

**STATISTICAL MACHINE LEARNING (WS2021)**  
**SEMINAR 7**

**Assignment 1.** (*Gambler's ruin*) Consider a random walk on the set  $L = \{0, 1, 2, \dots, a\}$  starting in some point  $x \in L$ . The position jumps by either  $\pm 1$  in each time step (with equal probabilities). The walk ends if either of the boundary states  $0, a$  is hit. Compute the probability  $u(x)$  to finish in state  $a$  if the process starts in state  $x$ .

*Hints:*

- (1) What are the values of  $u(0)$  and of  $u(a)$ ?
- (2) Find a difference equation for  $u(x)$ ,  $0 < x < a$  by relating it with  $u(x - 1)$  and  $u(x + 1)$ .
- (3) Translate the difference equation into a relation between the successive differences  $u(x + 1) - u(x)$  and  $u(x) - u(x - 1)$ .
- (4) Deduce that the solution is a linear function of  $x$  and find its coefficients from the boundary conditions  $u(0)$  and  $u(a)$ .

**Assignment 2.** Let us consider a Markov chain model for sequences  $s = (s_1, \dots, s_n)$  of length  $n$  with states  $s_i \in K$  from a finite set  $K$ . Its joint probability distribution is given by

$$p(s) = p(s_1) \prod_{i=2}^n p(s_i | s_{i-1}).$$

The conditional probabilities  $p(s_i | s_{i-1})$  and the marginal probability  $p(s_1)$  for the first element are known. Let  $A \subset K$  be a subset of states and let  $\mathcal{A} = A^n$  denote the set of all sequences  $s$  with  $s_i \in A$  for all  $i = 1, \dots, n$ .

**a)** Find an efficient algorithm for computing the most probable sequence in  $\mathcal{A}$ .

**b)** Find an efficient algorithm for computing the probability  $\mathbb{P}(\mathcal{A}) = \sum_{s \in \mathcal{A}} p(s)$ .

**Assignment 3.** Let us consider the following matching problem. Given a sequence  $x = (x_1, \dots, x_m)$  of points  $x_i \in \mathbb{R}^2$  and another sequence  $y = (y_1, \dots, y_n)$  of points in the same space, we want to find an optimal matching between them. (Notice that the sequences may have different length).

A matching  $\tau$  is encoded as a path in the graph  $(V, E)$ , with nodes  $V = \{(i, j) \mid i = 1, \dots, m, j = 1, \dots, n\}$  and edges connecting each node  $(i, j)$  with nodes  $(i + 1, j)$ ,  $(i, j + 1)$  and  $(i + 1, j + 1)$ . The path should start in  $(1, 1)$  and end in  $(m, n)$ . The cost of the matching  $\tau$  is the sum of costs of the traversed nodes, where the cost for the node  $(i, j)$  is the Euclidean distance  $\|x_i - y_j\|$ .

Explain how to find the optimal matching for a pair of point sequences  $x$  and  $y$  by dynamic programming. What is the run time complexity of your algorithm?

**Assignment 4.** Consider a hidden Markov model

$$p(x, s) = p(s_1) \prod_{i=2}^n p(s_i | s_{i-1}) \prod_{i=1}^n p(x_i | s_i),$$

where  $x = (x_1, \dots, x_n)$  is a sequence of features and  $s = (s_1, \dots, s_n)$  is a sequence of hidden states, with values  $s_i$  from a finite set  $K$ . Given a sequence of features  $x$  we want to predict the sequence of hidden states that has generated  $x$ .

**a)** The predictor should minimise the expected loss

$$R(x, h) = \sum_{s \in K^n} p(x, s) \ell(s, h(x)),$$

where  $\ell(s, s')$  is the Hamming distance between sequences  $s$  and  $s'$ , i.e.

$$\ell(s, s') = \sum_{i=1}^n \mathbb{1}[s_i \neq s'_i].$$

Show that the optimal predictor for this loss is given by

$$s_i^* = \arg \max_{k \in K} p(s_i = k | x),$$

i.e. predicting the sequence of most probable states.

**b)** The predictor in a) requires to compute the marginal posterior probabilities  $p(s_i = k | x)$  for all positions  $i$  and all states  $k \in K$ . Show how to compute them for an HMM by performing dynamic matrix-vector multiplications from left to right and from right to left and combining the results.

*Hint:* Your algorithm will in fact compute the probabilities  $p(s_i = k, x)$ . The required normalisation for the posterior probabilities  $p(s_i = k | x)$  can be postponed and done in the last step.