Lecture 4: Graphs, graph representation, basic graph processing

Robert Pěnička, Daniel Průša

Graph

Graph representations

# Lecture 4: Graphs, graph representation, basic graph processing BE5B33ALG — Algorithms

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### Introduction

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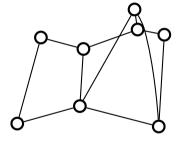
Graphs

Graph representations

- Graph G = (V, E)
- *V* is a set of nodes (vertices)
- ullet E is a set of edges (connections between nodes)
- Node (Vertex)

0

• Edge



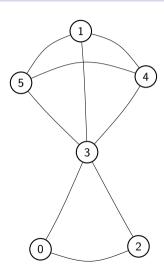
# Graphs: Nodes and Edges

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Graphs

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### □ Nodes, Vertices

- Servers, cities...
- Persons, people...
- Objects in comp. science
- ... etc.

### □ Edges

- Connections, roads...
- Personal relations
- Relations among objects
- ... etc.

# Usual graph representations

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Graph representation

### Node degrees & Lists of neighbours

- Less obvious, more effective
- 1D/2D array, vector, ArrayList...

Node indices	Degree	Neighbours indices
0	2	2, 3
1	3	3, 4, 5
2	2	0, 3
3	5	0, 2, 1, 4, 5
4	3	1, 3, 5
5	3	1, 3, 4

### **Adjacency matrix**

- Plain, obvious, less effective
- 2D array, matrix

Node indices	0	1	2		4	5
0	0	0	1	1	0	0
1	0	0	0	1	1	1
2	1	0	0	1	0	0
3	1	1	1	0	1	1
4	0	1	0	1	0	1
5	0	1	0	1 1 0 1	1	0

# Small graph zoo

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# Graphs

### Connected



# Disconnected



### Cycle

ullet N nodes and Nedges



### Path

- N nodes and N-1edges

### Tree

- Connected, no cycles
- ullet N nodes and N-1 edges



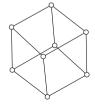
### Complete

- Connected
- N nodes and  $(N^2 N)/2$ edges



### Regular

• All node degrees are the same



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### Weighted graph

Each edge has its cost (length, weight)



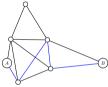
### Cycle in a graph

 Path whose first and last node are the same



### Path between A and B

Path visits each node at most once



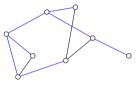
Bipartite graph

 Two-colorable; cycles only of even length; no edges inside partitions



### Spanning tree

Subgraph that is a tree and contains all nodes



### Complete bipartite graph

 $\begin{tabular}{ll} \bullet & M \ \ {\rm and} \ \ N \ \ {\rm nodes} \ \ {\rm in} \ \ {\rm partitions}; \\ M \times N \ \ {\rm edgescx} \end{tabular}$ 



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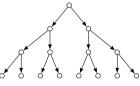
# Directed graph



# Rooted tree

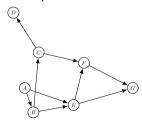


# Binary rooted tree



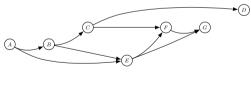
# Directed acyclic graph (DAG)

• No directed loops



### Topological order

• of the same DAG



# A few apparently innocuous problems related to graphs

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Graph representation Easy problem = a complete solution may be taught in bachelor courses.

Hard problem = a complete solution is unknown to this day.

(However, there often exist satisfactory approximate solutions.

Typically, they are quite advanced.)

### Clay Mathematics Institute

https://www.claymath.org/millennium-problems/

Offers prize \$1,000,000 for a complete solution of any of those hard questions.

The prize exists since the year 2000.

Nobody has claimed it yet :-(

# **Graph Connectivity**

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Graph representation Is there a path between any two nodes?

### Easy problem

Algorithm: DFS, BFS, Union-Find

**Complexity:** DFS, BFS O(|V| + |E|), Union-Find  $O(|E| \cdot \alpha(|V|))$ 

Yes, one connected component.

No,

four connected components.

# Connectivity

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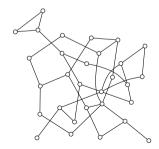
Graph representation

# Connectivity

Is there a path between any two nodes?

# Easy problem

# Is the graph connected?



# Connectivity

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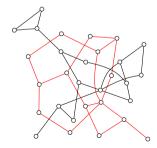
Graph representation

# Connectivity

Is there a path between any two nodes?

# Easy problem

Is the graph connected? (No!)



# Independence

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**Independence** – Maximum size of a set of nodes in which no two nodes are adjacent. **Hard problem in general** 



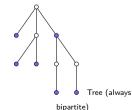




### Easy problem on graphs with some particular structure



Bipartite graph







Complete graph

# Independence

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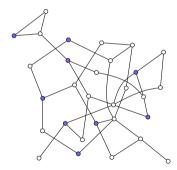
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Independence – Maximum size of a set of nodes in which no two nodes are adjacent.

Example: How many of them in this graph? more than 9?

# Hard problem



### Dominance

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**Dominance** – Maximum size of such set M of nodes that each node in the graph is either in M or is a neighbour of some node in M.

Example: A fire station must be located either in a village or in the immediately neighbouring village. How many fire stations are enough to serve the region?

### Hard problem







### Easy problem on graphs with some particular structure



Tree (apply Dynamic programming)





Complete graph

### Dominance

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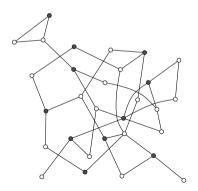
Graphs

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**Dominance** – Maximum size of such set M of nodes that each node in the graph is either in M or is a neighbour of some node in M.

Ex: A fire station must be located either in a village or in the immediately neighbour village. Can there be less than 11 fire stations to serve the region?

Hard problem



# Colorability, chromatic number

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Graphs

Graph representation

**Colorability, chromatic number** – Minimum number of colors needed to color each node so that no two neighbours have the same color.

Is 2 colors enough? - Easy problem. Graph must be bipartite.



2 colors, bipartite graph



2 colors for any tree



2 colors not enough in odd cycle



Graph not 2-colorable (odd cycle)

Is graph bipartite? Apply BFS. Mark by 1 all nodes in odd distance from the start and by 0 all nodes in even distance. If two nodes with the same mark are connected by an edge, the graph is not bipartite (two-colorable).

# Colorability, chromatic number

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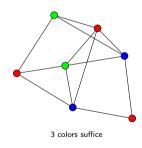
Graphs

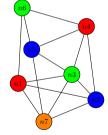
resentation:

### Colorability, chromatic number

Minimum number of colors needed to color each node so that any two neighbours have different color.

# Hard problem – Are 3 colors enough?









5 colors. Clique of size 5 (hard problem)

# Colorability, chromatic number

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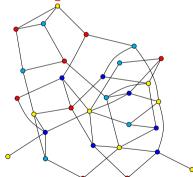
Graphs

resentation

# Colorability, chromatic number

Minimum number of colors needed to color each node so that any two neighbours have different color.

Hard problem – Are 3 colors enough?



4 colors are sufficient in

Maybe 3 colors would

this graph.

suffice too?

# Shortest and Longest Paths

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# Shortest paths

Minimum possible number of edges (nodes) on a path from A to B.

# Easy problem

Algorithms: BFS, Dijkstra, Bellman-Ford, Floyd-Warshall, Johnson...

**Complexities:** Polynomial, mostly less than  $O(|V|^3)$ .

# Longest paths

Typically, each node/edge can be visited at most once.

# Hard problem for general graphs

# Easy problem for trees and DAGs

Algorithm: Dynamic programming

Complexity: O(|V| + |E|)

# Minimum spanning tree

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### Minimum spanning tree

Minimum total cost (weight) of selected edges which connect all nodes in the graph. The selected edges form a tree.

# Easy problem

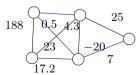
### Algorithms:

Prim's  $O(|V|^2)$  $O(|E| \cdot \log(|V|))$ 

with matrix representation with linked list representation and with binary heap

Kruskal's  $O(|E| \cdot \log(|V|))$ 

Borůvka's  $O(|E| \cdot \log(|V|))$ 



# Minimum spanning tree

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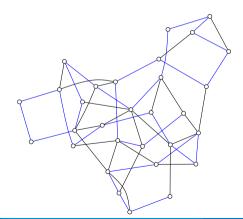
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# Minimum spanning tree

Minimum total cost (weight) of selected edges which connect all nodes in the graph. The selected edges form a tree.

# Easy problem



# Travelling salesman problem (TSP)

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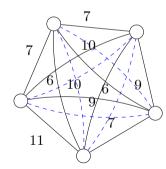
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### Travelling salesman problem (TSP)

Traverse a complete weighted graph, visit each node once and pay the minimum price for the journey = sum of costs of all visited edges.

### Hard problem



# Hamilton path / Hamilton cycle

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Graph representation

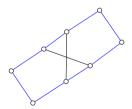
### Hamilton path

Is there a path in the graph which contains each node (exactly once)?

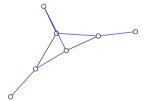
# Hamilton cycle

Is there a cycle in the graph which contains each node?

# Hard problem



Both Hamilton path and Hamilton cycle exist.



Only Hamilton path exists. There is no Hamilton cycle.



Neither a Hamilton path nor a Hamilton cycle exists.

### Euler trail

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**Euler trail** — A trail that visits every edge exactly once (vertices may be revisited).

Ex: Can a postman walk through each street in their region exactly once?

Easy problem Graph must be connected and it must contain at most two nodes of odd degree.

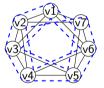
**Algorithm:** Hierholzer's O(|E|)











degree.

The trail is closed: all node degrees are even.

Euler trail does not exist, there are > 2 nodes with odd degree.

# Planar graph

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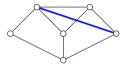
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Graphs

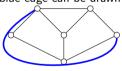
resentation

# Can the graph be drawn in a plane without crossing its edges?

Easy question (however, little bit more advanced)



The graph is planar, the blue edge can be drawn differently:





Not planar.
Non planar graphs "contain" either

- a complete graph on 5 nodes, or
- a complete bipartite graph on 3 and 3 nodes.

The planar graphs do not "contain" them.

### Algorithm:

Hopcroft and Tarjan, O(|V|)Boyer and Myrvold, O(|V|)

# Planar graph

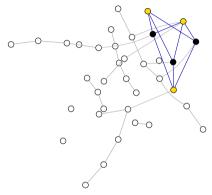
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Can the graph be drawn in a plane without crossing its edges?



### It is impossible here.

- Each black node is connected to each yellow node by a separate path.
- This part is a complete bipartite graph with partitions of size 3 and 3  $(K_{3,3})$ .
- Such a graph cannot be drawn in the plane without edge crossings.

# Clique number

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Graph representation

The size of the maximal clique, that is, of a subgraph which is complete — that is, of the subgraph where each node is connected to each other node.

Ex. Choose the largest group of your friends in which everybody knows each other.

### Hard problem



Clique number of all trees is 2. (Rather obviously.)

# Graph isomorphism

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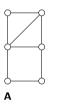
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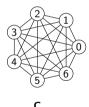
resentation

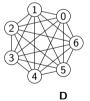
- Is the structure of two graphs identical?
- In other words, can one graph be drawn in such a way that it looks exactly as the other one?

### It is not known if this is a hard problem or an easy problem.









A and B are not isomorphic: the right central node in B has degree 5, but there is no such node in A, so their structures differ.

C and D are isomorphic: the nodes with the same labels correspond, and edges connect identical pairs.

# Partial recapitulation of the jungle of graph problems and their complexities

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# Easy problem

- Connectivity?
- Shortest path?
- Min. spanning tree?
- Euler trail?
- Planarity?

# "It depends..."

Colorability?

1,2 colors easy

3 or more colors hard

Isomorphism?

Trees, ciculants easy regular graphs, etc...

Longest path?

DAG, tree easy
General graph hard

# Hard problem

- Travelling salesman?
- Independence?
- Dominancy?
- Hamiltonicity?
- Clique number?

# Many more questions ... ?

Again, "it depends". There is no definite cookbook for determining the difficulty of a problem.

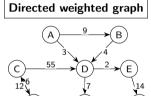
# Graph most usual representations — Directed weighted graph

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Graph

Graph representations



The representation is usually a more or less obvious combination of the methods in the previous cases — Weight matrix or linked list.

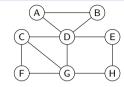
# Graph most usual representations — Undirected graph

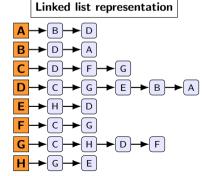
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Graph representations





		Adj						
	Α	В	С	D	Е	F	G	Н
Α	0	1	0	1	0	0	0	0
В	1	0	0	1	0	0	0	0
C	0	0	0	1	0	1	1	0
D	1	1	1	0	1	0	1	0
Ε	0	0	0	1	0	0	0	1
F	0	0	1	0	0	0	1	0
G	0	0	1	1	0	1	0	1
Н	0	0	0	0	1	0	1	0

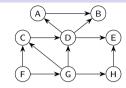
# Graph most usual representations — Directed graph

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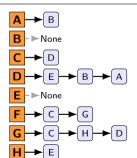
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Graph representations



# Linked list representation



# Adjacency matrix

	Α	В	c	D	Ε	F	G	н
Α	0 0 0 1 0 0 0	1	0	0	0	0	0	0
В	0	0	0	0	0	0	0	0
C	0	0	0	1	0	0	0	0
D	1	1	0	0	1	0	0	0
Ε	0	0	0	0	0	0	0	0
F	0	0	1	0	0	0	1	0
G	0	0	1	1	0	0	0	1
Н	0	0	0	0	1	0	0	0

# Graph most usual representations — Undirected weighted graph

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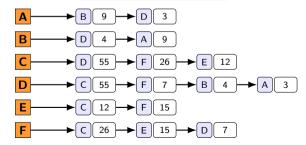
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Graph representations

# A 9 B C 55 D 7 E 15 F

### Linked list representation



# Weight (cost) matrix

	Α	В	0 0 0 55 12 26	D	E	F
Α	0	9	0	3	0	0
В	9	0	0	4	0	0
C	0	0	0	55	12	26
D	3	4	55	0	0	7
Ε	0	0	12	0	0	15
F	0	0	26	7	15	0

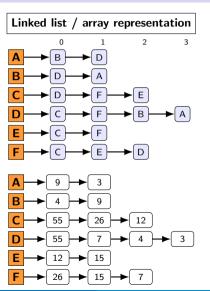
# Graph most usual representations — Undirected weighted graph — alternative rep.

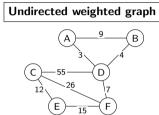
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The weights of edges are at the same index in the second list.

- + **Pro:** Simpler object or even no object at all in arrays.
- **Con:** Keeping lists in sync needs more care and caution in code.

# Graph most usual representations — Directed weighted graph

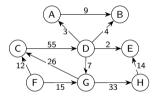
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# Directed weighted graph



The representation is usually a more or less obvious combination of the methods in the previous cases

— Weight matrix or linked list.