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**Question 1.** (10 points)

Let  $h, h'$  be two propositional clauses or conjunctions. Show that  $\text{lgg}(h, h') = \text{Lits}(h) \cap \text{Lits}(h')$  is a least general generalization of  $h, h'$ .

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**Question 2.** (15 points)

Determine if

1.  $h \subseteq_{\theta} h'$
2.  $h' \models h$

for

$$h = p(x, y) \wedge p(y, z) \wedge \neg p(x, z)$$
$$h' = p(a, b) \wedge p(b, c) \wedge p(c, d) \wedge \neg p(a, d)$$

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**Question 3.** (5 points)

Consider the following statements

1.  $X =$  non-self-resolving FOL clauses
2.  $X =$  contingent FOL clauses
3. There is no  $k \in \mathbb{N}$ ,  $x \in X$  such that  $h_k \models x$  and  $h_k \not\subseteq_{\theta} x$ , where  $h_k$  ( $k \in \mathbb{N}$ ) are the hypotheses of the generalization algorithm.

Decide for each of the implications  $1 \rightarrow 2$ ,  $1 \rightarrow 3$ ,  $2 \rightarrow 3$ , whether it is true. Change the relation  $h_k \models x$  in (3) so that all the implications you decided true are true when (1) and (2) assume conjunctions instead of clauses.

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**Question 4.** (1 points)

Find two different least general generalizations of  $p(a)$  and  $p(b) \vee p(c)$ , prove that they are indeed generalizations of the two clauses and prove that they are mutually  $\theta$ -equivalent. Explain why two least general generalizations of the same pair of clauses or conjunctions must be  $\theta$ -equivalent.

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**Question 5.** (10 points)

Explain why the proof of the mistake bound  $n$  of the generalization algorithm is no longer valid when the assumption on  $X$  is changed to  $X =$  non-self-resolving FOL conjunctions or non-self-resolving FOL conjunctions clauses, the relations  $\subseteq, \subset$  are changed to  $\subseteq_{\theta}, \subset_{\theta}$  (respectively), and we set  $n = |\mathcal{P}|$ . Show that the proof cannot be rectified, in particular that no finite mistake bound exists under said assumption even if  $\mathcal{F} = \emptyset$ .