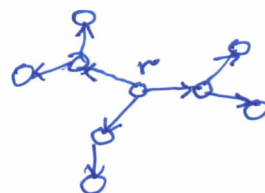


11. Hidden Markov Models on acyclic graphs

Notation

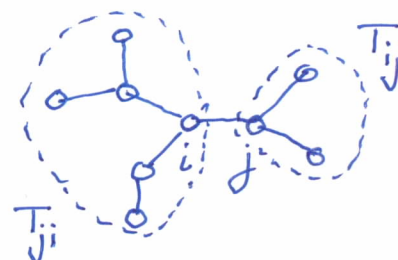
Let $T = (V, E)$ be an undirected, connected acyclic graph

- Fixing an arbitrary vertex $r \in V$, we denote by \vec{E}_r the edge set of the corresponding rooted digraph



- For each oriented edge ij , $\{ij\} \in \vec{E}$ we define the subtree T_{ij} by

$$V(T_{ij}) = \{m \in V \mid \text{path}(i, m) \ni j\}$$



1.4. The model

Definition 1a Let $T = (V, E)$ be an undirected tree and $s_i, i \in V$ be K -valued random variables. A p.d. for the random field $s \in K^{|V|}$ is a Markov model on T if

$$p(s) = p(s_r) \prod_{j \in \vec{E}_r} p(s_j | s_i)$$

holds for any choice $r \in V$ (root). ■

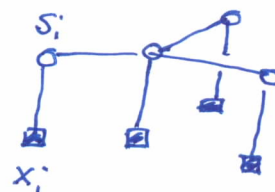
Definition 1b Equivalently, a p.d. $p(s)$ for $s \in K^{|V|}$ is a Markov model on T if

$$p(s) = \prod_{\{ij\} \in E} p(s_i, s_j) / \prod_{i \in V} p^{n_i-1}(s_i)$$

holds, where n_i is the degree of vertex $i \in V$. ■

An HMM on an undirected tree $T = (V, E)$ is a p.d. for pairs $s \in K^{|V|}$, $x \in F^{|V|}$, where F is a feature space and $p(x, s) = p(x|s)p(s)$, where

- $p(s)$ is a Markov model on T
- $p(x|s) = \prod_{i \in V} p(x_i | s_i)$
(conditional independence)



B. Inference tasks

Consider the following task: Given an observation field $x \in F^{|V|}$, compute

$$p(x) = \sum_{s \in K^{|V|}} p(x, s)$$

By substituting the model Def. 1a we get

$$p(x) = \sum_{s \in K^{|V|}} p(s_r) p(x_r | s_r) \prod_{j \in \vec{E}_r} p(s_j | s_i) p(x_j | s_j)$$

The problem has the form (fixed observation!)

$$\sum_{s \in K^{|V|}} \prod_{i \in V} \varphi_i(s_i) \prod_{j \in \vec{E}_r} f_{ij}(s_i, s_j)$$

where:

$$\varphi_i(s_i) = \begin{cases} p(s_r) p(x_r | s_r) & \text{if } i=r \\ 1 & \text{otherwise} \end{cases}$$

$$f_{ij}(s_i, s_j) = p(s_j | s_i) p(x_j | s_j)$$

The algorithm recomputes the φ -s starting from an arbitrary leaf $j \in V$. Let $ij \in \vec{E}_r$ denote the only incoming edge.

$$\varphi_i(s_i) := \varphi_i(s_i) \sum_{s_j \in K} \varphi_{ij}(s_i, s_j) \varphi_j(s_j)$$

Then the leaf j is removed. This is repeated until only the root r remains. Finally

$$p(x) = \sum_{k \in K} \varphi_r(k)$$

Complexity: $|K|^2 |E|$

Remark 1 The same approach is applied for solving the task

$$s_* \in \operatorname{argmax}_{s \in K^{|V|}} \log p(x, s)$$

Simply by replacing operations

$$x \mapsto +$$

$$+ \mapsto \max$$

C. Computing marginals

Consider the following task: Given an observation field $x \in F^{|V|}$, compute marginal prob's

$$p(x, s_i), \quad \forall i \in V, \quad \forall s_i \in K$$

Recall the algorithm for computing marginals of an HMM on a chain (see sec. 5)

Here:

• It follows from Def 16 that

$$p(x, s_i) = p(s_i) p(x_i | s_i) \prod_{j \in \mathcal{U}_i} p(x_{T_{ij}} | s_i)$$

- Let us denote $\varphi_{ij}(s_i) = p(x_{T_{ij}} | s_i)$. The φ -s fulfil the following system of equations

$$\varphi_{ij}(s_i) = \sum_{s_j} p(s_j | s_i) p(x_j | s_j) \prod_{\substack{l \in N_j \\ l \neq i}} \varphi_{je}(s_j)$$

- Two passes through all edges of T are needed to compute all of them \Rightarrow complexity $2|K|^2|E|$

1. Learning the tree structure

Given: i.i.d. sample of realisations $s \in K^{|V|}$

Task: estimate parameters and structure of the model

Denote by $\beta(s)$ the empirical p.d. assoc. with the sample.

MLE for known structure $T = (V, E)$

$$\sum_{s \in K^{|V|}} \beta(s) \log p(s) \stackrel{\text{Def. 16}}{=} \dots$$

$$= \sum_{\{i,j\} \in E} \sum_{s_i, s_j} \beta(s_i, s_j) \log p(s_i, s_j) - \sum_{i \in V} (n_i - 1) \sum_{s_i} \beta(s_i) \log p(s_i)$$

Maximiser given by $p(s_i, s_j) = \beta(s_i, s_j) \forall \{i,j\} \in E$

$$p(s_i) = \beta(s_i) \forall i \in V$$

It remains to solve

$$\sum_{\{i,j\} \in E} \sum_{s_i, s_j} \beta(s_i, s_j) \log \beta(s_i, s_j) - \sum_{i \in V} (n_i - 1) \sum_{s_i} \beta(s_i) \log \beta(s_i) \rightarrow$$

$\rightarrow \max_{E: \text{tree}}$

Both terms depend on the edge structure E . Rearrange:

$$\sum_{\{i,j\} \in E} \left\{ \sum_{s_i, s_j} \beta(s_i, s_j) \log \beta(s_i, s_j) - \sum_{s_i} \beta(s_i) \log \beta(s_i) - \sum_{s_j} \beta(s_j) \log \beta(s_j) \right\} + \\ + \sum_{i \in V} \sum_{s_i} \beta(s_i) \log \beta(s_i)$$

The last term does not depend on the edge structure. Denote the values of the curly brackets by h_{ij} :

$$\sum_{\{i,j\} \in E} h_{ij} \rightarrow \max_{E: \text{tree}}$$

\Rightarrow Maximum spanning tree problem. Can be solved in polyn. time

Remark 1 Similar problem for chain \rightarrow travelling salesman problem \rightarrow NP hard.