

**GRAPHICAL MARKOV MODELS (WS2014)**  
**5. SEMINAR**

**Assignment 1.** Consider the task of finding the most probable sequence of (hidden) states for a (Hidden) Markov model on a chain.

- a) Show that the Dynamic Programming approach applied for this task can be interpreted as an equivalent transformation (re-parametrisation) of the model.
- b) Show that the transformed functions (potentials) encode an explicit description of *all* optimisers of the problem

**Assignment 2.** Let  $(V, E)$  be an undirected tree and let  $s_i, i \in V$  be  $K$ -valued random variables associated with the vertices of that tree. Their joint distribution is a Markov model given by

$$p(s) = \frac{1}{Z} \prod_{ij \in E} g_{ij}(s_i, s_j),$$

where  $g_{ij}$  are non-negative factors.

- a) Find an algorithm for computing all marginal probabilities  $p(s_i), s_i \in K, i \in V$  of the model. *Hint:* See sec. 11 of the lecture, fill in details.
- b) Is it possible to parallelise the algorithm?
- c) Modify the algorithm such that it computes “max-marginals” of the distribution. The latter are defined by

$$m(s_i = k) = \max_{s: s_i=k} p(s).$$

**Assignment 3.** Consider a GRF for binary valued labellings  $x: V \rightarrow \{0, 1\}$  of a graph  $(V, E)$  given by

$$p(x) = \frac{1}{Z} \exp \left[ \sum_{ij \in E} u_{ij}(x_i, x_j) + \sum_{i \in V} u_i(x_i) \right].$$

Show that is is always possible to find an equivalent transformation (re-parametrisation)

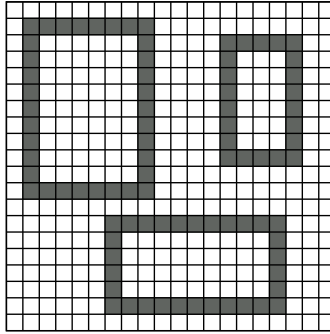
$$u_{ij} \rightarrow \tilde{u}_{ij}, \quad u_i \rightarrow \tilde{u}_i$$

such that the new pairwise functions  $\tilde{u}_{ij}$  have the form

$$\tilde{u}_{ij}(x_i, x_j) = \alpha_{ij} |x_i - x_j|$$

with some reals  $\alpha_{ij} \in \mathbb{R}$ .

**Assignment 4.** Consider the language  $L$  of all b/w images  $x: D \rightarrow \{b, w\}$  containing an arbitrary number of non-overlapping and non-touching one pixel wide rectangular frames (see figure).



**a)** Prove that  $L$  is not expressible by a locally conjunctive predicate

$$x \in L \quad \text{if and only if} \quad f(x) = \bigwedge_{c \in \mathcal{C}} f_c(x_c) = 1$$

with predicates  $f_c$ , defined on image fragments  $x_c$ , where  $c \subset D$  have bounded size  $|c| \leq m$ .

**b)** Show that  $L$  can be expressed by introducing a field  $s: D \rightarrow K$  of non-terminal symbols, a locally conjunctive predicate for them and pixel-wise predicates  $g$  “connecting” the non-terminal and terminal symbol in each pixel

$$x \in L \quad \text{if and only if} \quad \bigvee_{s \in K^D} \left[ \bigwedge_{c \in \mathcal{C}} f_c(s_c) \wedge \bigwedge_{i \in D} g(x_i, s_i) \right] = 1$$

Find a suitable structure  $\mathcal{C}$ , alphabet of non-terminal symbols  $K$  and predicates  $f_c, g$ .