

3. Recognising the generating model - computing the emission probability for a feature sequence

Let a, b index two HMMs on $F^n \times K^n$

$$a: p_a(s_1), p_a(s_i | s_{i-1}), p_a(x_i | s_i)$$

$$b: p_b(s_1), p_b(s_i | s_{i-1}), p_b(x_i | s_i)$$

Given: observed sequence $x \in F^n$, $x = (x_1, \dots, x_n)$

Question: which model has generated x ?

Any reasonable answer is based on $p_a(x), p_b(x) \rightarrow$

Task: compute

$$\begin{aligned} P(x) &= \sum_{s \in K^n} P(x, s) \\ &= \sum_{s_1 \in K} \dots \sum_{s_n \in K} p(s_1) p(x_1 | s_1) \prod_{i=2}^n [p(s_i | s_{i-1}) p(x_i | s_i)] \end{aligned}$$

$$\text{Denote: } P_i(s_{i-1}, s_i) = p(s_i | s_{i-1}) p(x_i | s_i)$$

$$\varphi_1(s_1) = p(s_1) p(x_1 | s_1)$$

$$\varphi_n(s_n) = 1$$

Expression under the sum has the form

$$\varphi_1(s_1) \cdot P_2(s_1, s_2) \cdot \dots \cdot P_n(s_{n-1}, s_n) \cdot \varphi_n(s_n)$$

Summation can be performed iteratively from right to left

$$\varphi_{i-1}(s_{i-1}) = \sum_{s_i \in K} P_i(s_{i-1}, s_i) \varphi_i(s_i)$$

$$P(x) = \sum_{s \in K^n} \varphi_1(s_1) \varphi_n(s_n)$$

or from left to right

$$\varphi_i(s_i) = \sum_{s_{i-1} \in K}^1 \varphi_{i-1}(s_{i-1}) P_i(s_{i-1}, s_i)$$

$$P(x) = \sum_{s_n \in K}^1 \varphi_n(s_n) \varphi_n(s_n)$$

In matrix-vector forms we have

$$P(x) = \vec{\varphi}_1 \cdot P_2 \cdot P_3 \cdot \dots \cdot P_n \cdot \vec{\varphi}_n$$

complexity: $n|K|^2$

Remark The intermediate results have the following statistical meaning

$$\varphi_i(s_i) = p(x_{i+1}, \dots, x_n | s_i)$$

$$\varphi_i(s_i) = p(x_1, \dots, x_i, s_i)$$

4. Recognising the most probable sequence of hidden states

We observe $x = (x_1, \dots, x_n)$ for a known tHMM

Question Which sequence $s = (s_1, \dots, s_n)$ has generated x ?

If the answer is $s^* = \operatorname{argmax}_{S \in K^n} p(x, s)$, we have to solve

$$s^* = \operatorname{argmax}_{S \in K^n} \log p(x, s)$$

$$= \operatorname{argmax}_{S \in K^n} \left\{ \log [p(s_1) p(x_1 | s_1)] + \sum_{i=2}^n \log [p(s_i | s_{i-1}) p(x_i | s_i)] \right\}$$

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Denote: $M_i(s_{i-1}, s_i) = \log [p(s_i | s_{i-1}) p(x_i | s_i)]$

$$\psi_i(s_i) = \log [p(s_i) p(x_i | s_i)]$$

$$\psi_n(s_n) \equiv 0$$

Now, translate algorithm from sec. 3 by replacing operators $x \mapsto +$, $+$ $\rightarrow \max$. E.g. maximising dynamically from right to left

$$\psi_{i-1}(s_{i-1}) = \max_{s_i} [M_i(s_{i-1}, s_i) + \psi_i(s_i)]$$

$$\max_{s \in K^n} \log p(x, s) = \max_{s_i} [\psi_i(s_i) + \psi_n(s_i)]$$

We are looking for $s^* = \operatorname{argmax}_s p(x, s) \Rightarrow$

introduce $p\tau_i : K \rightarrow K$ for $i=1, \dots, n-1$

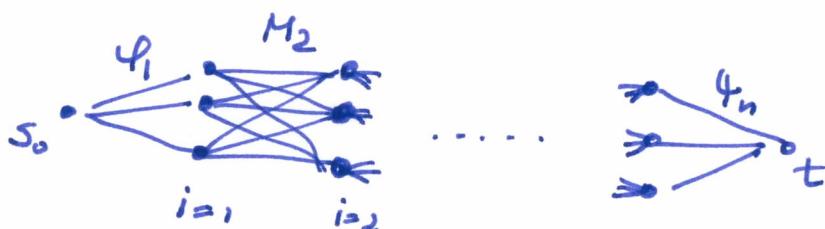
$$p\tau_{i-1}(s_{i-1}) = \operatorname{argmax}_{s_i} [M_i(s_{i-1}, s_i) + \psi_i(s_i)]$$

Backtrack an optimiser s^* starting from

$$s_1^* = \operatorname{argmax}_{s_1} [\psi_1(s_1) + \psi_n(s_1)]$$

Similarly, we can maximize from left to right

Both variant search the best path in the graph



Complexity: $n|K|^2$

5. Recognising the sequence of most probable hidden states

Is $s^* = \underset{s \in K^n}{\operatorname{argmax}} p(x, s)$ the best answer to the question posed in sec. 4?

Answer depends on $p(x, s)$ and the loss $C(s', s)$

$$s^* = \underset{s \in K^n}{\operatorname{argmin}} \sum_{s' \in K^n} p(x, s') C(s', s)$$

a) If $C(s', s) = \mathbb{1}\{s' \neq s\}$, then $s^* = \underset{s \in K^n}{\operatorname{argmax}} p(x, s)$

b) If $C(s', s) = \sum_{i=1}^n \mathbb{1}\{s'_i \neq s_i\}$, i.e. Hamming distance?

$$s^* = \underset{s \in K^n}{\operatorname{argmin}} \sum_{s' \in K^n} p(x, s') \sum_{i=1}^n (1 - \delta(s'_i, s_i))$$

$$= \underset{s \in K^n}{\operatorname{argmax}} \sum_{i=1}^n \sum_{s' \in K^n} p(x, s') \delta(s'_i, s_i)$$

$$= \underset{s \in K^n}{\operatorname{argmax}} \sum_{i=1}^n p(x, s_i)$$

$$\hookrightarrow \boxed{s_i^* = \underset{s_i}{\operatorname{argmax}} p(x, s_i)} \quad i = 1, \dots, n$$

Problem: Compute prob's $p(x, s_i = k) + i=1, \dots, n, \forall k \in K$

Answer: Remember quantities φ, ψ from sec. 3

$$\varphi_i(s_i) = p(x_1, \dots, x_i, s_i)$$

$$\psi_i(s_i) = p(x_{i+1}, \dots, x_n | s_i)$$

Since $p(x|s_i) = p(x_1, \dots, x_i|s_i) p(x_{i+1}, \dots, x_n|s_i)$ holds for an HMM (!), we have

$$p(x, s_i) = \psi_i(s_i) \psi_i(s_i)$$

Complexity: $2n|K|^2$