B4M36DS2, BE4M36DS2: Database Systems 2

https://cw.fel.cvut.cz/wiki/courses/b4m36ds2/

Lecture 9

Wide Column Stores: Cassandra

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Lecture Outline

Wide column stores

Introduction

Apache Cassandra

- Sharding, replication
- CAP theorem
- Data model
- Cassandra query language
 - DDL statements
 - DML statements

Wide Column Stores

Data model

- Column family
 - Table is a collection of similar rows (not necessarily identical)
- Row
 - Row is a collection of columns
 - Should encompass a group of data that is accessed together
 - Associated with a unique row key
- Column
 - Column consists of a column name and column value (and possibly other metadata records)
 - Scalar values, but also flat sets, lists or maps may be allowed

Apache Cassandra



Apache Cassandra











- http://cassandra.apache.org/
- Features
 - Open-source, high availability, linear scalability, sharding (spanning multiple datacenters), peer-to-peer configurable replication, tunable consistency, MapReduce support
- Developed by Apache Software Foundation
 - Originally at Facebook
- Implemented in Java
- Operating systems: cross-platform
- Initial release in 2008

Key Features and Architecture

Decentralized Peer-to-Peer Architecture

All nodes are equal; no master-slave hierarchy.

High Availability and Fault Tolerance:

- Data automatically replicated across multiple nodes.
- Ensures uptime during node or hardware failures.

Linear Horizontal Scalability:

- Seamless addition of nodes increases capacity and performance.
- No downtime is required for scaling operations.

Gossip Protocol:

- Nodes communicate state information periodically.
- Maintains cluster synchronization without a central coordinator.

Ring Topology:

- Logical arrangement of nodes for efficient data distribution.
- Simplifies scaling and data partitioning.

Data Partitioning

COUNTRY		POPULATION
USA	New York	8.000.000
USA	Los Angeles	4.000.000
FR	Paris	2.230.000
DE	Berlin	3.350.000
UK	London	9.200.000
AU	Sydney	4.900.000
DE	Nuremberg	500.000
CA	Toronto	6.200.000
CA	Montreal	4.200.000
FR	Toulouse	1.100.000
JP	Tokyo	37.430.000
IN	Mumbai	20.200.000

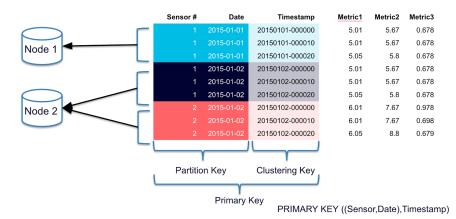
Partition Key

59 4.900.000 AU 12 CA 12 Montreal CA 45 Hashing Function Berlin DE Berlin 45 500,000 DE 500.000 Nuremberg USA New York USA



Source https://cassandra.apache.org/_/cassandra-basics.html

Partition and Clustering Key



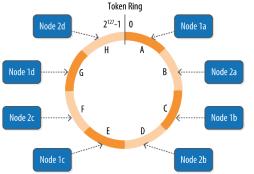
The Primary Key consists of two parts:

- Partition Key: (Sensor, Date) defines which node will store the data
- Clustering Key: Timestamp defines how data is sorted/organized within each partition

Source https://www.instaclustr.com/blog/cassandra-data-partitioning/

Data Partitioning

- Data Partitioning (Sharding):
 - Consistent Hashing: Distributes data based on partition key hashes.
 - Virtual Nodes (vnodes): Each physical node handles multiple token ranges (256 virtual nodes by default).
 - Partition Key Design: Ensures even data distribution and prevents hotspots.



Source: https://www.digihunch.com/2018/03/cassandra-architecture-summary/

Replication

Replication Strategies:

- Replication Factor (RF): Number of replicas for each piece of data.
- SimpleStrategy:

Suitable for single data center deployments. Replicates data to adjacent nodes in the ring.

NetworkTopologyStrategy:

Ideal for multi-data center clusters. Configures replication per data center.

Replica Placement:

- Distributes replicas across different racks and data centers.
- Enhances fault tolerance and data availability.

Impact on Performance and Availability:

- Higher RF increases data redundancy and fault tolerance.
- Balances storage costs with reliability and access speed.

Consistency, Synchronization, and CAP Theorem

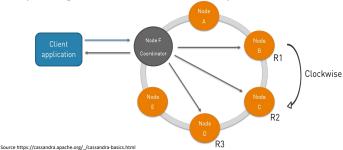
- Tunable Consistency Levels:
 - Levels Available: ANY, ONE, QUORUM, LOCAL_QUORUM, EACH QUORUM, ALL.
 - Per-Operation Configuration: Set consistency levels for individual reads and writes.
- Consistency and Synchronization Mechanisms:
 - Read Repair: Corrects inconsistencies during read operations.
 - Synchronous: Repairs during the read.
 - Asynchronous: Repairs in the background.
 - Hinted Handoff: Temporarily stores writes for unreachable nodes.
 - Replays hints when nodes recover.
- Anti-Entropy Repair:
 - **Full Repair:** Comprehensive data synchronization.
 - Incremental Repair: Targets recently changed data.

Consistency, Synchronization, and CAP Theorem

- Clock Synchronization and Conflict Resolution:
 - Uses timestamps (microsecond precision) to resolve write conflicts.
 - Last Write Wins: Most recent write overwrites previous ones.
- CAP Theorem Positioning:
 - AP System: Prioritizes Availability and Partition Tolerance.
 - Tunable Consistency: Allows configurations to approach CP as needed.
- Trade-offs and Implications:
 - Higher Consistency Levels:
 - Increased data accuracy.
 - Potentially higher latency and reduced availability.
 - Lower Consistency Levels:
 - Enhanced performance and availability.
 - Risk of reading stale or inconsistent data.

Cassandra Request Handling Flow

- Client sends a write request to Coordinator (Node F)
- Coordinator calculates token positions from its token map
- Based on token, Coordinator sends write request to all three replica nodes (R1-R3)
- Coordinator waits for the answer from one, two or three replicas depending on the Tunable Consistency level (ONE, QUORUM, ALL)
- Data is replicated clockwise: B(R1), C(R2), D(R3)
- For read request, Coordinator queries one, two or three replica nodes (R1-R3) depending on the Tunable Consistency level (ONE, QUORUM, ALL)



Database system structure

$$\mathsf{Instance} \to \mathsf{keyspaces} \to \mathsf{tables} \to \mathsf{rows} \to \mathsf{columns}$$

- Keyspace
- Table (column family)
 - Collection of (similar) rows
 - Rows do not need to have exactly the same columns
 - Table schema must be specified, yet can be modified later on
- Row
 - Collection of columns
 - Each row is uniquely identified by a compulsory primary key
- Column
 - Name-value pair + additional data

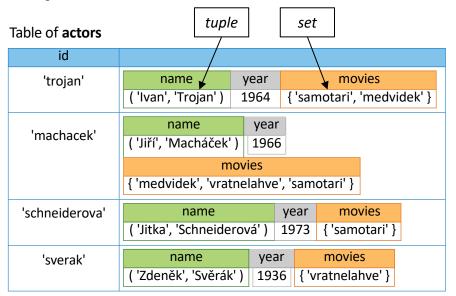
Column values

- Empty value
 - null
- Atomic values
 - Native data types such as text, integers, date, timestamp, ...
 - Tuples
 - Tuple of anonymous fields, each of <u>any</u> type (even different)
 - User-defined types (UDT)
 - Set of named fields of <u>any</u> type
- Collections
 - Lists, sets, and maps
 - Nested tuples, UDTs, or collections are also permitted, however, currently only in a frozen mode

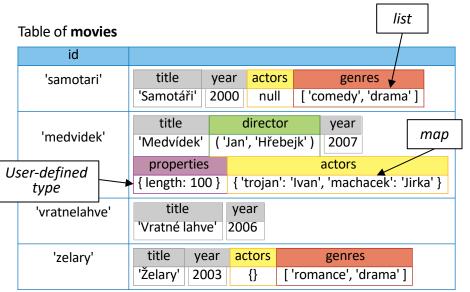
Collections

- List = ordered collection of values
 - This order is based on positions
 - Values do not need to be unique
- Set = collection of unique values
 - Values are internally <u>ordered</u> based on hash values
- Map = collection of key-value pairs
 - Keys must be unique
 - Pairs are internally <u>ordered</u> based on keys

Sample Data



Sample Data



Additional data associated with...

the whole column in case of atomic values, or each individual element of a collection

- Time-to-live (TTL)
 - After a certain period of time (number of seconds)
 a given column / element is automatically deleted
- Timestamp (writetime)
 - Timestamp of the last modification
 - Assigned automatically or manually as well
- Both the records can be queried
 - Limited support for collections and their elements, depending on the way of collection storage (frozen/nonfrozen)

Cassandra API

CQLSH

- Interactive command line shell
- bin/cqlsh (or via Docker container)
- Uses CQL (Cassandra Query Language)

Client drivers

- Officially supported by DataStax and community
- Available for various languages
 - Java, Python, Ruby, PHP, C++, Scala, Erlang, ...

Query Language

CQL = Cassandra Query Language

- Declarative query language
 - Inspired by SQL

DDL statements

- CREATE KEYSPACE creates a new keyspace
- CREATE TABLE creates a new table
- **...**

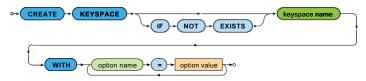
DML statements

- SELECT selects and projects rows from a single table
- INSERT inserts rows into a table
- UPDATE updates columns of rows in a table
- DELETE removes rows from a table
-

DDL Statements

Keyspaces

CREATE KEYSPACE



- Creates a new keyspace
- Replication option is mandatory
 - SimpleStrategy
 (single data centre, development/testing only)
 - NetworkTopologyStrategy (individual replication factor for each data center)

```
CREATE KEYSPACE moviedb
WITH replication = {'class': 'SimpleStrategy', 'replication_factor': 3}
```

Keyspaces

USE

Changes the current keyspace



DROP KEYSPACE

Removes a keyspace, all its tables, data etc.



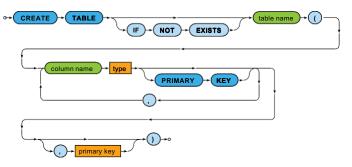
ALTER KEYSPACE

Modifies options of an existing keyspace

Tables

CREATE TABLE

- Creates a new table within the current keyspace
- Each table must have exactly one primary key specified



- None of the columns is compulsory (except the primary key)
- The primary key consists of partition key and (optionally) clustering key

Tables

Examples: tables for actors and movies

```
CREATE TABLE IF NOT EXISTS actors (
  id TEXT PRIMARY KEY.
  name TUPLE<TEXT, TEXT>,
  year SMALLINT,
  movies SET<TEXT>
CREATE TABLE IF NOT EXISTS movies (
  id TEXT.
  title TEXT,
  director TUPLE<TEXT, TEXT>,
 year SMALLINT,
  actors MAP<TEXT, TEXT>,
  genres LIST<TEXT>,
  countries SET<TEXT>,
  properties details,
  PRIMARY KEY (id)
```

Primary Keys

Primary keys have two parts

- Compulsory partition key
 - At least one column
 - Defines how individual rows are distributed between shards
- Optional clustering columns
 - Defines the order in which individual rows are locally stored by each shard

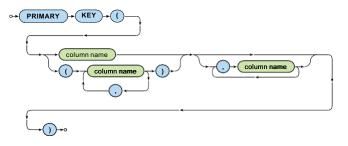
Primary Key Types:

- Simple primary key:
 - Can consist of a single column that becomes the partition key
 - May have no clustering columns
- Composite primary key:
 - Can have a composite partition key consisting of multiple columns
 - May include clustering columns

Primary Keys

Table-level primary key definition

- The first column / all columns in the embedded parentheses become the partition key
- All the remaining ones (if any) form the clustering columns

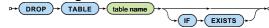


PRIMARY KEY(user_id, timestamp, action)
PRIMARY KEY((user_id, year), month, day)

Tables

DROP TABLE

Removes a table together with all data it contains



TRUNCATE TABLE

Preserves a table but removes all data it contains



ALTER TABLE

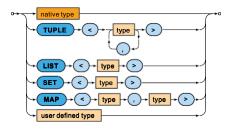
 Allows to alter, add or drop table columns, modify table properties and settings, manage secondary indexes, change caching option

It is recommended to add IF EXISTS to all these commands

Data Types

Types of columns

- Native types
- Tuples
- Collection types: lists, sets, and maps
- User-defined types



Native Data Types

Native types

- tinyint, smallint, int, bigint
 - Signed numbers (1B, 2B, 4B, 8B)
- varint
 - Arbitrary-precision integer
- decimal
 - Variable-precision decimal
- float, double
 - Floating point numbers (4B, 8B)
- boolean
 - Boolean values true and false

Native Data Types

Native types

- text (preferable), varchar
 - UTF8 encoded string (identical types)
 - Enclosed in single quotes (not double quotes)
 - Escaping sequence: ' '
- ascii
 - ASCII encoded string
- date, time, timestamp
 - Dates, times and timestamps
 - Supports timezone
 - E.g. '2016-12-05', '2016-12-05 09:15:00', '2016-12-05 09:15:00+0300'

Native Data Types

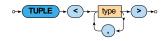
Native types

- counter 8B signed integer
 - Only 2 operations supported: incrementing and decrementing
 - I.e. value of a counter cannot be set to a particular number
 - Restrictions in usage
 - Counters cannot be a part of a primary key
 - Either all table columns (outside the primary key) are counters, or none of them
 - TTL is not supported
 - Counters do not support secondary indexes
 - ..
- blob arbitrary bytes
- inet IP address (both IPv4 and IPv6)
- ..

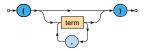
Tuple Data Types

Tuples

Declaration



Literals



• E.g. ('Jiří', 'Macháček')

When working with non-ASCII characters in tuples:

- Ensure cluster uses UTF-8 encoding
- Verify client app handles Unicode properly
- Maintain consistent encoding across all system components

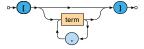
Collection Data Types

Lists

Declaration



Literals



• E.g. ['comedy', 'drama']

Note: Lists maintain order and allow duplicates. Consider performance impact for large lists.

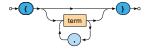
Collection Data Types

Sets

Declaration



Literals



E.g. { 'medvidek', 'vratnelahve', 'samotari' }

Note: Sets ensure unique values, unordered collection. Efficient for membership testing.

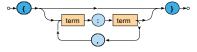
Collection Data Types

Maps

Declaration



Literals



• E.g. { 'machacek': 'Robert Landa' }

Note: Maps are key-value pairs. Keys must be unique. Efficient for key lookups.

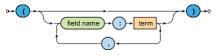
User-Defined Data Types

User-defined types (UDT)

Definition



- E.g. CREATE TYPE details (length SMALLINT, annotation TEXT)
- Literals



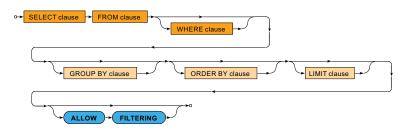
E.g. { length: 100 }

Note: UDTs allow creating complex, reusable data structures. Support nesting.

DML Statements

SELECT statement

Selects matching rows from a single table



Notes:

WHERE clause **must** include the complete partition key and can include only clustering columns (in the order they are defined). Other columns require ALLOW FILTERING (not recommended). ALLOW FILTERING may significantly impact performance.

Clauses of SELECT statements

- SELECT columns or values to appear in the result
- FROM single table to be queried
- WHERE filtering conditions to be applied on table rows
- GROUP BY columns to be used for grouping of rows
- ORDER BY criteria defining the order of rows in the result
- LIMIT number of rows to be included in the result

Examples

```
SELECT id, title, actors
FROM movies
WHERE id = 'medvidek'
```

```
SELECT id, title, actors
FROM movies
WHERE year = 2000 AND genres CONTAINS 'comedy'
```

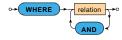
FROM clause

- Defines the table to be queried
 - From the current / selected keyspace
- Limited JOIN operations are supported:
 - Only within the same partition
 - Both tables must share the same partition key
 - Best practice is still data denormalization



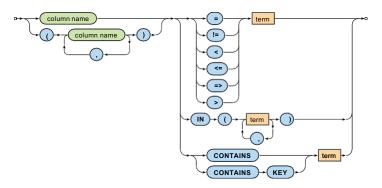
WHERE clause

 One or more relations a row must satisfy in order to be included in the query result



- · Conditions are limited:
 - Primary key columns can be queried directly
 - Non-key columns require secondary indexes
 - Partition key supports:
 - Equality (=)
 - IN operator
 - Clustering columns support:
 - Equality (=)
 - Inequality (<, >, <=, >=)
 - IN operator
 - ALLOW FILTERING enables non-indexed queries (not recommended)

WHERE clause: relations



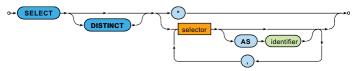
WHERE clause: relations

- Comparisons
 - =,!=,<,<=,=>,>
- IN
 - Returns true when the actual value is one of the enumerated
- CONTAINS
 - May only be used on collections (lists, sets, and maps)
 - Returns true when a collection contains a given element
- CONTAINS KEY
 - May only be used on maps
 - Returns true when a map contains a given key

Require secondary index for efficient queries

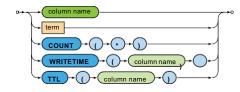
SELECT clause

- Defines columns or values to be included in the result
 - * = all the table columns
 - Aliases can be defined using AS
 - Supports aggregate functions (COUNT, MIN, MAX, AVG, SUM)
 - Collection functions supported (LIST, SET, MAP operations)
 - Supports user-defined functions (UDFs)



- DISTINCT duplicate rows are removed
 - DISTINCT operates only on partition key columns
 - Performance consideration: DISTINCT operations are memory-intensive

SELECT clause: **selectors**



COUNT(*)

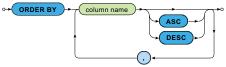
- Number of all the rows in a group (see aggregation)
- Functions SUM, AVG, MIN, MAX are also available

WRITETIME and TTL

- Selects modification timestamp / remaining time-to-live of a given column
- Cannot be used on collections and their elements
- Cannot be used in other clauses (e.g. WHERE)
- WRITETIME returns timestamp in microseconds since epoch
- Can only be used on one column per query

ORDER BY clause

- Defines the order of rows returned in the query result
- ORDER BY can only be used on clustering columns defined in table schema
- The order must match the clustering order defined in CREATE TABLE



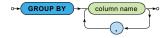
LIMIT clause

Limits the number of rows returned in the query result



GROUP BY clause

- Groups rows of a table according to certain columns
- Only groupings induced by primary key columns are allowed!



 When a non-grouping column would be accessed directly in the SELECT clause (i.e. without being wrapped by an aggregate function), the first value encountered will always be returned

GROUP BY clause: aggregates

- Native aggregates
 - COUNT(column)
 - Number of all the values in a given column
 - null values are ignored
 - MIN(column), MAX(column)
 - Minimal / maximal value in a given column
 - SUM(column)
 - Sum of all the values in a given column
 - AVG(column)
 - Average of all the values in a given column
- User-defined aggregates

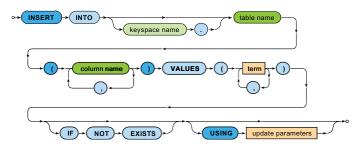
ALLOW FILTERING modifier

- By default, only non-filtering queries are allowed
 - I.e. queries where
 the number of rows read ~ the number of rows returned
 - Such queries have predictable performance
 - They will execute in a time that is proportional to the amount of data returned
- ALLOW FILTERING enables (some) filtering queries
 - May cause full table scan
 - Not recommended for production use
 - Better alternative: proper data modeling
- Best Practice
 - Use only on small tables or testing environments

Insertions

INSERT statement

- Inserts a new row into a given table
 - When a row with a given primary key already exists, it is updated
- Values of at least primary key columns must be set
- Names of columns must always be explicitly enumerated



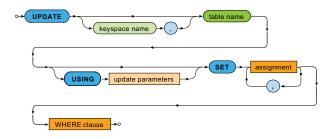
Insertions

Example

```
INSERT INTO movies (id, title, director, year, actors, genres) VALUES (
'stesti',
'Štěstí',
('Bohdan', 'Sláma'),
2005,
{ 'vilhelmova': 'Monika', 'liska': 'Toník' },
[ 'comedy', 'drama' ]
)
USING TTL 86400
```

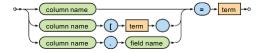
UPDATE statement

- Updates existing rows within a given table
 - When a row with a given primary key does not yet exist, it is inserted
- At least all primary key columns must be specified in the WHERE clause



UPDATE statement: **assignments**

- Describe modifications to be applied
- Allowed assignments:
 - Value of a whole column is replaced
 - Value of a list or map element is replaced
 - Items of lists are numbered starting with 0
 - Value of a user-defined type field is replaced



Examples

```
UPDATE
movies SET
year = 2006,
director = ('Jan', 'Svěrák'),
actors = { 'machacek': 'Robert Landa', 'sverak': 'Josef Tkaloun' },
genres = [ 'comedy' ],
countries = { 'CZ' }
WHERE id = 'vratnelahve'

UPDATE
movies SET
```

```
movies SET
actors['vilhelmova'] = 'Helenka',
genres[1] = 'comedy',
properties.length = 99
WHERE id = 'vratnelahve'
```

Note: Due to consistency issues, working with list indices (genres[1] = 'comedy') in Cassandra is not recommended. Use add/remove operations instead

Examples: modification of collection elements

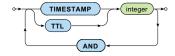
```
UPDATE movies
SET
  actors = actors + { 'vilhelmova': 'Helenka' },
  genres = [ 'drama' ] + genres,
  countries = countries + { 'SK' }
WHERE id = 'vratnelahve'
```

```
UPDATE movies
SET
  actors = actors - { 'vilhelmova', 'landovsky' },
  genres = genres - [ 'drama', 'sci-fi' ],
  countries = countries - { 'SK' }
WHERE id = 'vratnelahve'
```

Insertions and Updates

Update parameters

- TTL: time-to-live
 - 0, null or simply missing for persistent values
- TIMESTAMP: writetime

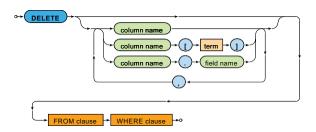


Only newly inserted / updated values are really affected

Deletions

DELETE statement

Removes the matching rows /
 Preserves these rows but removes the selected columns /
 Preserves these columns but removes elements of collections
 or fields of UDT values



Lecture Conclusion

Cassandra

Wide column store

Cassandra query language

- DDL statements
- DML statements
 - SELECT, INSERT, UPDATE, DELETE