# Path Planning

Jan Faigl

Department of Computer Science
Faculty of Electrical Engineering
Czech Technical University in Prague

Lecture 03

Robotic Exploration and Data Collection Planning



#### Overview of the Lecture

- Part 1 Path Planning
  - Introduction to Path Planning
  - Notation and Terminology
  - Path Planning Methods
- Part 2 Grid and Graph based Path Planning Methods
  - Grid-based Planning
  - DT for Path Planning
  - Graph Search Algorithms
  - Path Planning based on Reaction-Diffusion Process



# Part I

Part 1 - Path and Motion Planning



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## Outline

Introduction to Path Planning

Notation and Terminology

Path Planning Methods



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# Robot Motion Planning - Motivational problem

■ How to transform high-level task specification (provided by humans) into a low-level description suitable for controlling the actuators?

To develop algorithms for such a transformation.

The motion planning algorithms provide transformations how to move a robot (object) considering all operational constraints

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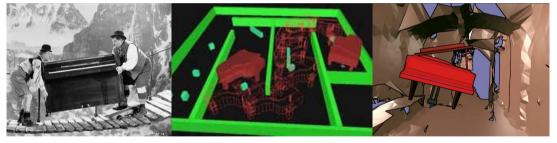
It encompasses several disciples, e.g., mathematics, robotics, computer science, control theory, artificial intelligence, computational geometry, etc.



## Piano Mover's Problem

#### A classical motion planning problem

Having a CAD model of the piano, model of the environment, the problem is how to move the piano from one place to another without hitting anything.



Basic motion planning algorithms are focused primarily on rotations and translations.

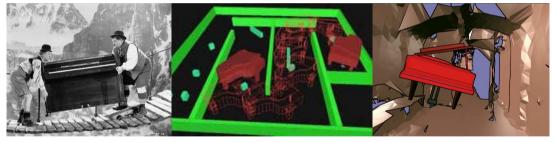
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- Moreover, we also need a context about the problem and realistic assumptions.

The plans have to be admissible and feasible

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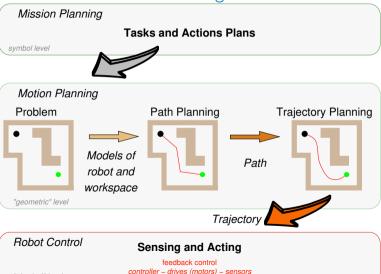
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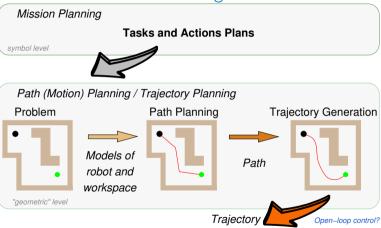
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"physical" level



Robot Control

#### Sensing and Acting

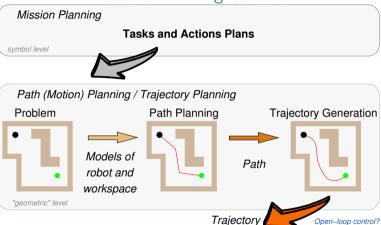
feedback control

controller – drives (motors) – sensors

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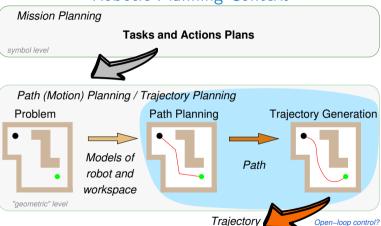
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feedback control
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Sources of uncertainties because of real environment



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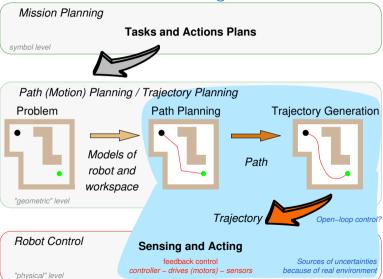
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Sources of uncertainties because of real environment







#### Real Mobile Robots

In a real deployment, the problem is more complex.

- The world is changing.
- Robots update the knowledge about the environment.

localization, mapping and navigation

 New decisions have to be made based on the feedback from the environment.

> Motion planning is a part of the mission replanning loop.



Josef Štrunc, Bachelor thesis, CTU, 2009.

An example of robotic mission:

Multi-robot exploration of unknown environment.

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laxing constraints and considering realistic assumptions.



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## Outline

Introduction to Path Planning

Notation and Terminology

■ Path Planning Methods



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■ W – World model describes the robot workspace and its boundary determines the obstacles  $O_i$ .

world,  $W = \mathbb{R}^2$ 

- A Robot is defined by its geometry, parameters (kinematics) and it is controllable by the motion plan.
- $\mathcal{C}$  Configuration space ( $\mathcal{C}$ -space)

A concept to describe possible configurations of the robot. The robot's configuration completely specify the robot location in  $\mathcal W$  including specification of all degrees of freedom.

E.g., a robot with rigid body in a plane 
$$C = \{x, y, \phi\} = \mathbb{R}^2 \times S^1$$

- Let  $\mathcal{A}$  be a subset of  $\mathcal{W}$  occupied by the robot,  $\mathcal{A} = \mathcal{A}(q)$ .
- $\blacksquare$  A subset of  $\mathcal{C}$  occupied by obstacles is

$$C_{abs} = \{ a \in \mathcal{C} : \mathcal{A}(a) \cap \mathcal{O}_i, \forall i \}$$

Collision-free configurations are

$$C_{free} = C \setminus C_{obs}$$



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# Path / Motion Planning Problem

■ Path is a continuous mapping in C-space such that

$$\pi: [0,1] \to \mathcal{C}_{free}$$
, with  $\pi(0) = q_0$ , and  $\pi(1) = q_f$ .

■ **Trajectory** is a path with explicit parametrization of time, e.g., accompanied by a description of the motion laws ( $\gamma: [0,1] \to \mathcal{U}$ , where  $\mathcal{U}$  is robot's action space).

$$[T_0, T_f] \ni \mathbf{t} \leadsto \mathbf{\tau} \in [0, 1] : q(\mathbf{t}) = \pi(\mathbf{\tau}) \in \mathcal{C}_{free}$$

The path planning is the determination of the function  $\pi(\cdot)$ .

Additional requirements can be given:

- Smoothness of the path;
- Kinodynamic constraints, e.g., considering friction forces;
- Optimality criterion shortest vs fastest (length vs curvature).
- Path planning planning a collision-free path in *C*-space.
- Motion planning planning collision-free motion in the state space.



It includes dynamics.

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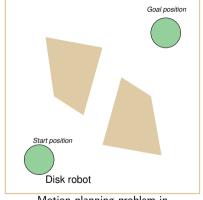


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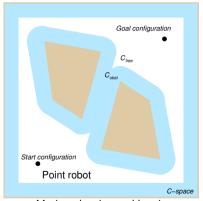
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# Planning in C-space

Robot path planning for a disk-shaped robot with a radius  $\rho$ .



Motion planning problem in geometrical representation of  $\mathcal{W}$ .



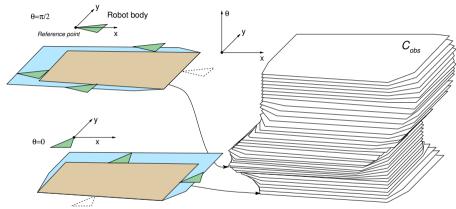
Motion planning problem in C-space representation.

 $\mathcal{C}$ -space has been obtained by enlarging obstacles by the disk  $\mathcal{A}$  with the radius  $\rho$ .

By applying Minkowski sum:  $\mathcal{O} \oplus \mathcal{A} = \{x + y \mid x \in \mathcal{O}, y \in \mathcal{A}\}.$ 



# Example of $C_{obs}$ for a Robot with Rotation



A simple 2D obstacle ightarrow has a complicated  $\mathcal{C}_{obs}.$ 

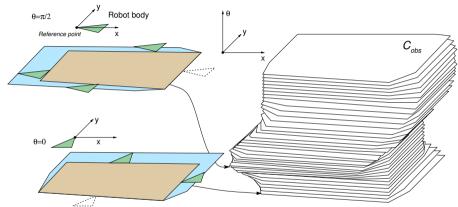
Deterministic algorithms exist.

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**Explicit** representation of  $C_{free}$  is impractical to compute.



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# Representation of C-space

How to deal with continuous representation of C-space?

#### Continuous Representation of C-space

#### Discretization

processing critical geometric events, (random) sampling roadmaps, cell decomposition, potential field

# **Graph Search Techniques**BFS, Gradient Search, A\*



## Outline

Introduction to Path Planning

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Path Planning Methods



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# Planning Methods - Overview (selected approaches)

■ Point-to-point path/motion planning.

Multi-goal path/motion/trajectory planning later

- Roadmap based methods Create a connectivity graph of the free space.
  - Visibility graph;

(Complete but impractical)

- Cell decomposition;
- Voronoi graph.
- Discretization into a grid-based (or lattice-based) representation (Resolution complete)
- Potential field methods (Complete

(Complete only for a "navigation function", which is hard to compute in general)

Classic path planning algorithms

#### Randomized sampling-based methods

- Creates a roadmap from connected random samples in  $C_{free}$ .
- Probabilistic roadmaps.

Samples are drawn from some distribution.



Very successful in practice.

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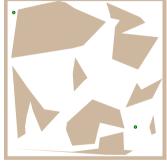
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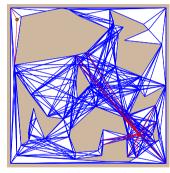
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- Compute visibility graph.
- 2. Find the shortest path.

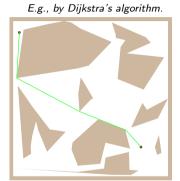


Problem

# Visibility Graph



Visibility graph

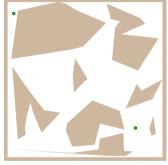


Found shortest path

- Naïve all segments between n vertices of the map  $O(n^3)$ ;
- Using rotation trees for a set of segments  $-O(n^2)$ . M. H. Overmars and E. Welzl, 1988

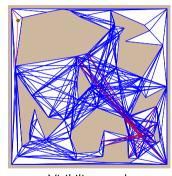


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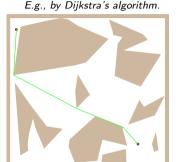


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Visibility graph



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Constructions of the visibility graph:

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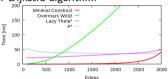
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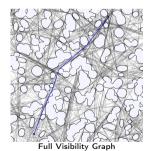


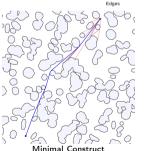
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#### Minimal Construct: Efficent Shortest Path in Polygonal Maps

- Minimal Construct algorithm computes visibility graph during the A\* search instead of first computation of the complete visibility graph and then finding the shortest path using A\* or Dijkstra algorithm.
- Based on A\* search with line intersection tests are delayed until they become necessary.
- The intersection tests are further accelerated using bounding boxes.



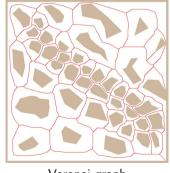




Marcell Missura, Daniel D. Lee, and Maren Bennewitz (2018): Minimal Construct: Efficient Shortest Path Finding for Mobile Robots in Polygonal Maps. IROS.

# Voronoi Graph

- 1. Roadmap is Voronoi graph that maximizes clearance from the obstacles.
- 2. Start and goal positions are connected to the graph.
- 3. Path is found using a graph search algorithm.





Voronoi graph

Path in graph

Found path



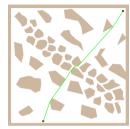
# Visibility Graph vs Voronoi Graph

## Visibility graph

Shortest path, but it is close to obstacles. We have to consider safety of the path.

An error in plan execution can lead to a collision.

Complicated in higher dimensions



## Voronoi graph

- It maximize clearance, which can provide conservative paths.
- Small changes in obstacles can lead to large changes in the graph.
- Complicated in higher dimensions.

A combination is called Visibility-Voronoi – R. Wein, J. P. van den Berg, D. Halperin, 2004.



For higher dimensions we need other types of roadmaps.

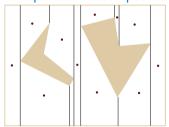
#### Cell Decomposition

1. Decompose free space into parts.

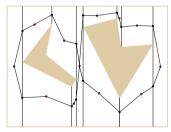
Any two points in a convex region can be directly connected by a segment.

- 2. Create an adjacency graph representing the connectivity of the free space.
- 3. Find a path in the graph.

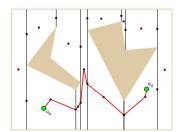
#### Trapezoidal decomposition



Centroids represent cells



Connect adjacency cells



Find path in the adjacency graph

Other decomposition (e.g., triangulation) are possible.



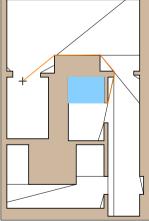
# Shortest Path Map (SPM)

Speedup computation of the shortest path towards a particular goal location  $p_g$  for a polygonal domain  $\mathcal{P}$  with n vertices.

- A partitioning of the free space into cells with respect to the particular location  $p_g$ .
- **Each** cell has a vertex on the shortest path to  $p_g$ .
- Shortest path from any point p is found by determining the cell (in  $O(\log n)$  using point location alg.) and then travesing the shortest path with up to k bends, i.e., it is found in  $O(\log n + k)$ .
- Determining the SPM using "wavefront" propagation based on continuous Dijkstra paradigm.

Joseph S. B. Mitchell: A new algorithm for shortest paths among obstacles in the plane, Annals of Mathematics and Artificial Intelligence, 3(1):83–105, 1991.

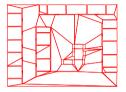
- SPM is a precompute structure for the given  $\mathcal{P}$  and  $p_g$ ;
  - single-point query.

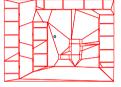


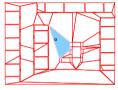
A similar structure can be found for two-point query, e.g., H. Guo, A. Maheshwari, J.-R. Sack, 2008.

#### Point Location Problem

• For a given partitioning of the polygonal domain into a discrete set of cells, determine the cell for a given point *p*.







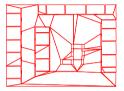
Masato Edahiro, Iwao Kokubo and Takao Asano: A new point-location algorithm and its practical efficiency: comparison with existing algorithms, ACM Trans. Graph., 3(2):86–109, 1984.

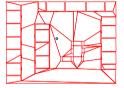
It can be implemented using interval trees – slabs and slices.

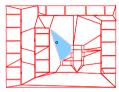


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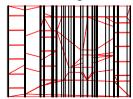






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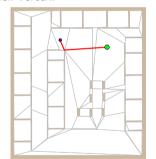


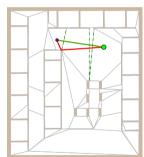


Point location problem, SPM and similarly problems are from the Computational Geometry field.

#### Approximate Shortest Path and Navigation Mesh

- We can use any convex partitioning of the polygonal map to speed up shortest path queries.
  - 1. Precompute all shortest paths from map vertices to  $p_g$  using visibility graph.
  - 2. Then, an estimation of the shortest path from p to  $p_g$  is the shortest path among the one of the cell vertex.

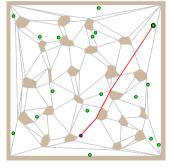


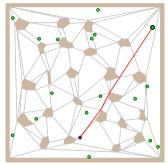


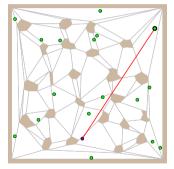
■ The estimation can be further improved by "ray-shooting" technique combined with walking in triangulation (convex partitioning). (Faigl, 2010)

#### Path Refinement

- Testing collision of the point p with particular vertices of the estimation of the shortest path.
  - Let the initial path estimation from p to  $p_g$  be a sequence of k vertices  $(p, v_1, \dots, v_k, p_g)$ .
  - We can iteratively test if the segment  $(p, v_i)$ ,  $1 < i \le k$  is collision free up to  $(p, p_g)$ .







Path over  $v_0$ 

Path over  $v_1$ 

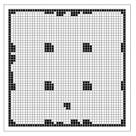
Full refinement

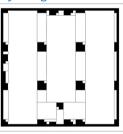
With the precomputed structures, an estimate of the shortest path is determined in units of microseconds.

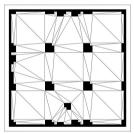
#### Navigation Mesh

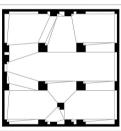
- In addition to robotic approaches, fast shortest path queries are studied in computer games.
- There is a class of algorithms based on navigation mesh.
  - A supporting structure representing the free space.

It usually originated from the grid based maps, but it is represented as CDT – Constrained Delaunay triangulation.









Grid mesh

Merged grid mesh

CDT mesh

Merged CDT mesh

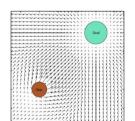
■ E.g., Polyanya algorithm based on navigation mesh and best-first search.

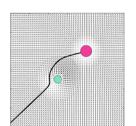
M. Cui, D. Harabor, A. Grastien: Compromise-free Pathfinding on a Navigation Mesh, IJCAI 2017, 496–502. https://bitbucket.org/dharabor/pathfinding

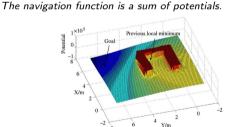


#### Artificial Potential Field Method

- The idea is to create a function f that will provide a direction towards the goal for any configuration of the robot.
- Such a function is called navigation function and  $-\nabla f(q)$  points to the goal.
- Create a potential field that will attract robot towards the goal  $q_f$  while obstacles will generate repulsive potential repelling the robot away from the obstacles.







Such a potential function can have several local minima.



#### Avoiding Local Minima in Artificial Potential Field

Consider harmonic functions that have only one extremum

$$\nabla^2 f(q) = 0.$$

Finite element method with defined Dirichlet and Neumann boundary conditions.







J. Mačák, Master thesis, CTU, 2009

# Part II

Part 2 - Grid and Graph based Path Planning Methods



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#### Outline

- Grid-based Planning
- DT for Path Planning
- Graph Search Algorithms
- Path Planning based on Reaction-Diffusion Process



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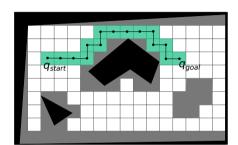
### Grid-based Planning

- A subdivision of  $C_{free}$  into smaller cells.
- **Grow obstacles** can be simplified by growing borders by a diameter of the robot.
- Construction of the planning graph G = (V, E) for V as a set of cells and E as the neighbor-relations.

4-neighbors and 8-neighbors



A grid map can be constructed from the so-called occupancy grid maps.
 E.g., using thresholding.



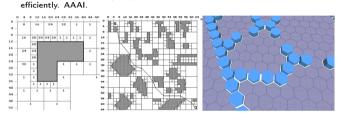


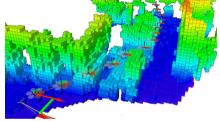


Grid-based Planning DT for Path Planning Graph Search Algorithms RD-based Planning

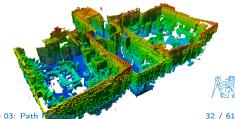
#### Grid-based Environment Representations

Hiearchical planning with coarse resolution and re-planning on finer resolution.
 Holte, R. C. et al. (1996): Hierarchical A \*: searching abstraction hierarchies



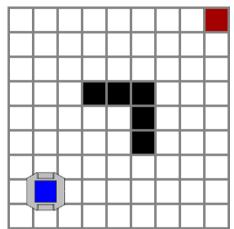


- Octree can be used for the map representation.
- In addition to squared (or rectangular) grid a hexagonal grid can be used.
- 3D grid maps OctoMap https://octomap.github.io.
- Memory grows with the size of the environment.
- Due to limited resolution it may fail in narrow passages of  $C_{free}$ .



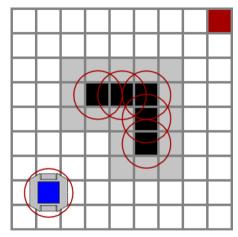
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- Wave-front propagation using path simplication
- Initial map with a robot and goal.
- Obstacle growing.
- Wave-front propagation "flood fill".
- Find a path using a navigation function.
- Path simplification.
  - "Ray-shooting" technique combined with Bresenham's line algorithm.
  - The path is a sequence of "key" cells for avoiding obstacles.





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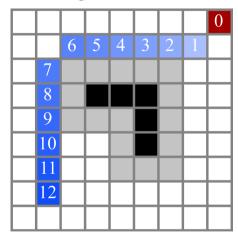


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ı	8	7	6	5	4	3	2	1	0
	8	7	6	5	4	3	2	1	1
	8	7						2	2
	8	8						3	3
	9	9						4	4
	10	10	10	10				5	5
	11	11	11	10				6	6
		12	11	10	9	8	7	7	7
			11	10	9	8	8	8	8

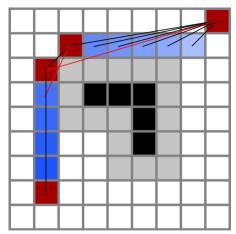


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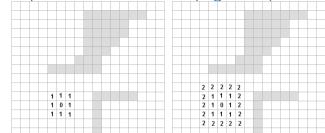


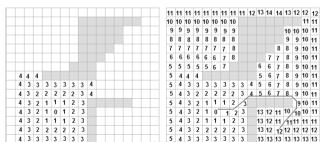
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  - The path is a sequence of "key" cells for avoiding obstacles.





#### Example – Wave-Front Propagation (Flood Fill)

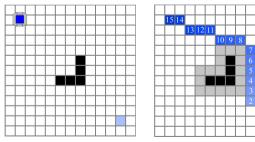


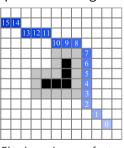


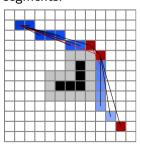


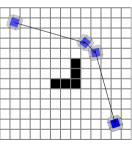
# Path Simplification

- The initial path is found in a grid using 8-neighborhood.
- The rayshoot cast a line into a grid and possible collisions of the robot with obstacles are checked.
- The "farthest" cells without collisions are used as "turn" points.
- The final path is a sequence of straight line segments.









Obtacle growing, wave-front propagation

Ray-shooting

Simplified path

Initial and goal locations

- Filling a grid by a line with avoding float numbers.

```
• A line from (x_0, y_0) to (x_1, y_1) is given by y = \frac{y_1 - y_0}{x_1 - x_0}(x - x_0) + y_0.
     CoordsVector& bresenham(const Coords& pt1, const Coords& pt2,
                                                                               int twoDy = 2 * dv:
            CoordsVector& line)
                                                                               int twoDvTwoDx = twoDv - 2 * dx: //2*Dv - 2*Dx
                                                                       28
                                                                               int e = twoDv - dx: //2*Dv - Dx
        // The pt2 point is not added into line
                                                                       29
                                                                                int v = v0;
        int x0 = pt1.c; int y0 = pt1.r;
                                                                       30
                                                                                int xDraw, yDraw;
        int x1 = pt2.c; int v1 = pt2.r;
                                                                       31
                                                                                for (int x = x0: x != x1: x += xstep) {
 6
        Coords p;
                                                                       32
                                                                                   if (steep) {
                                                                       33
        int dx = x1 - x0:
                                                                                      xDraw = v;
        int dv = v1 - v0:
                                                                       34
                                                                                      vDraw = x:
        int steep = (abs(dy) >= abs(dx));
                                                                       35
                                                                                   } else {
10
                                                                       36
        if (steep) {
                                                                                      xDraw = x:
11
           SWAP(x0. v0):
                                                                       37
                                                                                      vDraw = v:
12
           SWAP(x1, y1);
                                                                       38
13
           dx = x1 - x0; // recompute Dx, Dy
                                                                       39
                                                                                   p.c = xDraw:
14
           dy = y1 - y0:
                                                                       40
                                                                                   p.r = vDraw:
15
                                                                       41
                                                                                   line.push_back(p); // add to the line
16
                                                                       42
                                                                                   if (e > 0) {
        int xstep = 1:
17
                                                                                      e += twoDvTwoDx: //E += 2*Dv - 2*Dx
        if (dx < 0) {
                                                                       43
18
           xstep = -1:
                                                                       44
                                                                                      y = y + ystep;
19
           dx = -dx:
                                                                       45
                                                                                  } else {
20
                                                                       46
                                                                                      e += twoDv: //E += 2*Dv
21
        int ystep = 1;
                                                                       47
22
        if (dv < 0) {
                                                                       48
23
           vstep = -1:
                                                                       49
                                                                                return line:
24
                                                                       50
           dv = -dv:
25
```



Grid-based Planning

#### Outline

- Grid-based Planning
- DT for Path Planning
- Graph Search Algorithms
- Path Planning based on Reaction-Diffusion Process



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### Distance Transform based Path Planning

- For a given goal location and grid map compute a navigational function using wave-front algorithm, i.e., a kind of potential field.
  - The value of the goal cell is set to 0 and all other free cells are set to some very high value.
  - For each free cell compute a number of cells towards the goal cell.
  - It uses 8-neighbors and distance is the Euclidean distance of the centers of two cells, i.e., EV=1 for orthogonal cells or  $EV = \sqrt{2}$  for diagonal cells.
  - The values are iteratively computed until the values are changing.
  - The value of the cell c is computed as

$$cost(c) = \min_{i=1}^{8} \left( cost(c_i) + EV_{c_i,c} \right),$$

where  $c_i$  is one of the neighboring cells from 8-neighborhood of the cell c.

- The algorithm provides a cost map of the path distance from any free cell to the goal cell.
- The path is then used following the gradient of the cell cost.

Jarvis, R. (2004): Distance Transform Based Visibility Measures for Covert Path Planning in Known but Dynamic Environments.

# Distance Transform Path Planning

#### Algorithm 1: Distance Transform for Path Planning

```
for y := 0 to yMax do
    for x := 0 to xMax do
         if goal [x,v] then
             cell [x,y] := 0;
         else
             cell [x,y] := xMax * yMax; //initialization, e.g., pragmatic of the use longest distance as \infty ;
repeat
    for v := 1 to (vMax - 1) do
         for x := 1 to (xMax - 1) do
             if not blocked [x,y] then | cell [x,y] := cost(x, y);
    for y := (yMax-1) downto 1 do
         for x := (xMax-1) downto 1 do
```



until no change;

Grid-based Planning

# Distance Transform based Path Planning – Impl. 1/2

```
Grid& DT::compute(Grid& grid) const
                                                                                           for (int r = H - 2; r > 0; --r) {
      2
                                                                             36
                                                                                           for (int c = W - 2; c > 0; --c) {
              static const double DIAGONAL = sqrt(2);
                                                                                               if (map[r][c] != FREESPACE) {
              static const double ORTOGONAL = 1:
                                                                             38
                                                                                                  continue:
              const int H = map.H:
                                                                             39
                                                                                               } //obstacle detected
              const int W = map.W:
                                                                                               double t[4]:
              assert(grid.H == H and grid.W == W. "size");
                                                                             41
                                                                                               t[1] = grid[r + 1][c] + ORTOGONAL:
              bool anyChange = true;
                                                                                               t[0] = grid[r + 1][c + 1] + DIAGONAL;
              int counter = 0:
                                                                                               t[3] = grid[r][c + 1] + ORTOGONAL:
     10
              while (anyChange) {
                                                                                               t[2] = grid[r + 1][c - 1] + DIAGONAL:
                                                                             44
     11
                 anyChange = false;
                                                                             45
                                                                                               double pom = grid[r][c];
     12
                 for (int r = 1: r < H - 1: ++r) {
                                                                             46
                                                                                               bool s = false:
     13
                    for (int c = 1: c < W - 1: ++c) {
                                                                                               for (int i = 0: i < 4: i++) {
     14
                       if (map[r][c] != FREESPACE) {
                                                                             48
                                                                                                  if (pom > t[i]) {
     15
                                                                                                     pom = t[i];
                          continue:
     16
                       } //obstacle detected
                                                                                                     s = true:
     17
                       double t[4];
                                                                             51
     18
                       t[0] = grid[r - 1][c - 1] + DIAGONAL:
                                                                             52
     19
                      t[1] = grid[r - 1][c] + ORTOGONAL;
                                                                             53
                                                                                               if (s) {
     20
                       t[2] = grid[r - 1][c + 1] + DIAGONAL:
                                                                             54
                                                                                                  anyChange = true;
     21
                      t[3] = grid[r][c - 1] + ORTOGONAL;
                                                                             55
                                                                                                  grid[r][c] = pom;
                      double pom = grid[r][c];
     23
                       for (int i = 0: i < 4: i++) {
                                                                             57
                          if (pom > t[i]) {
     24
                                                                             58
     25
                             pom = t[i];
                                                                             59
                                                                                        counter++;
     26
                             anvChange = true;
                                                                             60
                                                                                      } //end while any change
     27
                                                                             61
                                                                                      return grid;
     28
                                                                             62
                       if (anyChange) {
                                                                          A boundary is assumed around the rectangular map
     30
                          grid[r][c] = pom:
     31
     32
     33
Jan Faigl, 2022
```

# Distance Transform based Path Planning – Impl. 2/2

■ The path is retrived by following the minimal value towards the goal using min8Point().

```
Coords& min8Point(const Grid& grid, Coords& p)
         double min = std::numeric limits<double>::max():
         const int H = grid.H;
         const int W = grid.W;
         Coords t:
         for (int r = p.r - 1; r \le p.r + 1; r++) {
            if (r < 0 \text{ or } r >= H) \{ \text{ continue}; \}
            for (int c = p.c - 1; c \le p.c + 1; c++) {
10
11
                if (c < 0 \text{ or } c >= W) \{ \text{ continue; } \}
12
                if (min > grid[r][c]) {
13
                   min = grid[r][c]:
14
                   t.r = r; t.c = c;
15
16
17
18
         p = t;
19
         return p:
20
```

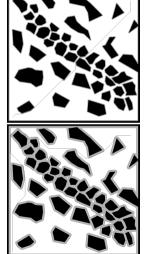
```
CoordsVector& DT::findPath(const Coords& start, const Coords&
           goal, CoordsVector& path)
23
     ł
24
        static const double DIAGONAL = sqrt(2);
25
        static const double ORTOGONAL = 1:
26
        const int H = map.H:
27
        const int W = map.W:
28
        Grid grid(H, W, H*W); // H*W max grid value
29
        grid[goal.r][goal.c] = 0;
30
        compute(grid):
31
32
        if (grid[start.r][start.c] >= H*W) {
33
           WARN("Path has not been found"):
34
        } else {
35
           Coords pt = start:
36
           while (pt.r != goal.r or pt.c != goal.c) {
37
              path.push_back(pt);
              min8Point(grid, pt);
39
40
           path.push_back(goal);
41
42
        return path;
43
```

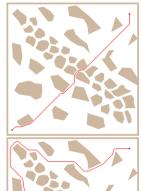


### DT Example

lacksquare  $\delta=$  10 cm, L= 27.2 m

•  $\delta = 30$  cm, L = 42.8 m









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#### Outline

- Grid-based Planning
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# Graph Search Algorithms

- The grid can be considered as a graph and the path can be found using graph search algorithms.
- The search algorithms working on a graph are of general use, e.g.,
  - Breadth-first search (BFS);
    - Depth first search (DFS);
  - Dijktra's algorithm;
  - A\* algorithm and its variants.
- There can be grid based speedups techniques, e.g.,
  - Jump Search Algorithm (JPS) and JPS<sup>+</sup>.
- There are many search algorithms for on-line search, incremental search and with any-time and real-time properties, e.g.,
  - Lifelong Planning A\* (LPA\*).

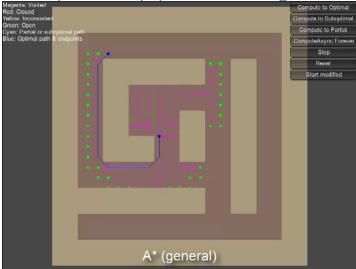
Koenig, S., Likhachev, M. and Furcy, D. (2004): Lifelong Planning A\*. AlJ.

■ E-Graphs — Experience graphs

Phillips, M. et al. (2012): E-Graphs: Bootstrapping Planning with Experience Graphs. RSS.



# Examples of Graph/Grid Search Algorithms





#### A\* Algorithm

• A\* uses a user-defined h-values (heuristic) to focus the search.

Peter Hart, Nils Nilsson, and Bertram Raphael, 1968

• Prefer expansion of the node *n* with the lowest value

$$f(n) = g(n) + h(n),$$

where g(n) is the cost (path length) from the start to n and h(n) is the estimated cost from n to the goal.

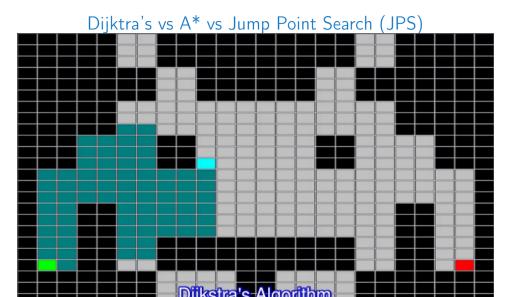
- h-values approximate the goal distance from particular nodes.
- Admissibility condition heuristic always underestimate the remaining cost to reach the goal.
  - Let  $h^*(n)$  be the true cost of the optimal path from n to the goal.
  - Then h(n) is admissible if for all n:  $h(n) \le h^*(n)$ .
  - E.g., Euclidean distance is admissible.
    - A straight line will always be the shortest path.
- Dijkstra's algorithm h(n) = 0.



### A\* Implementation Notes

- The most costly operations of A\* are:
  - Insert and lookup an element in the closed list;
  - Insert element and get minimal element (according to f() value) from the open list.
- The closed list can be efficiently implemented as a hash set.
- The open list is usually implemented as a priority queue, e.g.,
  - Fibonacii heap, binomial heap, k-level bucket;
  - **binary heap** is usually sufficient with O(logn).
- Forward A\*
  - 1. Create a search tree and initiate it with the start location.
  - 2. Select generated but not yet expanded state s with the smallest f-value, f(s) = g(s) + h(s).
  - 3. Stop if s is the goal.
  - 4. Expand the state s.
  - 5. Goto Step 2.

Similar to Dijktra's algorithm but it uses f(s) with the heuristic h(s) instead of pure g(s).



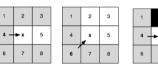


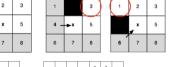
### Jump Point Search Algorithm for Grid-based Path Planning

Jump Point Search (JPS) algorithm is based on a macro operator that identifies and selectively expands only certain nodes (jump points).

Harabor, D. and Grastien, A. (2011): Online Graph Pruning for Pathfinding on Grid Maps. AAAI.

- Natural neighbors after neighbor prunning with forced neighbors because of obstacle.
- Intermediate nodes on a path connecting two jump points are never expanded.









No preprocessing and no memory overheads while it speeds up A\*.

https://harablog.wordpress.com/2011/09/07/jump-point-search/

■ JPS<sup>+</sup> is optimized preprocessed version of JPS with goal bounding.

https://github.com/SteveRabin/JPSPlusWithGoalBounding

http://www.gdcvault.com/play/1022094/JPS-Over-100x-Faster-than



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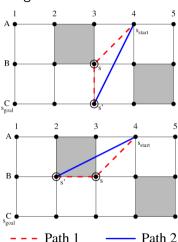
### Theta\* – Any-Angle Path Planning Algorithm

- Any-angle path planning algorithms simplify the path during the search.
- Theta\* is an extension of A\* with LineOfSight().
  Nash, A., Daniel, K., Koenig, S. and Felner, A. (2007): Theta\*: Any-Angle Path Planning on Grids. AAAI.

#### **Algorithm 2:** Theta\* Any-Angle Planning

 Path 2: considers path from start to parent(s) and from parent(s) to s' if s' has line-of-sight to parent(s).

http://aigamedev.com/open/tutorials/theta-star-any-angle-paths/



### Theta\* Any-Angle Path Planning Examples

■ Example of found paths by the Theta\* algorithm for the same problems as for the DT-based examples on Slide 42.







 $\delta = 30$  cm. L = 40.3 m

The same path planning problems solved by DT (without path smoothing) have  $L_{\delta=10}=27.2$  m and  $L_{\delta=30}=42.8$  m, while DT seems to be significantly faster.

■ Lazy Theta\* — reduces the number of line-of-sight checks.

Nash, A., Koenig, S. and Tovey, C. (2010): Lazy Theta\*: Any-Angle Path Planning and Path Length Analysis in 3D. AAAI.

http://aigamedev.com/open/tutorial/lazy-theta-star/

#### Outline

- Grid-based Planning
- DT for Path Planning
- Graph Search Algorithms
- Path Planning based on Reaction-Diffusion Process



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#### Reaction-Diffusion Processes Background

- Reaction-Diffusion (RD) models dynamical systems capable to reproduce the autowayes.
- Autowaves a class of nonlinear waves that propagate through an active media.
   At the expense of the energy stored in the medium, e.g., grass combustion.
- lacksquare RD model describes spatio-temporal evolution of two state variables  $u=u(ec{x},t)$  and

$$v = v(\vec{x},t)$$
 in space  $\vec{x}$  and time  $t$  
$$\dot{u} = f(u,v) + D_u \triangle u$$
 
$$\dot{v} = g(u,v) + D_v \triangle v$$

where  $\triangle$  is the Laplacian.

This RD-based path planning is informative, just for curiosity.



#### Reaction-Diffusion Background

FitzHugh-Nagumo (FHN) model

$$\dot{u} = \varepsilon \left( u - u^3 - v + \phi \right) + D_u \triangle u$$

$$\dot{v} = \left( u - \alpha v + \beta \right) + D_v \triangle u$$

where  $\alpha, \beta, \epsilon$ , and  $\phi$  are parameters of the model.

■ Dynamics of RD system is determined by the associated *nullcline configurations* for  $\dot{u}$ =0 and  $\dot{v}$ =0 in the absence of diffusion, i.e.,

$$\varepsilon (u - u^3 - v + \phi) = 0,$$
  

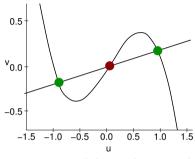
$$(u - \alpha v + \beta) = 0,$$

which have associated geometrical shapes.



Grid-based Planning DT for Path Planning Graph Search Algorithms RD-based Planning

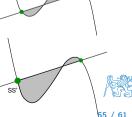
### Nullcline Configurations and Steady States



- Nullclines intersections represent:
  - Stable States (SSs):
  - Unstable States.
  - Bistable regime

The system (concentration levels of (u, v) for each grid cell) tends to be in SSs.

- We can modulate relative stability of both SS.
  - "preference" of SS+ over SS-
- System moves from  $SS^-$  to  $SS^+$ , if a small perturbation is introduced.
- The SSs are separated by a mobile frontier a kind of traveling frontwave (autowaves).



### RD-based Path Planning - Computational Model

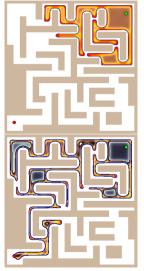
- Finite difference method on a Cartesian grid with Dirichlet boundary conditions (FTCS).
  discretization → grid based computation → grid map
- External forcing introducing additional information
   i.e., constraining concentration levels to some specific values.
- Two-phase evolution of the underlying RD model.
  - 1. Propagation phase
    - Freespace is set to  $SS^-$  and the start location  $SS^+$ .
    - Parallel propagation of the frontwave with non-annihilation property.

Vázquez-Otero and Muñuzuri, CNNA (2010)

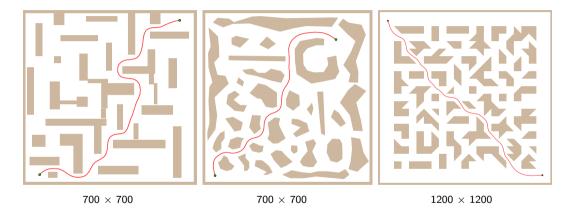
Terminate when the frontwave reaches the goal.

#### 2. Contraction phase

- Different nullclines configuration.
- Start and goal positions are forced towards SS<sup>+</sup>.
- SS<sup>-</sup> shrinks until only the path linking the forced points remains.



#### Example of Found Paths

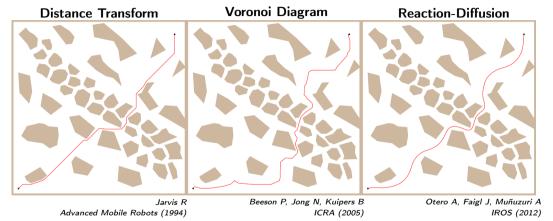


■ The path clearance maybe adjusted by the wavelength and size of the computational grid.

Control of the path distance from the obstacles (path safety).



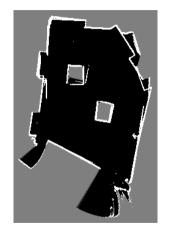
### Comparison with Standard Approaches



■ RD-based approach provides competitive paths regarding path length and clearance, while they seem to be smooth.



#### Robustness to Noisy Data





Vázquez-Otero, A., Faigl, J., Duro, N. and Dormido, R. (2014): Reaction-Diffusion based Computational Model for Autonomous Mobile Robot Exploration of Unknown Environments. International Journal of Unconventional Computing (IJUC).



# Summary of the Lecture



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### Topics Discussed

- Motion and path planning problems
  - Path planning methods overview
  - Notation of configuration space
- Path planning methods for geometrical map representation
  - Shortest-Path Roadmaps
    - Voronoi diagram based planning
    - Cell decomposition method
- Distance transform can be utilized for kind of *navigational function* 
  - Front-Wave propagation and path simplification
- Artificial potential field method
- Graph search (planning) methods for grid-like representation
  - Dijsktra's, A\*, JPS, Theta\*
  - Dedicated speed up techniques can be employed to decreasing computational burden, e.g., JPS
  - Grid-path can be smoothed, e.g., using path simplification or Theta\* like algorithms
- Unconventional reaction-diffusion based planning (informative)
- Next: Robotic Information Gathering Mobile Robot Exploration

