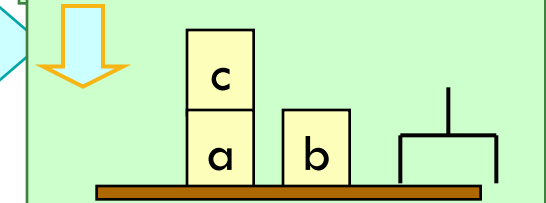
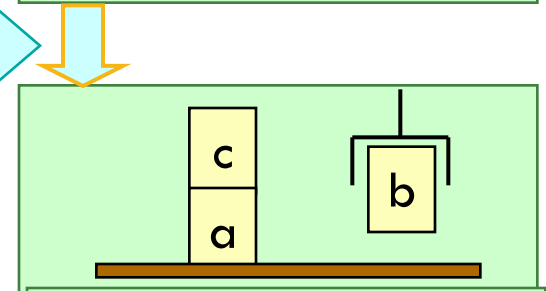
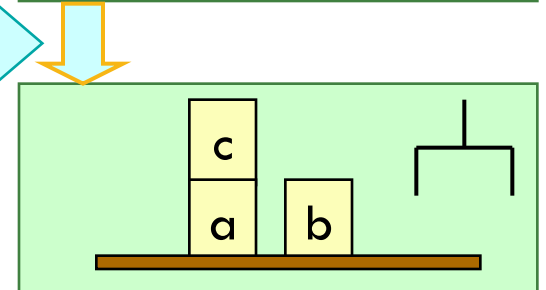
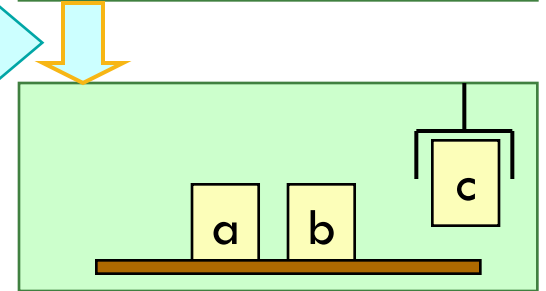
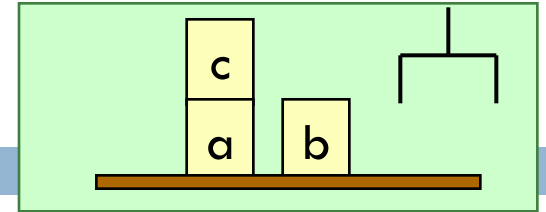
Precond: $\text{ontable}(x)$, $\text{clear}(x)$, handempty

Effects: $\neg\text{ontable}(x)$, $\neg\text{clear}(x)$,
 $\neg\text{handempty}$, $\text{holding}(x)$

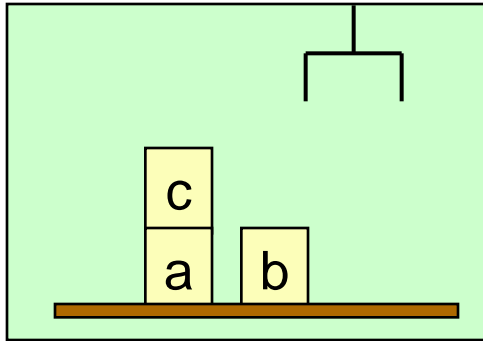
putdown(x)

Precond: $\text{holding}(x)$

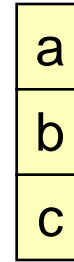
Effects: $\neg\text{holding}(x)$, $\text{ontable}(x)$,
 $\text{clear}(x)$, handempty



Sussman Anomaly



- Initial State
- Sub goals:
 - 1) Put A on B
 - 2) Put B on C



Goal

Interchanging Variables Repeated

- Objective: Interchange the values of variables v_1 and v_2 .
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- $g = \{\text{value}(v_1, 5), \text{value}(v_2, 3)\}$
- $\text{assign}(v, w, x, y)$
 - ▣ precondition: $\text{value}(v, x), \text{value}(w, y)$
 - ▣ effects: $\neg \text{value}(v, x), \text{value}(v, y)$

How to Handle Problems like These?

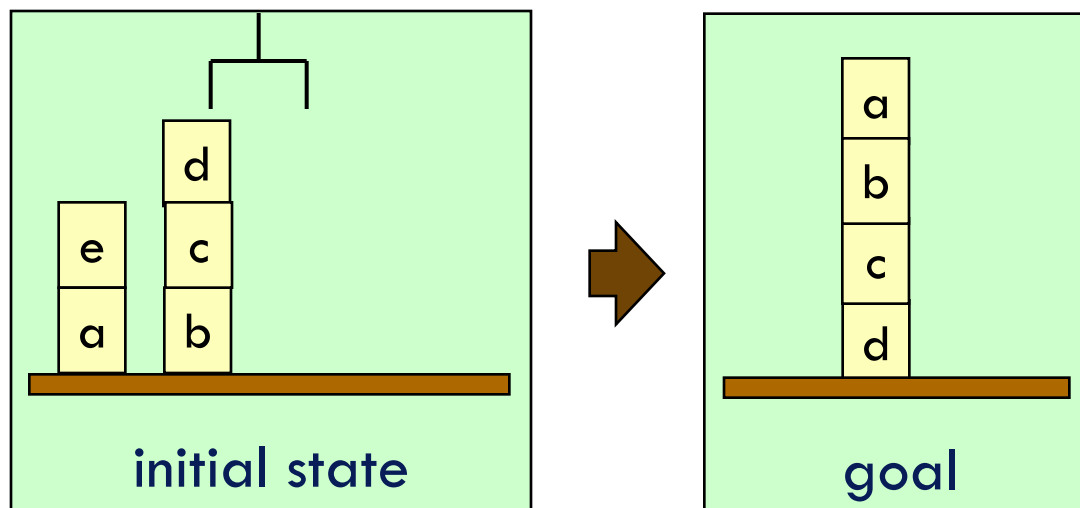
- Several ways:
 - ▣ Do something other than state-space search
 - ▣ Use forward or backward state-space search, with *domain-specific* knowledge to prune the search space
 - Can solve both problems quite easily this way
 - Example: block stacking using forward search

Domain-specific knowledge

- A blocks-world planning problem $P = (O, s_0, g)$ is solvable if s_0 and g satisfy some simple consistency conditions
 - g should not mention any blocks not mentioned in s_0
 - a block cannot be on two other blocks at once
- If P is solvable, can easily construct a solution of length $O(2m)$, where m is the number of blocks
 - ▣ Move all blocks to the table, then build up stacks from the bottom
 - Can do this in time $O(n)$
- With additional domain-specific knowledge can do even better ...

Additional Domain-Specific Knowledge

- A block x needs to be moved if any of the following is true:
 - s contains $ONtable(x)$ and g contains $ON(x,y)$ - see a below
 - s contains $ON(x,y)$ and g contains $ONtable(x)$ - see d below
 - s contains $ON(x,y)$ and g contains $ON(x,z)$ for some $y \neq z$, see C below
 - s contains $ON(x,y)$ and y needs to be moved - see e below



Domain – specific Algorithm

loop

if there is a clear block x such that

x needs to be moved **and**

x can be moved to a place where it won't need
to be moved

then move x to that place

else if there is a clear block x such that

x needs to be moved

then move x to the table

else if the goal is satisfied

then return the plan

else return failure

repeat

STRIPS Planning Task

- ▣ Monkey and Banana



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