Value Iteration and extensions

Jan Mrkos

PUI Tutorial Week 9

- Review of MDP concepts
- Value Iteration algorithm
- VI extensions

Value function of a policy

Look at the following definition of a value function of a policy for inifnite-horizon MDP. It contains multiple mistakes, correct them on a piece of paper:

Def: Value function of a policy for infinite-horizon MDP

Assume infinite horizon MDP with $\gamma \in [0, 100]$. Then let Value function of a policy π for every state $s \in S$ be defined as

$$V^{\pi}(s) = \sum_{s' \in S} R(s, \pi(s), s') R(s, \pi(s), s') + \gamma \pi(s')$$

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Basic algorithm for finding solution of Bellman Equations iteratively:

- **(**) initialize V_0 arbitrarily for each state, e.g to 0, set n = 0
- **2** Set n = n + 1.
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 V_n(s) = max_{a∈A} ∑_{s'∈S} T(s, a, s')[R(s, a, s') + γV_{n-1}(s')]
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Basic version uses 2 arrays to store state values.

Gridworld domain. Assuming $\gamma = 1$. Using VI and initializing $\forall s \ V_0(s) = 0$, calculate V_1, V_2, V_3 for the nine states around the +10 tile. Domain rules:

- Moving into edges gives -1 reward
- Moving onto marked tiles gives corresponding reward



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- GOTO 2.

Question: Does it converge? How fast? When do we stop?

Def: Residual

Residual of value function V_n *from* V_{n+1} *at state* $s \in S$ is defined by:

 $\operatorname{Res}^{V_n}(s) = |V_n(s) - V_{n+1}(s)|$

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Residual of value function V from V' is given by:

$$Res^{V_n} = ||V_n - V_{n+1}||_{\infty} = \max_{s} |V_n(s) - V_{n+1}(s)|$$

Stopping criterion: When residual of consecutive value functions is below low value of ϵ :

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However, this does not imply ϵ distance of value of greedy policy from optimal value function.

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In case of discounted ($\gamma < 1$) infinite-horizon MDPs:

$$|V_n, V^*$$
as above $\implies orall s \ |V_n(s) - V^*(s)| < 2rac{\epsilon \gamma}{1-\gamma}$

Algorithm 1: Value Iteration

1 initialize V_0 arbitrarily for each state, e.g to 0 while $Res^V > \epsilon$ do

- 2 pick some state *s*
- 3 Bellman backup $V(s) \leftarrow \max_{a \in A} \sum_{s' \in S} T(s, a, s') [R(s, a, s') + \gamma V(s')]$
- 4 Update residual at $s \operatorname{Res}^V(s) = |V_{old}(s) V_{new}(s)|$
- 5 return greedy policy π^V ;

Algorithm 2: Value Iteration

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Question: What is the greedy policy?

Algorithm 3: Value Iteration

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- 5 **return** greedy policy π^V ;

Question: What is the greedy policy?

• Greedy policy π_n^V is the policy given as argmax of V_n .

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Question: What are the memory requirements of VI?

• Value of each state needs to be stored twice

Another beautiful MDP example



• All undeclared rewards are -1

Task: Initialize VI with negative distance to S_5 and calculate first 3 iterations of VI, with state ordering S_0 to S_5

Algorithm 4: Asynchronous VI

- 1 initialize V_0 arbitrarily for each state, e.g to 0 while $Res^V > \epsilon$ do
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Algorithm 6: Asynchronous VI

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Algorithm 7: Asynchronous VI

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Question: Memory requirements compared to VI? Question: Convergence condition?

- \bullet Asymptotic as VI under condition that every state visited ∞ often.
- Question: How to pick s in 2.1?
 - Simplest is Gauss-Seidel VI, that is run AVI over all states iteratively

Build priority queue of states to update → Prioritized Sweeping VI. Update states in the order of the queue. Priority function:

$$\mathsf{priority}_{PS}(s) \leftarrow \max\{\mathsf{priority}_{PS}(s), \max_{a \in A} \{T(s, a, s') Res^V(s')\}\}$$

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EXAMPLE ON BOARD

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EXAMPLE ON BOARD

Convergence?

- If all states start with non-zero priority
- OR If you interleave regular VI sweeps with Prioritized VI

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Question: Why acyclic partitioning? EXAMPLE ON BOARD Thank you for participating in the tutorials :-)

Please fill out the feedback form \rightarrow



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