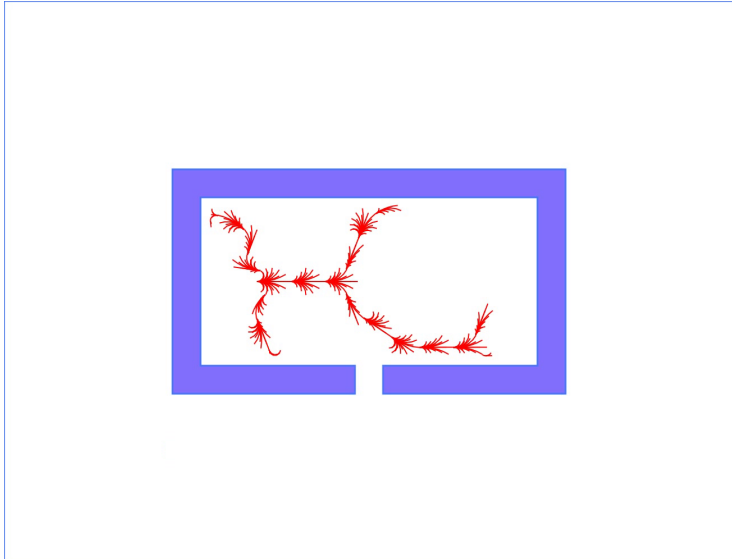


Motion planning: sampling-based planners I

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Planning for “car-like” vehicle



Motion/path planning

- Finding of collision-free trajectory/path for a robot
- Formulation using the configuration space \mathcal{C}
- \mathcal{C} is continuous \rightarrow conversion to a discrete representation (graph) \rightarrow graph search
- Combinatorial path planning
 - Require an explicit representation of \mathcal{C}_{obs}
 - For point/disc robots (if \mathcal{C} is same as \mathcal{W})
 - Visibility graphs, Voronoi diagrams, ...

Motion planning

Continuous problem

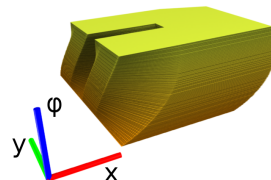
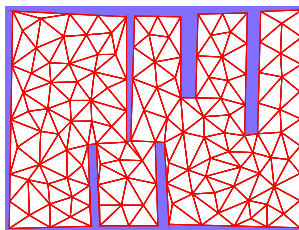
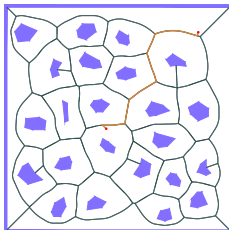
\mathcal{C} -space specification,
 $q_{\text{init}}, q_{\text{goal}}$

Discretization of \mathcal{C}

Explicit construction
random/deterministic
sampling of \mathcal{C}

Graph search

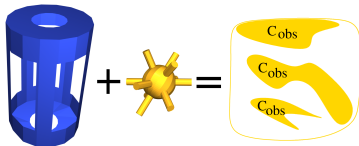
Dijkstra, A*, D*, ...



- Configuration space \mathcal{C} has as many dimensions as DOFs of the robot
- Obstacles \mathcal{C}_{obs} are given implicitly!

$$\mathcal{C}_{\text{obs}} = \{q \in \mathcal{C} \mid \mathcal{A}(q) \cap \mathcal{O} \neq \emptyset\}$$

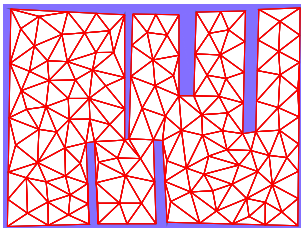
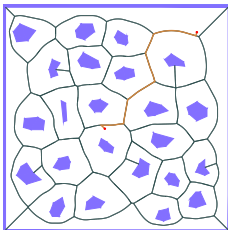
- \mathcal{C}_{obs} depends both on robot and obstacles!




- Generally, explicit geometry/shape of \mathcal{C}_{obs} is not available
- Problem of enumerating configurations in \mathcal{C}_{obs}
- Problem of enumerating “surface” configurations of \mathcal{C}_{obs}

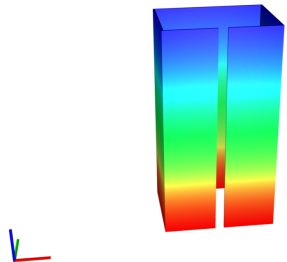
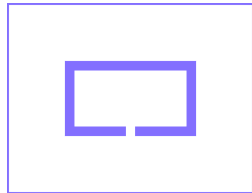
Problem of enumerating “surface” configurations of \mathcal{C}_{obs}

- We cannot generally/easy/fast say what are surface/boundary configurations of \mathcal{C}_{obs}
- This precludes combinatorial path planners (e.g., Visibility Graphs, Voronoi diagrams, Cell-decompositions, ...) to be used for high-dimensional \mathcal{C} -space
 - they require surface/boundary of \mathcal{C}_{obs}

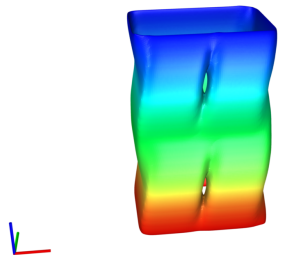


Configuration space: example I

- Map: 1000×700 units
- Robot: rectangle $20 \times a$ units
- $q = (x, y, \varphi)$
- \mathcal{C} visualized for $0 \leq \varphi < 2\pi$
- $\varphi = 0 \rightarrow$  $\leftarrow \varphi = 2\pi$




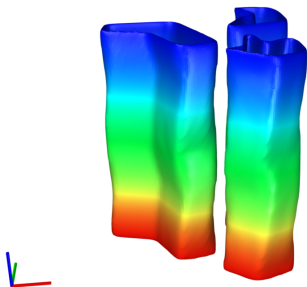
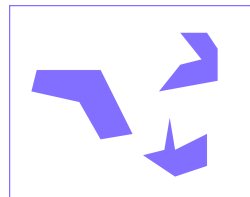
$a = 1$



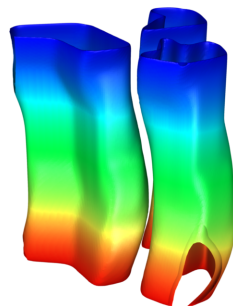
$a = 100$

Configuration space: example II

- Map: 2000×1600 units
- $q = (x, y, \varphi)$
- \mathcal{C} visualized for $0 \leq \varphi < 2\pi$
- $\varphi = 0 \rightarrow$  $\leftarrow \varphi = 2\pi$




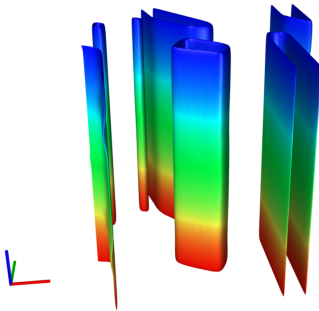
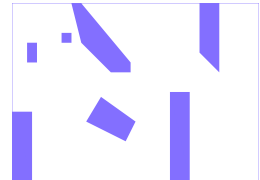
\mathcal{A} : rectangle 20×100 units



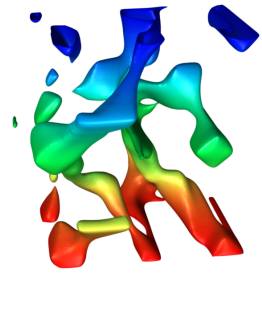
\mathcal{A} : equilateral triangle, side 100 units
(right-bottom "hole" caused by rendering clip)

Configuration space: example III

- Map: 5000×3000 units
- $q = (x, y, \varphi)$
- \mathcal{C} visualized for $0 \leq \varphi < 2\pi$
- $\varphi = 0 \rightarrow$  $\leftarrow \varphi = 2\pi$



\mathcal{A} : rectangle 20×100 units



\mathcal{A} : "u"-robot



- \mathcal{C} -space is usually high-dimensional in practical applications
 - Discretization not reasonable due to memory/time limits
- Non-trivial mapping between the shape of robot \mathcal{A} and obstacles \mathcal{O}
 - Simple obstacles in \mathcal{W} may be quite complex in \mathcal{C}
- Narrow passages (we will discuss later)

Early methods (combinatorial path planners)

- Designed for 2D/3D workspaces for point robots, complete, optimal (some), deterministic
- Limited only to special cases
- In late 1980s, these methods have become impractical

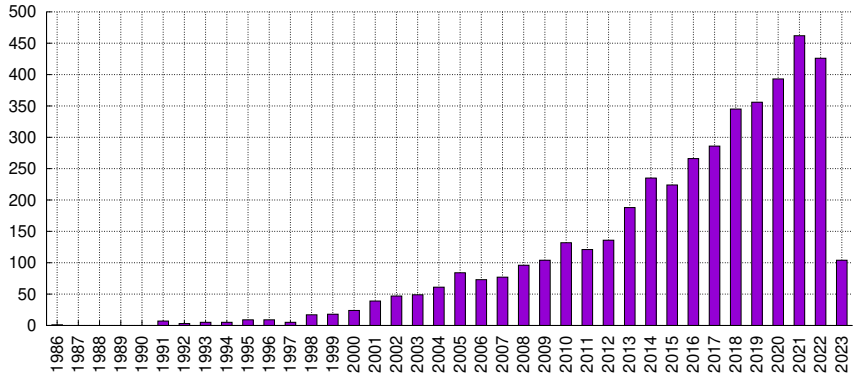
But general path/planning requires search in \mathcal{C} -space!

- If you are desperate, flip a coin \rightarrow randomization!

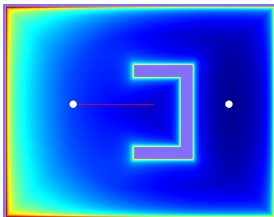
- Dijkstra's algorithm, 1959
- A*, 1968
- Configuration space, 1983
- Era of combinatorial planning, 1980s–1990s
- First planners using randomization, early 1990s
- Probabilistic roadmaps (PRM), 1995
- Rapidly-exploring Random Tree (RRT), 1998

- Dijkstra, E. W. "A Note on Two Problems in Connexion with Graphs", *Numerische Mathematik* 1, no. 1 (December 1959): 269–71.
- P. E. Hart, N. J. Nilsson and B. Raphael, "A Formal Basis for the Heuristic Determination of Minimum Cost Paths," in *IEEE Transactions on Systems Science and Cybernetics*, vol. 4, no. 2, pp. 100-107, July 1968,
- Lozano-Perez, "Spatial Planning: A Configuration Space Approach," in *IEEE Transactions on Computers*, vol. C-32, no. 2, pp. 108-120, Feb. 1983,

Publications with "sampling-based [path|motion] planning"



- Randomized path planner (RPP), 1991
 - Discrete workspace
 - Several potential fields for different control points of the robot
 - Gradient descent (GD) is performed for a selected point
 - If the goal is reached, the algorithm terminates
 - Otherwise, a different control point is selected and GD continues there
 - The escape from a local minimum is performed by a random walk



- ZZZ planner (1990)
 - Uses two planners: global and local
 - Global planner randomly places random goals in $\mathcal{C}_{\text{free}}$
 - Local planner uses potential field to connect these goals

• B. Glavina. Solving findpath by combination of goal-directed and randomized search. In IEEE International Conference on Robotics and Automation (ICRA), 1718-1723, 1990.

- Ariadne's clew algorithm (1998)
 - Two phase tree-based planner
 - Exploration phase: adds new configuration to tree rooted at q_{init}
 - Search phase: attempts to connect known (tree) configuration to q_{goal}
 - Both phases are solved using a genetic algorithm

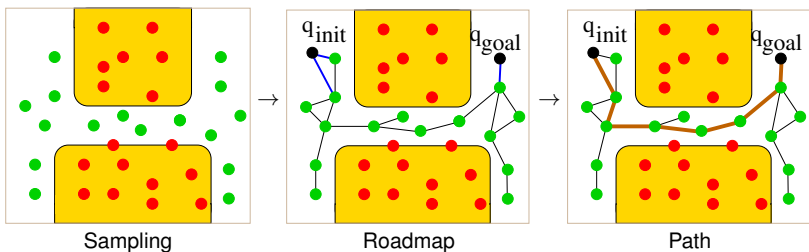
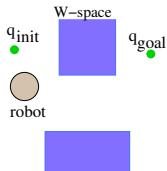
• E. Mazer and J. M. Ahuactzin and P. Bessiere; The Ariadne's Clew Algorithm, Journal of Artificial Intelligence Research, vol 9, 1998, 295-316

- Horsch planner (1994)
 - The first roadmap-based approach: generate random samples in $\mathcal{C}_{\text{free}}$
 - Connect samples by straight-line if possible
 - If the roadmap is disconnected, a random ray is shoot from one of its vertices
 - A contact configuration is added to the roadmap and connected with its nearest neighbors

• Horsch, T. and Schwarz, F. and Tolle, H.; Motion planning with many degrees of freedom-random reflections at C-space obstacles; IEEE International Conference on Robotics and Automation (ICRA), 1994

Main idea:

- \mathcal{C} is randomly sampled
- Each sample is a configuration $q \in \mathcal{C}$
- The samples are classified as free ($q \in \mathcal{C}_{\text{free}}$) or non-free ($q \in \mathcal{C}_{\text{obs}}$) using collision detection
- Free samples are stored and connected, if possible, by a “local planner”
- Result of sampling-based planning is a “roadmap” — graph
- The roadmap is the discretized image of $\mathcal{C}_{\text{free}}$
- Graph-search in the roadmap



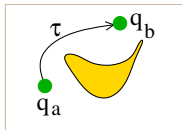
- Sampling-based planning can solve any problem formulated using \mathcal{C} -space
- ✓ Robots of arbitrary shapes
 - Robot shape is considered in collision detection
 - Collision detection is used as a “black-box”
 - Single-body or multi-body robots allowed
- ✓ Robots with many-DOFs
 - Because the search is realized directly in \mathcal{C} -space
 - Dimension of \mathcal{C} is determined by the DOFs
- ✓ Kinematic, dynamic and task constraints can be considered
 - It depends on the employed local planner

- Given configurations $q_a \in \mathcal{C}_{\text{free}}$ and $q_b \in \mathcal{C}_{\text{free}}$, the local planner attempts to find a path τ :

$$\tau : [0, 1] \rightarrow \mathcal{C}_{\text{free}}$$

such that $\tau(0) = q_a$ and $\tau(1) = q_b$

τ must be collision free!



Control-theory approach: special cases

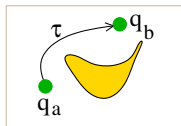
- We can assume that q_a and q_b are “near” without obstacles
- Two-point boundary value problem (BVP)
- Local planner is designed as a controller
- But problems are with obstacles!

Generally:

- The definition of “local planning” is same as motion planning
- same complexity as motion planning!

Exact local planners

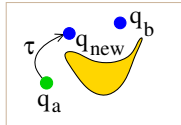
- For certain systems, BVP can be solved analytically
- Example: car-like without backward motions \rightarrow Dubins car



Exact local planner

Approximate local planners

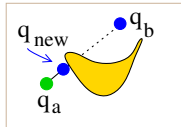
- Path τ connects q_a with q_{new} that is near-enough from q_b
- Computation e.g. using forward motion model and integration over time Δt



Approximate

Straight-line local planners

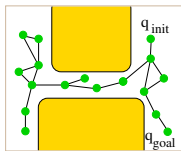
- Connects q_a and q_b by line-segment
- Check the collisions of the line-segment
- Connect q_a with the first contact configuration q_{new} or with q_b if no collision occurs
- Suitable for systems without kinematic/dynamic constraints



Straight-line

Multi-query methods

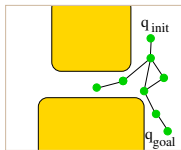
- Can find paths between multi start/goal queries
- Requires to build a roadmap covering whole $\mathcal{C}_{\text{free}}$
- Probabilistic Roadmaps (PRM) + many derivatives
- ✓ good for frequent planning and replanning
- ✗ sometimes slower construction



Multi-query roadmap

Single-query methods

- The roadmap is built only to answer a single start/goal query
- The sampling of \mathcal{C} terminates if the query can be answered
- Tree-based planners: Rapidly-exploring Random Trees (RRT), Expansive-space Tree (EST) + their variants
- ✓ Practically faster for single-query
- ✗ Any subsequent planning requires novel search of \mathcal{C}
- ✗ Slow for multi-query planning



Single-query roadmap

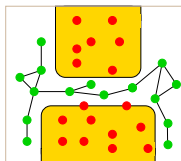
- Two-phase method: learning phase and query phase

Learning phase

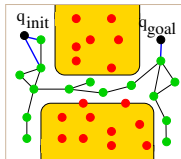
- Random samples are generated in \mathcal{C}
- Samples are classified as free/non-free; free samples are stored
- Each sample is connected to its near neighbors by a local planner
- Final roadmap may contain cycles

Query phase:

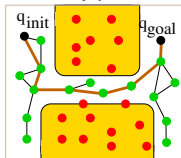
- Answers path/motion planning from $q_{init} \in \mathcal{C}_{free}$ to $q_{goal} \in \mathcal{C}_{free}$
- q_{init} and q_{goal} are connected to their nearest neighbors in the roadmap (using local planner)
- Graph-search of the roadmap



Learning phase



Query phase



Path

- L. E. Kavraki, P. Svestka, et al., "Probabilistic roadmaps for path planning in high-dimensional configuration spaces, ". IEEE Trans. on Robotics and Automation, 12(4), 1996.

- Simultaneous sampling + roadmap expansion
- q_{rand} is connected to each graph component only once
- Roadmap is a tree structure

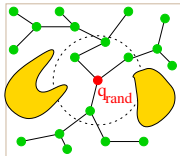
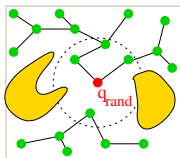
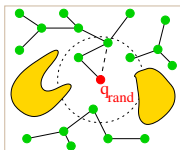
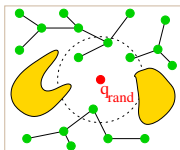
```

1   $V = \emptyset; E = \emptyset$  // vertices and edges
2   $G = (V, E)$  // empty roadmap
3  while  $|V| < n$  do
4       $q_{\text{rand}}$  = generate random sample in  $\mathcal{C}$ 
5      if  $q_{\text{rand}}$  is collision-free then
6           $G.addVertex(q_{\text{rand}})$ 
7          foreach  $q \in V.neighborhood^*(q_{\text{rand}})$  do
8              if not  $G.sameComponent(q_{\text{rand}}, q) \wedge connect(q_{\text{rand}}, q)$ 
9                  then
4               $G.addEdge(q_{\text{rand}}, q)$ 

```

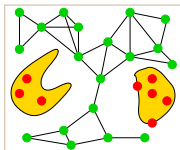
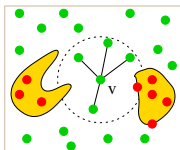
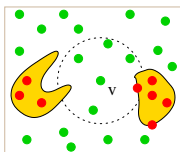
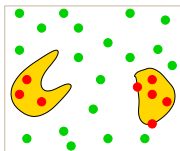
- $neighborhood^*$ returns q by increasing distance from q_{rand}

• L. E. Kavraki, P. Svestka, et al., "Probabilistic roadmaps for path planning in high-dimensional configuration spaces,". IEEE Trans. on Robotics and Automation, 12(4), 1996.



- Separate sampling and roadmap connection
- Each node is connected to its nearest neighbors
- Roadmap can contains cycles

```
1  $V = \emptyset; E = \emptyset$  // vertices and edges
2 while  $|V| < n$  do // generating n collision-free
  samples
3    $q_{\text{rand}} = \text{generate random sample in } \mathcal{C}$ 
4   if  $q_{\text{rand}}$  is collision-free then
5      $V = V \cup \{q_{\text{rand}}\}$ 
6 foreach  $v \in V$  do // connecting samples to roadmap
7    $V_n = V.\text{neighborhood}(v)$ 
8   foreach  $u \in V_n, u \neq v$  do
9     if connect( $u, v$ ) then // local planner
10     $E = E \cup \{(u, v)\}$ 
11  $G = (V, E)$  // final roadmap
```



• S. Karaman, and E. Frazzoli. "Sampling-based algorithms for optimal motion planning." The international journal of robotics research 30.7 (2011): 846-894.

- Behavior of sPRM is mostly influenced by $V.neighborhood$ function
- Several variants were proposed and analyzed

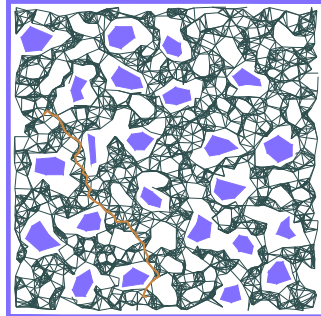
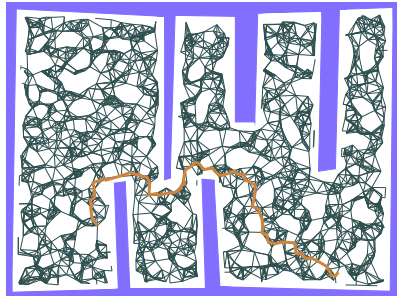
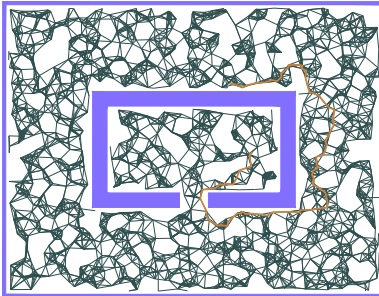
k -nearest sPRM (aka k -sPRM)

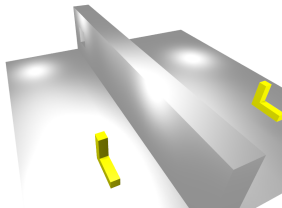
- $V.neighborhood$ provides k nearest neighbors from q_{rand}
- Probabilistically complete if $k \neq 1$
- Is not asymptotically optimal
- Usually $k = 15$

Variable radius sPRM

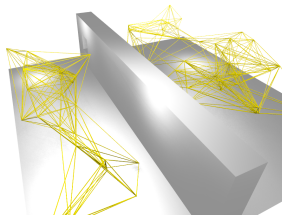
- $V.neighborhood$ returns nearest neighbors of q_{rand} within a radius r
- The choice of r influences completeness and optimality of sPRM
- Most important — PRM* planner

sPRM example 2D \mathcal{W}

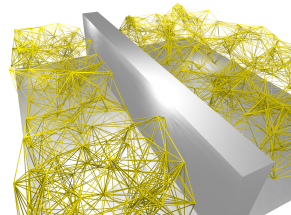




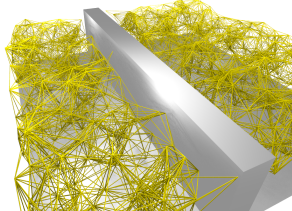
start/goal



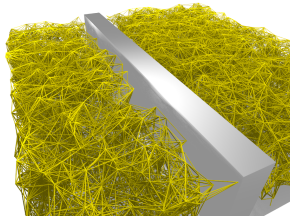
$n = 100$



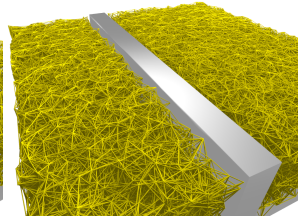
$n = 700$



$n = 1500$



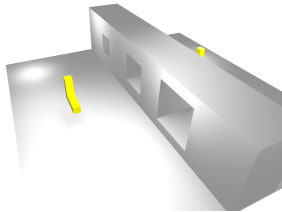
$n = 8000$



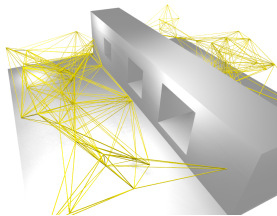
$n = 50000$

The wall contains one window, but no path found with 50k samples

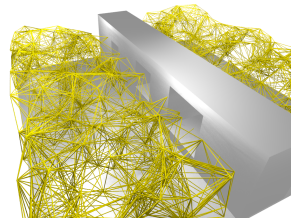
sPRM example 3D \mathcal{W}



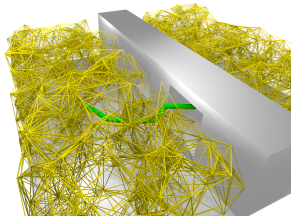
start/goal



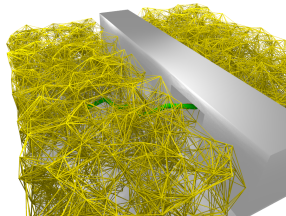
$n = 100$



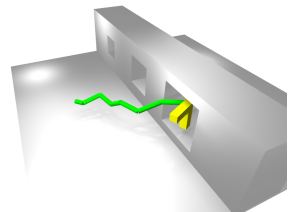
$n = 1000$



$n = 2100$



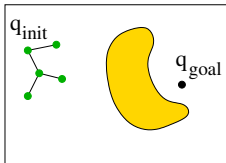
$n = 4100$



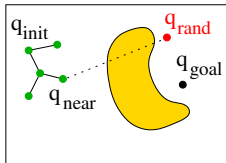
solution

- Incremental search of \mathcal{C}
- Collision-free configurations are stored in tree \mathcal{T}
- \mathcal{T} is rooted at q_{init}
- Tree is expanded towards random samples q_{rand}
- The search terminates if tree is close enough to q_{goal} , or after l_{max} iterations

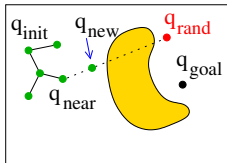
```
1 initialize tree  $\mathcal{T}$  with  $q_{init}$ 
2 for  $i = 1, \dots, l_{max}$  do
3    $q_{rand} =$  generate randomly in  $\mathcal{C}$ 
4    $q_{near} =$  find nearest node in  $\mathcal{T}$  towards
      $q_{rand}$ 
5    $q_{new} =$  localPlanner from  $q_{near}$  towards
      $q_{rand}$ 
6   if canConnect( $q_{near}, q_{new}$ ) then
7      $\mathcal{T}.addNode(q_{new})$ 
8      $\mathcal{T}.addEdge(q_{near}, q_{new})$ 
9     if  $\rho(q_{new}, q_{goal}) < d_{goal}$  then
10      return path from  $q_{init}$  to  $q_{goal}$ 
```



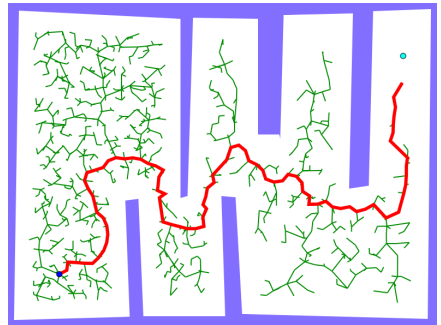
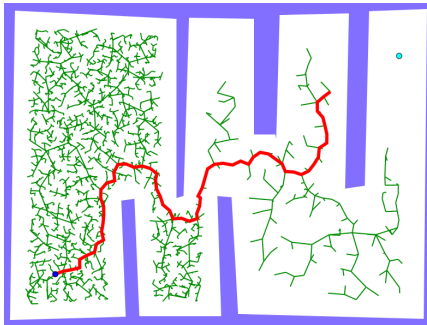
Tree



Sampling

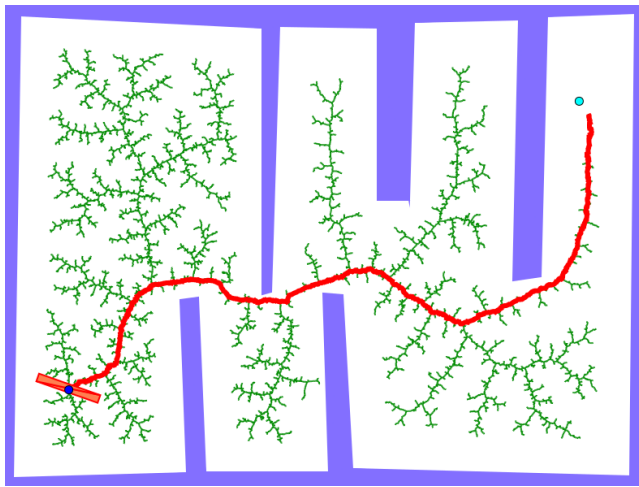


Tree extension

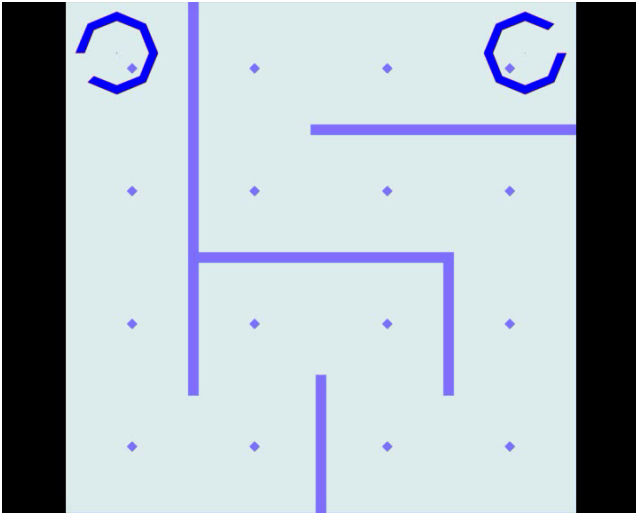


- 2D robot, rotation allowed \rightarrow 3D \mathcal{C}
- Why the tree does not “touch” the obstacles?

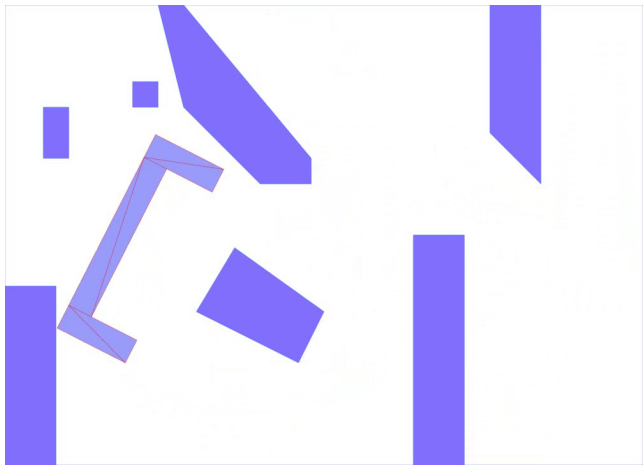
RRT example in 2D \mathcal{W}

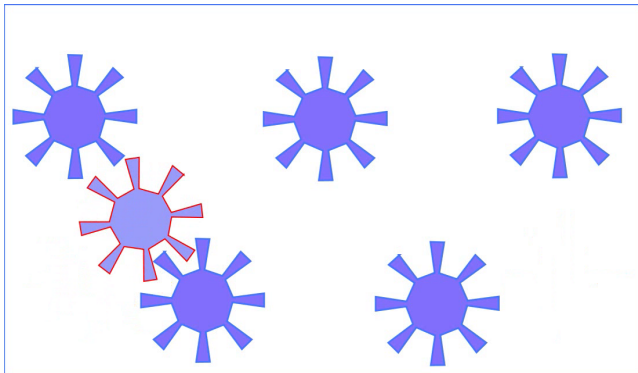


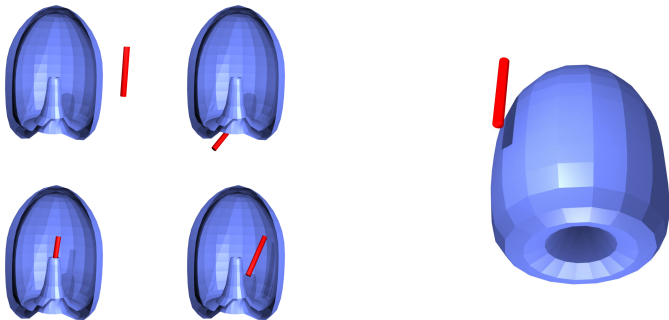
RRT example in 2D \mathcal{W}



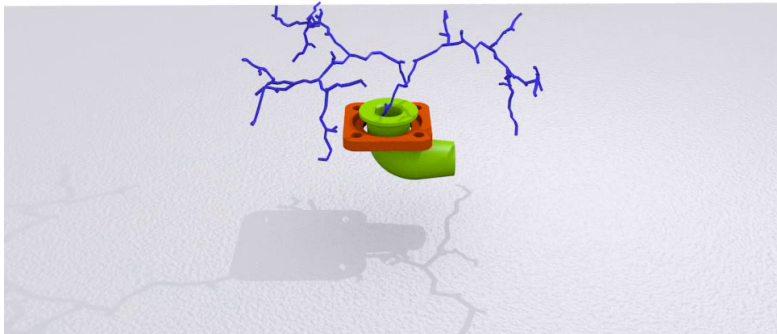
RRT example in 2D \mathcal{W}







- 3D Bugtrap benchmark
parasol.tamu.edu/groups/amatogroup/benchmarks/
- 3D robot in 3D space \rightarrow 6D \mathcal{C}

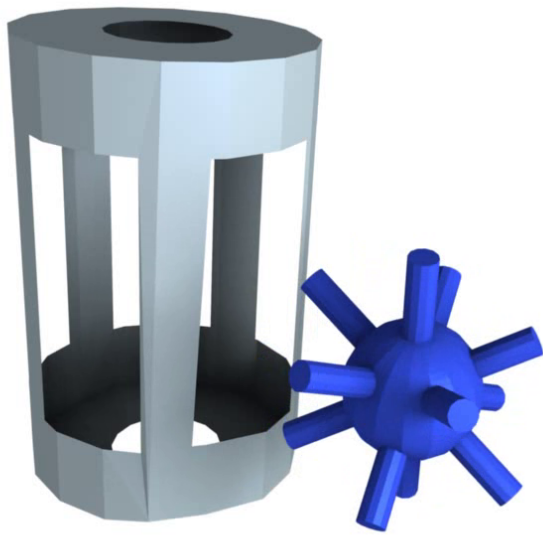


- 3D Flange benchmark

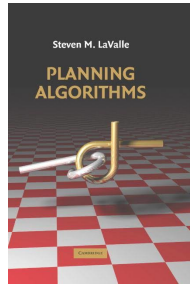
parasol.tamu.edu/groups/amatogroup/benchmarks/

- 3D robot in 3D space \rightarrow 6D \mathcal{C}

RRT example in 3D \mathcal{W}



RRT example in 3D \mathcal{W}



Straight-line expansion: make the line-segment S from q_{near} to q_{rand}

Variants:

A If S is collision-free, expand the tree only by

$$q_{\text{new}} = q_{\text{rand}}$$

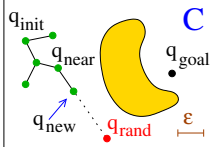
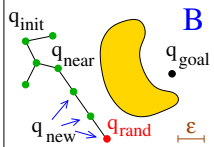
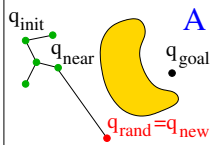
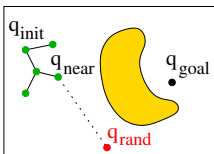
- Creates long segments, fast exploration of \mathcal{C}
- Requires nearest-neighbor search to consider point-segment distance
- Requires connection in the middle of line-segment

B If S is collision-free, discretize S and expand the tree by all points on S

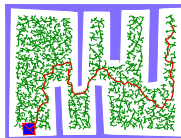
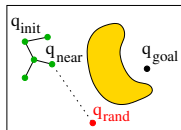
- Most used, enables fast nearest-neighbor search

C Find configuration $q_{\text{new}} \in S$ at the distance ε from q_{near} . Expand tree by q_{new} if it's collision-free

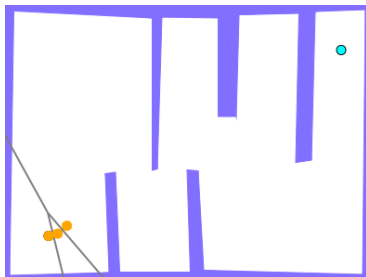
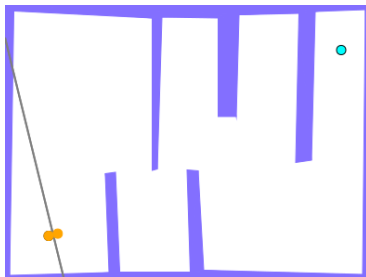
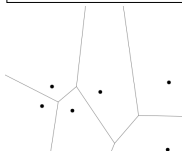
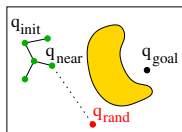
- Basic RRT, slower growth than **B**
- Enables fast nearest-neighbor search



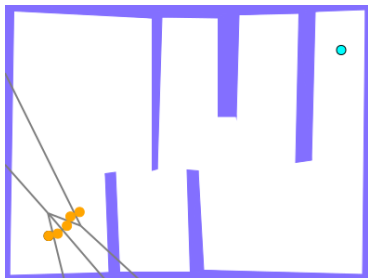
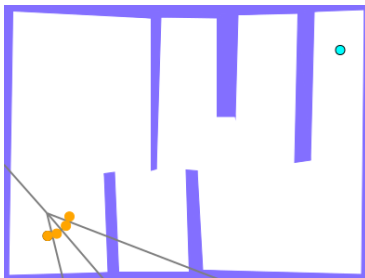
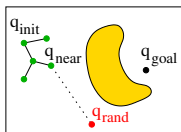
- RRT builds a tree \mathcal{T} of collision-free configurations
 - \mathcal{T} is rooted at q_{init}
 - \mathcal{T} is without cycles
 - Path from q_{init} to q_{goal} :
 - Find nearest node $q'_{goal} \in \mathcal{T}$ towards q_{goal}
 - Start at q'_{goal} and follow predecessors to q_{init}
 - Existing \mathcal{T} can answer queries starting at q_{init}
 - if goal is not in/near current \mathcal{T} , \mathcal{T} is further grown
 - Non-optimal
 - Probabilistically complete
-
- Why the tree does not grow to itself?
 - Why does it “rapidly” explore the \mathcal{C} -space?
... because of Voronoi bias!



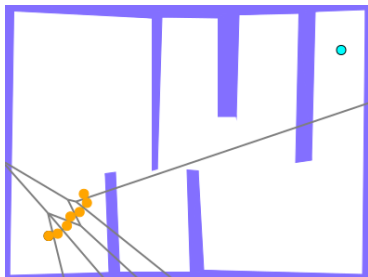
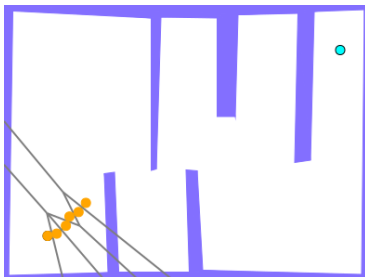
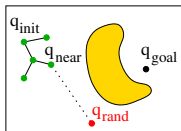
- RRT prefers to expand \mathcal{T} towards unexplored areas of \mathcal{C}
- This is caused by **Voronoi bias**:
 - q_{rand} is generated **uniformly** in \mathcal{C}
 - \mathcal{T} is expanded from **nearest** node in \mathcal{T} **towards** q_{rand}
 - The probability that a node $q \in \mathcal{T}$ is selected for the expansion is proportional to the area/volume of it's Voronoi cell
- Voronoi bias is implicit (caused by the nearest-rule selection)



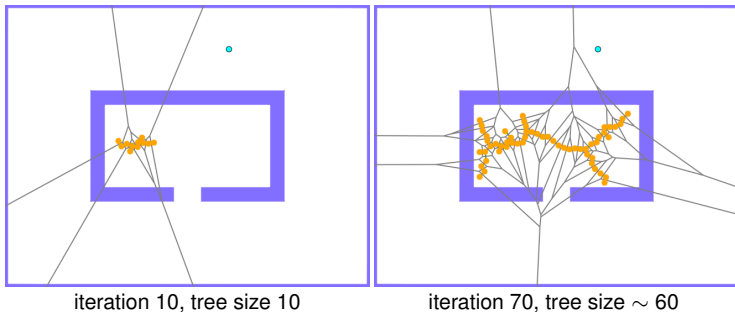
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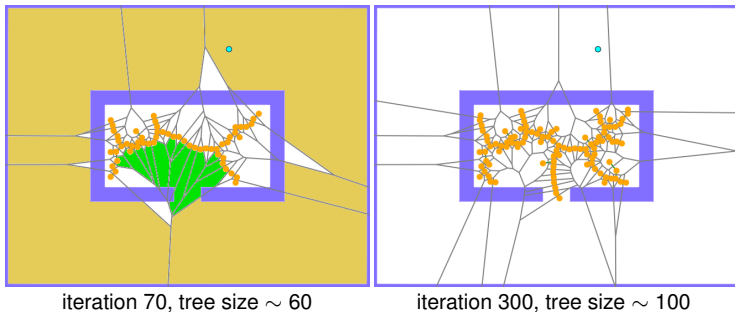


- Nearest-neighbors/Voronoi bias do not respect obstacles!
- If a node having large Voronoi cells is near an obstacle \rightarrow tree expansion is blocked at this node



- Tree grows well until iteration 70
- Yellow: areas with high prob. of being selected for expansion
- Green: areas that show be selected for expansion so the tree can escape the obstacle
- The tree does not expand much until iteration 300!

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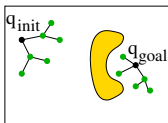


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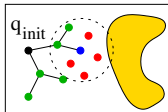
- Builds two trees \mathcal{T}_i and \mathcal{T}_g (from q_{init} and q_{goal})
- Weight $w(q)$ is computed for each configuration q
- Nodes are selected for expansion with probability $w(q)^{-1}$
- Expansion of one tree \mathcal{T} :

```
1  $q' =$  select node from  $\mathcal{T}$  with probability  $w(q)^{-1}$ 
2  $Q = k$  random points around
    $q' : Q = \{q \in \mathcal{C}_{free} \mid \varrho(q, q') < d\}$ 
3 foreach  $q \in Q$  do
4    $w(q) =$  compute weight of the sample  $q$ 
5   if  $rand() < w(q)^{-1}$  and  $connect(q, q')$  then
6      $\mathcal{T}.addNode(q)$ 
7      $\mathcal{T}.addEdge(q', q)$ 
```

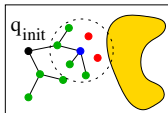
- $w(q)$ is the number of nodes in \mathcal{T} around q
- Both \mathcal{T}_i and \mathcal{T}_g grow until they approach each other
- Trees are connected using local planner between their nearest nodes



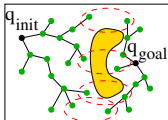
\mathcal{T}_i and \mathcal{T}_g



q' , samples Q

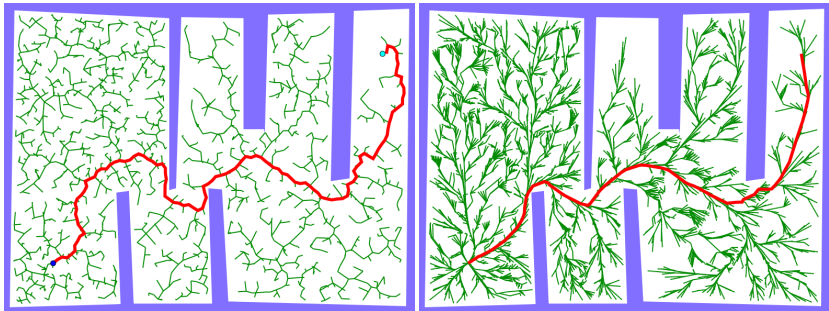


connected, ignored



pairs for tree connection

- PRM/RRT/EST do not consider any optimality criteria
- Only sPRM is asymptotically optimal
- PRM* and RRT* are new planners for which asymptotic optimality was proven



RRT

RRT*

• S. Karaman, and E. Frazzoli. "Sampling-based algorithms for optimal motion planning." The international journal of robotics research 30.7 (2011): 846-894.

- PRM* is an improved version of sPRM
- PRM* uses “optimal” radius $r(n)$ for searching the nearest neighbors depending on the actual number of nodes n :

$$r(n) = \gamma_{PRM} \left(\frac{\log(n)}{n} \right)^{\frac{1}{d}}$$
$$\gamma_{PRM} > \gamma_{PRM}^* = 2 \left(1 + \frac{1}{d} \right)^{\frac{1}{d}} \left(\frac{\mu(\mathcal{C}_{\text{free}})}{\zeta_d} \right)^{\frac{1}{d}}$$

- d is the dimension of \mathcal{C}
- $\mu(\mathcal{C}_{\text{free}})$ is the volume of $\mathcal{C}_{\text{free}}$
- ζ_d is the volume of the unit ball in the d -dimensional Euclidean space
- r decays with n
- r depends also on the problem instance! — why?

PRM* algorithm

- Same as for sPRM, just the line 7 is changed to:
 $V_n = V.\text{neighborhood}(v, r(n))$, where $n = |V|$

- Variant of PRM* that uses k -nearest neighbors definitions

$$k = k_{PRM} \log(n)$$

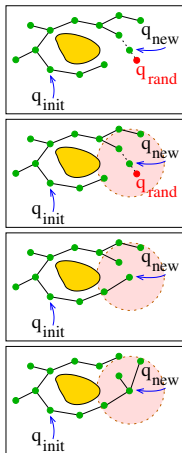
$$k_{PRM} > k_{PRM}^* = e \left(1 + \frac{1}{d} \right)$$

- The constant k_{PRM}^* depends only on d and not on the problem instance (compare it to γ_{PRM}^*)
- $k_{PRM} = 2e$ is a valid choice for all problem instances

k -nearest PRM* algorithm (aka k -PRM*)

- Same as for sPRM, just the line 7 is changed to:
 $V_n = k$ -nearest neighbors from V , $k = k_{PRM} \log(n)$

- Optimal version of RRT
- For each node, a cost of the path from q_{init} to that node is established
- RRT* has improved tree expansion and nearest-neighbor search
- Tree expansion by node q_{new}
 - Parent of q_{new} is optimized to minimize cost at q_{new}
 - After q_{new} is connected to tree, nodes in its vicinity are “rewired” via q_{new} if it improves their cost
- Nearest-neighbor search
 - Number of nearest-neighbors varies similarly to PRM*

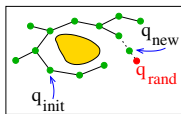


• S. Karaman, and E. Frazzoli. "Sampling-based algorithms for optimal motion planning." The international journal of robotics research 30.7 (2011): 846-894.

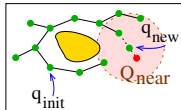
```

1 initialize tree  $\mathcal{T}$  with  $q_{init}$ 
2 for  $i = 1, \dots, l_{max}$  do
3    $q_{rand}$  = generate randomly in  $\mathcal{C}$ 
4    $q_{near}$  = find nearest node in  $\mathcal{T}$  towards  $q_{rand}$ 
5    $q_{new}$  = localPlanner from  $q_{near}$  towards  $q_{rand}$ 
6   if  $q_{new}$  is collision-free then
7      $Q_{near} = \mathcal{T}.neighborhood(q_{new}, r)$ 
8      $\mathcal{T}.addNode(q_{new})$  // new node to tree
9      $q_{best} = q_{near}$  // best parent of  $q_{new}$  so far
10     $c_{best} = cost(q_{near}) + cost(line(q_{near}, q_{new}))$ 
11    foreach  $q \in Q_{near}$  do
12       $c = cost(q) + cost(line(q, q_{new}))$ 
13      if  $canConnect(q, q_{new})$  and  $c < c_{best}$  then
14         $q_{best} = q$  // new parent of  $q_{new}$  is  $q$ 
15         $c_{best} = c$  // its cost
16     $\mathcal{T}.addEdge(q_{best}, q_{new})$  // tree connected to
17     $q_{new}$ 
18    foreach  $q \in Q_{near}$  do // rewiring
19       $c = cost(q_{new}) + cost(line(q_{new}, q))$ 
20      if  $canConnect(q_{new}, q)$  and  $c < cost(q)$  then
        change parent of  $q$  to  $q_{new}$ 

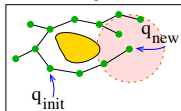
```



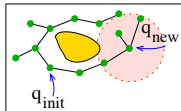
lines 3–5



line 7



lines 10–16



lines 17–20

- $cost(line(q_1, q_2))$ is cost of path from q_1 to q_2 (path by the local planner)
- $cost(q), q \in \mathcal{T}$ is cost of the path from q_{init} to q (path in \mathcal{T})
- nearest neighbors Q_{near} are searched within radius r depending on the number of nodes n in the tree:

$$r = \min \left\{ \gamma_{RRT}^* \left(\frac{\log(n)}{n} \right)^{\frac{1}{d}}, \eta \right\}$$

$$\gamma_{RRT}^* = 2 \left(1 + \frac{1}{d} \right)^{\frac{1}{d}} \left(\frac{\mu(\mathcal{C}_{free})}{\zeta_d} \right)^{\frac{1}{d}}$$

- d is the dimension of \mathcal{C}
- $\mu(\mathcal{C}_{free})$ is the volume of \mathcal{C}_{free}
- ζ_d is the volume unit ball in the d -dimensional Euclidean space
- η is constant given by the used local planner
- r decays with n
- r depends also on the problem instance

Alternative k -nearest RRT* (aka k -RRT*)

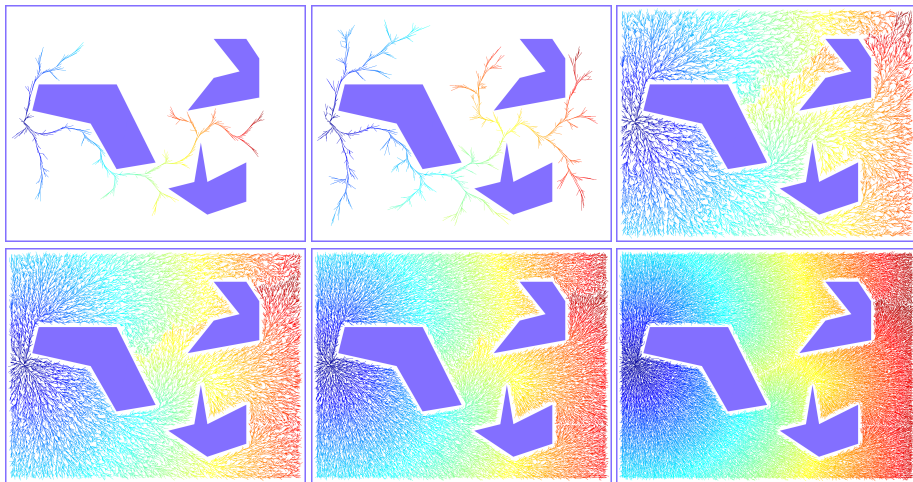
- k -nearest neighbors are selected for parent search and rewiring

$$k = k_{RRT} \log(n)$$

$$k_{RRT} > k_{RRT}^* = e \left(1 + \frac{1}{d} \right)$$

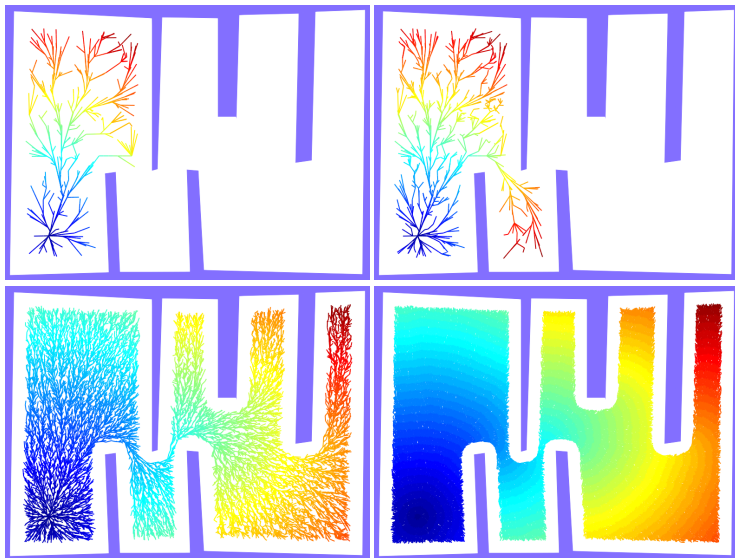
- n is the number of nodes in \mathcal{T}
- k -RRT* has same implementation as RRT* just line 7 is changed to $Q_{near} = \text{find } k \text{ nearest neighbors in } \mathcal{T} \text{ towards } q_{new}$

RRT*: example in 2D \mathcal{W}



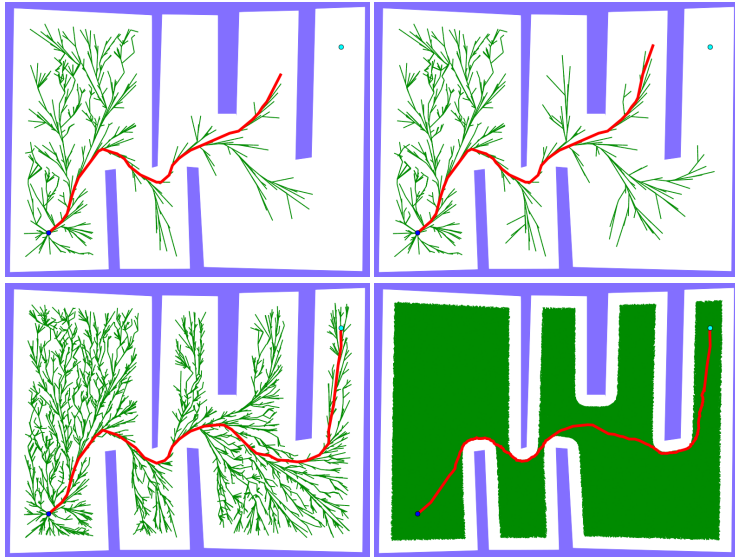
Rectangle robot, rotation allowed \rightarrow 3D \mathcal{C}

RRT*: example in 2D \mathcal{W}



2D rectangle robot \rightarrow 3D \mathcal{C} . The colormap shows the path length from q_{init} .
But is it really good?

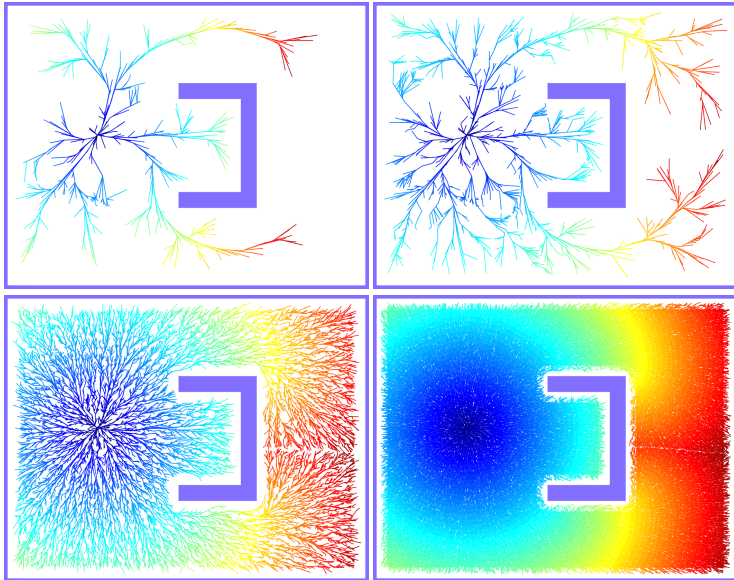
RRT*: example in 2D \mathcal{W}



2D rectangle robot \rightarrow 3D \mathcal{C}

Depicted path demonstrates the slow convergence of the path quality

RRT*: example in 2D \mathcal{W}



Algorithm	Probabilistic completeness	Asymptotic optimality
RRT	Yes	No
PRM	Yes	No
sPRM	Yes	Yes
k -sPRM	No if $k = 1$	No
PRM* / k -PRM*	Yes	Yes
RRT* / k -RRT*	Yes	Yes

- If you don't need optimal solution, stay with RRT/PRM
 - RRT is faster than RRT*
 - RRT is way easier for implementation than RRT* (if we need an efficient implementation)
 - Path quality of RRT can be improved by fast post-processing
 - Asymptotic optimality is just asymptotic!
- slow convergence of path quality

- Sampling-based planning randomly samples \mathcal{C}
- Samples are classified as free/non-free, free samples are stored
- Multi-query vs. single-query planners
- PRM/RRT/EST and their optimal variants PRM* and RRT*