## Introduction to classical planning Problem representations

Michaela Urbanovská

PUI Tutorial Week 1

- Classical planning (M. Urbanovska), probabilistic planning (J. Mrkos)
- Each part one assignment 50 points total
- Classical planning assignment  $\rightarrow$  programming your own **planning** system
  - C, C++, Java
  - more details later
  - Deadline: half of April (TBD)
- Zapocet: 25 out of 50 points to pass
- Exam: 25 out of 50 points to pass
- Final grade: zapocet points + exam points
- Problems? Questions?
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- STRIPS, FDR problem definitions
- PDDL, compilations
- Relaxation heuristics
- Landmark heuristics
- Abstraction heuristics
- LP based heuristics
- Machine learning in planning
- …and much more!

## Lecture check

• Any questions regarding the lecture?

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# when your lecturer asks if you have any questions



- General problem solving
- Basically can solve all your problems

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- Problem

- General problem solving
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- Problem + representation

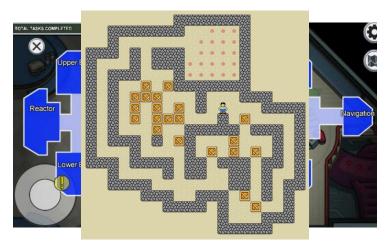
- Basically can solve all your problems
- Problem + representation + solver

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- Problem + representation + solver = **solution**

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#### • General problem solving

• Basically can solve all your problems

#### • Problem + representation + solver = solution



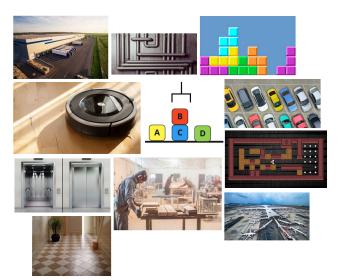
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## Planning benchmarks

For example...

- Airport
- Depot
- Sokoban
- Blocksworld
- Elevators
- Floortile
- Parking
- Pipesworld
- Tetris
- Tidybot
- Woodworking



- Often inspired by real-world problems
- Sometimes even modeled real-world problems
- Problems with interesting properties to test performance of algorithms heuristics
- Doesn't have to correlate with real-world performance

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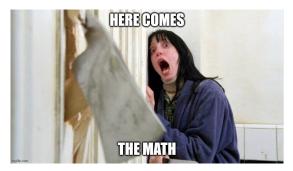
...this is not the class with robots.

There are many different types of planning... In this part of the course  $\rightarrow$  classical planning

- Fully defined environment
- Deterministic actions
- Domain-independence

- STRIPS
- FDR
- Specify the model, capture the structure of the problem

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### STRIPS problem $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$

• 
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 (facts)

• 
$$O = \{o_1, o_2, \dots o_m\}$$
 (operators)

•  $s_{init} \subseteq F$  (initial state)

•  $c(o_i) \in \mathbb{R}^+$  (cost function)

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#### STRIPS operator $o = \langle pre(o), add(o), del(o) \rangle$

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- $add(o) \subseteq F$  (set of add effects)
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- $pre(o) \subseteq F$  (set of preconditions)
- $add(o) \subseteq F$  (set of add effects)
- $del(o) \subseteq F$  (set of delete effects)
- operators are well-formed
  - $add(o) \cap del(o) = \emptyset$
  - $pre(o) \cap add(o) = \emptyset$

#### Applicable operator

Operator *o* is applicable in state *s* if  $pre(o) \subseteq s$ . **Resulting state**  $res(o, s) = (s \setminus del(o)) \cup add(o)$ . State *s* is a **goal state** iff  $s_{goal} \subseteq s$ .

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#### Sequence of applicable operators

**Sequence of operators**  $\pi = \langle o_1, o_2, \dots o_n \rangle$  is applicable in state  $s_0$  if there are states  $s_1, s_2, \dots s_n$  such that  $o_i$  is applicable in  $s_{i-1}$  and  $s_i = res(o_i, s_{i-1})$  for  $1 \le i \le n$ .

- $res(\pi, s_0) = s_n$  (result of the applied operator sequence  $\pi$ )
- $c(\pi) = \sum_{o \in \pi} c(o)$  (cost of applying the operator sequence  $\pi$ )

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- $c(\pi) = \sum_{o \in \pi} c(o)$  (cost of applying the operator sequence  $\pi$ ) Sequence  $\pi$  is called a **plan** if  $s_{goal} \subseteq res(\pi, s_{init})$ .
  - $\pi$  is an **optimal plan** is  $c(\pi)$  is the minimal cost over all plans

#### Reachable state

State *s* is **reachable** if there exists an applicable sequence of operators  $\pi$  such that  $res(\pi, s_{init} = s)$ . Set of all reachable states is denoted  $\mathcal{R}_{\Pi}$ .

## Problem definitions



# THERE'S MORE

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**PUI** Tutorial 1

#### FDR problem $P = \langle \mathcal{V}, \mathcal{O}, s_{init}, s_{goal}, c \rangle$

- $\mathcal{V} = \{V_1, V_2, \dots, V_n\}$  (finite set of variables)
- $\mathcal{O} = \{o_1, o_2, \dots o_m\}$  (set of operators)
- s<sub>init</sub> (initial state)
- s<sub>goal</sub> (goal state)
- $c(o_i) \in \mathbb{R}^+$

#### FDR problem $P = \langle \mathcal{V}, \mathcal{O}, s_{init}, s_{goal}, c \rangle$

- $\mathcal{V}$  (finite set of variables)
  - $V \in \mathcal{V}$  (variable)
  - $D_V$  (finite domain of variable V)
- s (state) is partial variable assignment over  ${\cal V}$ 
  - vars(s) =  $V \in \mathcal{V}$  assigned in s
  - s[V] = value of V in s
  - s is consistent with s' if s[V] = s'[V] for all V ∈ vars(s')
  - atom V = v is true in s if s[V] = v

#### FDR operator $o = \langle pre(o), eff(o) \rangle$

- $\mathcal{O}$  (set of operators)
  - pre(o) = partial assignment over V (preconditions)
  - eff(o) = partial assignment over V (effects)
  - V = v cannot be in both pre(o) and eff(o)

## Problem in FDR

#### Applicable operator

## Operator *o* is applicable in state *s* if pre(o) is **consistent** with *s*. **Resulting state** $res(o, s) = \begin{cases} eff(o)[V], & \text{if } V \in vars(eff(o))) \\ s[V], & \text{otherwise} \end{cases}$

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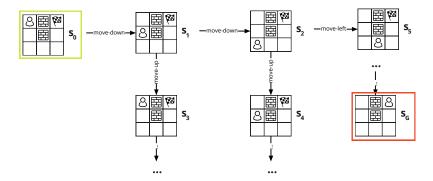
•  $c(\pi) = \sum_{o \in \pi} c(o)$  (cost of applying the operator sequence  $\pi$ )

Sequence  $\pi$  is called a **plan** if  $res(\pi, s_{init})$  is consistent with  $s_{goal}$ .

•  $\pi$  is an **optimal plan** is  $c(\pi)$  is the minimal cost over all plans

## Transition system

- Both STRIP and FDR have a notion of state and operator
  - $s_0$  is the initial state which gets expanded by using  $o \in O$  creating new state  $s' \rightarrow$  transition system



- Exhaustive search can be used but can be very slow
- Improvement? Heuristic function!
  - Additional information
  - Gives plan cost estimate for a state
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#### s-plan

Let  $\Pi$  denote a STRIPS planning task. Sequence of operators  $\pi$  is an **s-plan** iff  $\pi$  is applicable in *s* and  $res(\pi, s)$  is a goal state. Heuristic  $h : \mathcal{R}_{\Pi} \mapsto \mathbb{R} \cup \{\infty\}$  estimates costs of optimal s-plans. Heuristic function properties

- *h*\* **optimal** heuristic
- *h* is admissible iff  $h(s) \le h^*(s)$  for every  $s \in \mathcal{R}_{\Pi}$
- *h* is **goal-aware** iff  $h(s) \leq 0$  for every reachable goal state *s*
- *h* is safe iff  $h(s) = \infty$  implies  $h^* = \infty$  (there's no plan)
- h is consistent iff h(s) ≤ h(res(o, s)) + c(o) for all reachable states s ∈ R<sub>Π</sub> and o ∈ O applicable in s

Which statement holds?

- If *h* is both goal-aware and save, then *h* is admissible.
- If *h* is both goal-aware and consistent, then *h* is admissible.
- If *h* is both safe and consistent, then *h* is admissible.

#### Goals for today

• think that planning is useful

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 $\bullet$  think that planning is useful  $\checkmark$ 

#### Goals for today

- $\bullet$  think that planning is useful  $\checkmark$
- know STRIPS and FDR problem definition
- know what is a plan
- know what is a heuristic and why we need it
- know heuristic function properties

For more details and structure check out

• Notes on Classical Planning by Daniel Fiser

## The End



Feedback form

