AE4M33RZN, Fuzzy logic: Tutorial examples

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Assignment

On the universe $\Delta = \{a, b, c, d\}$ there is a fuzzy set

$$\mu_{\rm A} = \{(a; {\rm 0.3}), (b; {\rm 1}), (c; {\rm 0.5})\}.$$

Find its horizontal representation.

$$R_{A}(\alpha) = \begin{cases} X & \alpha = 0, \\ \{a, b, c\} & \alpha \in (0; 0.3), \\ \{b, c\} & \alpha \in (0.3; 0.5), \\ \{b\} & \alpha \in (0.5; 1). \end{cases}$$

Assignment

The fuzzy set A has a horizontal representation

$$R_{A}(\alpha) = \begin{cases} \{a, b, c, d\} & \alpha \in \langle 0, 1/3 \rangle \\ \{a, d\} & \alpha \in (1/3, 1/2) \\ \{d\} & \alpha \in (1/2, 2/3) \\ \emptyset & \alpha \in (2/3, 1) \end{cases}$$

Find the vertical representation.

Solution

$$\begin{split} \mu_A(a) &= \sup \bigl\{ \alpha \in \langle \mathbf{o}, \mathbf{1} \rangle : \ a \in \mathbb{R}_A(\alpha) \bigr\} = \sup \langle \mathbf{o}, \mathbf{1/2} \rangle = \mathbf{1/2} \\ \mu_A(b) &= \mathbf{1/3}, \quad \mu_A(c) = \mathbf{1/3}, \quad \mu_A(d) = \mathbf{2/3}, \quad \text{therefore} \\ \mu_A &= \bigl\{ (a, \mathbf{1/2}), (b, \mathbf{1/3}), (c, \mathbf{1/3}), (d, \mathbf{2/3}) \bigr\} \end{split}$$

Assignment

On the universe $\Delta = \mathbb{R}$ there is a fuzzy set A:

$$\mu_{A}(x) = \begin{cases} x & x \in \langle 0; 1 \rangle \\ 2 - x & x \in \langle 1; 1.5 \rangle \\ 0 & \text{otherwise} \end{cases}$$

Find its horizontal represenation.

Solution

$$R_{A}(\alpha) = \begin{cases} \mathbb{R} & \alpha = 0 \\ \langle \alpha; 1.5 \rangle & \alpha \in (0; 0.5) \\ \langle \alpha; 2 - \alpha \rangle & \alpha \in (0.5; 1) \end{cases}$$

Assignment

The fuzzy set A has a horizontal representation

$$R_A(\alpha) = \begin{cases} \mathbb{R} & \alpha = 0 \\ \langle \alpha^2, \mathbf{1} \rangle & \text{otherwise.} \end{cases}$$

Find the vertical representation.

$$\mu_A(\mathbf{x}) = \begin{cases} \sqrt{\mathbf{x}} & \mathbf{x} \in (\mathbf{0}, \mathbf{1}) \\ \mathbf{0} & \text{otherwise.} \end{cases}$$

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Assignment

Decide if the following function is a fuzzy conjunction.

$$\alpha \wedge \beta = \begin{cases} \alpha & \beta = 1, \\ \beta & \alpha = 1, \\ \alpha\beta & \alpha\beta \ge 1/10, \max(\alpha, \beta) < 1, \\ 0 & \text{otherwise.} \end{cases}$$

Solution

The first two possibilities ensure the boundary condition. Comutativity and monotonicity are trivially satisfied. If one of the arguments is 1, the associativity as well. For α , β , γ < 1 the associativity follows from

$$\alpha \wedge \left(\beta \wedge \gamma\right) = \begin{cases} \alpha\beta\gamma & \alpha\beta\gamma \geq 1/10. \\ 0 & \text{otherwise,} \end{cases}$$

We get the same for $(\alpha \land \beta) \land \gamma$.

It is always a fuzzy conjunction (interpretable as an algebraic conjunction, in which we ignore small values, e.g. for filtering small values).

Assignment

Prove that the Sugeno class of negations are all fuzzy negations

$$_{S_{\lambda}}^{\neg} \alpha = \frac{1-\alpha}{1+\lambda\alpha}$$
 for $\lambda \in (-1,\infty,)$.

Solution

When canceling out terms in a fraction, we use $1 + \lambda \alpha > 0$.

• Involutivity: Assuming $\lambda > -1$:

$$\overline{s_{\lambda}} \, \overline{s_{\lambda}} \, \alpha = \frac{\mathbf{1} - \frac{\mathbf{1} - \alpha}{\mathbf{1} + \lambda \alpha}}{\mathbf{1} + \lambda \, \frac{\mathbf{1} - \alpha}{\mathbf{1} + \lambda \alpha}} = \frac{\mathbf{1} + \lambda \alpha - \mathbf{1} + \alpha}{\mathbf{1} + \lambda \alpha + \lambda - \lambda \alpha} = \frac{\alpha \, (\lambda + \mathbf{1})}{\lambda + \mathbf{1}} = \alpha$$

• Non-increasing: If $\alpha \leq \beta$, then

$$\frac{1-\alpha}{1+\lambda\alpha} \ge \frac{1-\beta}{1+\lambda\beta}$$
$$(1+\lambda\beta)(1-\alpha) \ge (1+\lambda\alpha)(1-\beta)$$
$$\beta \ge \alpha$$

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Assignment

Decide if the following function is a fuzzy conjunction.

$$\alpha \circ \beta = \begin{cases} \alpha \beta & \alpha \beta \ge \text{ o,o1 or } \max(\alpha, \beta) = 1, \\ \text{o} & \text{otherwise} \end{cases}$$

Assignment

Decide if the following function is a fuzzy conjunction.

$$\bigwedge_{O}: \langle O, 1 \rangle^2 \to \langle O, 1 \rangle$$

$$\alpha \wedge \beta = \begin{cases} \min(\alpha, \beta) & \alpha + \beta \ge 1 \\ 0 & \text{otherwise} \end{cases}$$

Assignment

Decide, if for all $\alpha, \beta \in [0,1]$ holds $(\alpha \overset{\circ}{\vee} \beta) \overset{\wedge}{\underset{L}{\wedge}} (\alpha \overset{\circ}{\vee} \overset{\neg}{\underset{S}{\cap}} \beta) = \alpha$, where the disjunction $\overset{\circ}{\vee}$ is

- 1. standard, $\overset{S}{\lor}$,
- 2. algebraic, ♦.
- 3. Łukasiewicz, $\overset{\text{L}}{\forall}$,

Solution

- 1. No. Counterexample: $\alpha = 0.5$, $\beta = 0.1$.
- 2. Yes: $\max(o, (\alpha + \beta \alpha\beta) + (\alpha + (1 \beta) \alpha(1 \beta)) 1) =$ $= \max(o, \alpha) = \alpha$
- 3. No. Counterexample: $\alpha = 0.9$, $\beta = 0.8$.

Assignment

Decide if the function $\wedge : \langle 0, 1 \rangle^2 \rightarrow \langle 0, 1 \rangle$

$$\alpha \wedge \beta = \begin{cases} \alpha \beta & \alpha + \beta \ge 1 \\ 0 & \text{otherwise} \end{cases}$$

is a fuzzy conjunction.

Assignment

Decide if the $(\alpha \wedge \alpha) \vee (\alpha \wedge \alpha) \leq \alpha$ holds for

- 1. standard, $\overset{S}{\vee}$
- 2. algebraic, ♦
- 3. Łukasiewicz, $\overset{\text{L}}{\forall}$

Assignment

Decide, which equalities hold:

1.
$$(\alpha \underset{S}{\wedge} \alpha) \overset{L}{\vee} (\alpha \underset{S}{\wedge} \beta) = \alpha \underset{S}{\wedge} (\alpha \overset{L}{\vee} \beta)$$

2.
$$(\alpha \wedge \alpha) \stackrel{S}{\vee} (\alpha \wedge \beta) = \alpha \wedge (\alpha \stackrel{S}{\vee} \beta)$$

3.
$$\alpha \overset{S}{\vee} (\alpha \underset{L}{\wedge} \beta) = \alpha \underset{L}{\wedge} (\alpha \overset{S}{\vee} \beta)$$

Justify your conclusions.

Assignment

Decide, which equalities hold:

1.
$$(\alpha \wedge \beta) \wedge \gamma = \alpha \wedge (\beta \wedge \gamma)$$

2.
$$_{S}(\alpha \stackrel{A}{\vee} \beta) = _{S} \alpha \stackrel{A}{\wedge} _{L} \stackrel{B}{\vee} \beta$$

3.
$$(\alpha \wedge \alpha) \stackrel{L}{\vee} _{\stackrel{\circ}{S}} \alpha = (_{\stackrel{\circ}{S}} \alpha \wedge _{\stackrel{\circ}{L}} _{\stackrel{\circ}{S}} \alpha) \stackrel{L}{\vee} \alpha$$

Justify your conclusions.

Assignment

Verify that $\alpha \land (\alpha \stackrel{\mathbb{R}}{\underset{\circ}{\circ}} \beta) = \alpha \land \beta$ holds for

- 1. algebraic ops.
- 2. standard ops.

Solution

We will show the solution for algebraic operations. The other ones are similar.

$$\alpha \wedge (\alpha \xrightarrow{\mathbb{R}} \beta) = \left\{ \begin{array}{l} \alpha \cdot \mathbf{1} = \alpha \text{ for } \alpha \leq \beta \\ \alpha \cdot \frac{\beta}{\alpha} = \beta \text{ for } \alpha > \beta \end{array} \right\} = \min(\alpha, \beta) = \alpha \wedge \beta$$

Assignment

Complete the table, so that R is a S-partial order.

R	а	b	с	d
а				
b	0.5			
С		0.3		
d		0.2		

Solution

Reflexivity implies 1's on the diagonal. S-partial order implies o's to non-zero elements symmetric over the main diagonal.

R	а	b	С	d
а	1	О	x'	y'
b	0.5	1	О	О
С	x	0.3	1	z'
d	y	0.2	Z	1

The transitivity implies e.g. $R(3,2) \ {}^{\wedge}_{S} \ R(2,1) \le R(3,1)$, which translates into a condition $\min(0.3,0.5) \le x$ Using this and similar conditions, we derive the subspace of all solutions: $z \le 0.2$, $x \ge 0.3$, $y \ge 0.2$, $\min(y,z') \le a$, x' = y' = 0.

Ex: Jim revisited

We will use the Lukasiewicz logic in the following examples ($\Box = \Box$, ...).

$$\langle jim : Male | o.9 \rangle$$
 (1)

$$\langle jim : Female | o.2 \rangle$$
 (2)

$$\langle \mathsf{Male} \sqcap \mathsf{Female} \sqsubseteq \bot | \mathbf{1} \rangle$$
 (3)

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The interpretation domain is
$$\Delta^{\mathcal{I}_1} = \Delta^{\mathcal{I}_2} = \{j\}, jim^{\mathcal{I}_1} = jim^{\mathcal{I}_2} = j$$
. Male $\mathcal{I}_1 = \{(j; \mathbf{o}.9)\}$ Male $\mathcal{I}_2 = \{(j; \mathbf{o}.9)\}$ Female $\mathcal{I}_3 = \{(j; \mathbf{o}.2)\}$

Ex: Jim revisited (check your knowledge)

Let's check the interpretation against the definitions...

$\mathcal{I} \vDash \tau$	$ au_{(1)}$	$\tau_{(2)}$	$ au_{(3)}$
$\mathcal{I}_{\scriptscriptstyle 1}$?	?	?
\mathcal{I}_{2}	?	?	?

Ex: Jim revisited (check your knowledge)

Let's check the interpretation against the definitions...

$\mathcal{I} \vDash \tau$	$ au_{(1)}$	$\tau_{(2)}$	$ au_{(3)}$
$\mathcal{I}_{\scriptscriptstyle 1}$	yes	no	yes
\mathcal{I}_{2}	yse	yes	no

Ex: Jim revisited (in fuzzyDL)

Let's change the weights and encode the example in fuzzyDL:

```
(instance jim Male 0.4)
(instance iim Female 0.2)
(1-implies (and Male Female) *bottom* 0.9)
(min-instance? jim Male)
(max-instance? iim Male)
(min-instance? iim Female)
(max-instance? jim Female)
```

Let $\langle jim : \mathsf{Male} \, | \, \alpha \rangle$ and $\langle jim : \mathsf{Female} \, | \, \beta \rangle$, what are the bounds on α and β ? fuzzyDL shows that 0.4 $\leq \alpha \leq$ 0.9 and 0.2 $\leq \beta \leq$ 0.7. Why?

Ex: Smokers

Recall the motivational example from the first lecture:

$$\langle \text{symmetric}(\text{friend}) \rangle \qquad (4)$$

$$\langle (anna, bill) : \text{friend} | 1 \rangle \qquad (5)$$

$$\langle (bill, cloe) : \text{friend} | 1 \rangle \qquad (6)$$

$$\langle (cloe, dirk) : \text{friend} | 1 \rangle \qquad (7)$$

$$\langle anna : \text{Smoker} | 1 \rangle \qquad (8)$$

$$\langle \exists \text{friend} \cdot \text{Smoker} \sqsubseteq \text{Smoker} | 0.7 \rangle \qquad (9)$$

What are the bounds on $\langle i : Smoker \rangle$ for $i \in \{anna, bill, cloe, dirk\}$?

Ex: Smokers

What changes if we add

$$\langle dirk : \neg Smoker | o.7 \rangle$$
 (10)

What are the bounds on $\langle i : \neg Smoker \rangle$ for $i \in \{anna, bill, cloe, dirk\}$?

Ex: Smokers (in fuzzyDL)

```
(implies (some friendOf Smoker) Smoker 0.7)
(summetric friendOf)
(related anna bill friendOf)
(related bill cloe friendOf)
(related cloe dirk friendOf)
(instance anna Smoker)
(instance dirk (not Smoker) 0.7)
(min-instance? anna Smoker)
(min-instance? bill Smoker)
(min-instance? cloe Smoker)
(min-instance? dirk Smoker)
(max-instance? anna Smoker)
(max-instance? bill Smoker)
(max-instance? cloe Smoker)
(max-instance? dirk Smoker)
```

The domain $\Delta^{\mathscr{I}}$ is an unordered set. This is good for modelling cathegorical data: e.g. colors, people, ...

General idea: Extended interpretation

But we also need to include real numbers \mathbb{R} . The fuzzy description logic with concrete datatypes \mathcal{RGHE} —Oses "abstract objects" and "concrete objects":

$$\Delta^{\mathcal{I}} = \Delta_a^{\mathcal{I}} \cup {\rm I\!R}$$

• Concrete individuals, are interpreted as objects from ${\rm I\!R}.$

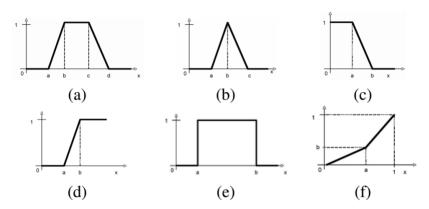
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- Concrete concepts, are interpreted as subsets from IR.
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All non-concrete notions are called abstract.

Concrete data types: New concepts



(a) Trapezoidal function; (b) Triangular function; (c) L-function; (d) R-function; (e) Crisp interval; (f) Linear function.

```
(related adam bob parent) (related adam eve parent)
(define-fuzzy-concept around23 triangular(0,100, 18,23,26))
(define-fuzzu-concept moreTh17 right-shoulder(0.100, 13.21))
(instance bob (some age around23) 0.9)
(instance eve (some age moreTh17))
(define-fuzzy-concept young left-shoulder(0,100, 17,25))
(define-concept YoungPerson (some age young))
(min-instance? eve YoungPerson) (max-instance? eve YoungPerson)
(min-instance? bob YoungPerson) (max-instance? bob YoungPerson)
(min-instance? adam (all parent YoungPerson))
(max-instance? adam (all parent YoungPerson))
(min-instance? adam (some parent YoungPerson))
(max-instance? adam (some parent YoungPerson))
```

1. What are the bounds on α from $\langle eve : YoungPerson | \alpha \rangle$?

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Start by drawing the concept around 23, then construct an interpretation. How much freedom do you have when constructing the interpretation?

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Start by drawing the concept around 23, then construct an interpretation. How much freedom do you have when constructing the interpretation?

2. Let fuzzyDL reasoner give you both bounds on $\langle i: YoungPerson | \beta_i \rangle$ for $i \in \{eve, bob\}$.

How do you infer the bounds on $\langle adam : YoungPerson | \gamma \rangle$?

- 1. The buyer wants a passenger that costs less than €26000.
- 2. If there is an alarm system in the car, then he is satisfied with paying no more than €22300, but he can go up to €22750 with a lesser degree of satisfaction.
- The driver insurance, air conditioning and the black color are important factors.
- 4. Preferably the price is no more than €22000, but he can go to €24000 to a lesser degree of satisfaction.

- 1. The seller wants to sell no less than €22000.
- 2. Preferably the buyer buys the **insurance plus** package.
- 3. If the color is black, then it is highly possible the car has an air-conditioning.

This can be formalized in fuzzy description logic.

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We have the background knowledge:

 $\langle InsurancePlus = DriverInsurance \sqcap TheftInsurance | 1 \rangle$

The buyer's preferences:

1. B = PassengerCar ⊓ ∃ price · ≤ 26000

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- 3. B_2 = DriverInsurance,

- 1. $B = PassengerCar \sqcap \exists price \cdot \leq 26000$
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The buyer's preferences:

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- 4. $B_5 = \exists \text{ price } \cdot l.sh.(22000, 24000)$

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The buyer's preferences:

- 1. $B = PassengerCar \sqcap \exists price \cdot \leq 26000$
- 2. $B_1 = \text{AlarmSystem} \mapsto \exists \text{ price } \cdot \text{ l.sh.}(22300, 22750)$
- 3. $B_2 = \text{DriverInsurance}, B_3 = \text{AirCondition}, B_4 = \exists \text{color} \cdot \text{Black}$
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- 1. $S = PassengerCar \sqcap \exists price \cdot \geq 22000$
- 2. $S_1 = InsurancePlus$
- 3. $S_2 = (0.5 \ (\exists \ color \cdot \ Black) \mapsto AirCondition)$

We know that S and B are hard constraints and $B_{1..5}$ and $S_{1..2}$ are soft preferences. All the concepts can be "summed up":

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Buy =
$$B \sqcap (o.1B_1 + o.2B_2 + o.1B_3 + o.4B_4 + o.2B_5)$$

and

$$Sell = S \sqcap (0.6S_1 + 0.4S_2)$$

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A good choice of \square can make B a hard constraint.

Optimal match

$$bsb(K, Buy \sqcap Sell)$$

Finds the optimal match between a seller and a buyer. (Finds an ideal, imaginary car that maximizes satisfaction of both parties.)

Particular car

$$bdb(K, \langle audiTT : Buy \sqcap Sell \rangle)$$

Finds the degree of satisfaction for a particuklar car audiTT.