AE4M33RZN, Fuzzy description logic: fuzzyDL reasoner

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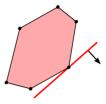
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Linear programming in a nutshell

Imagine a 2D space with a convex polygon in the space (x, y). Given constraints $4x + y \ge 6, y \le 8, ...,$ minimize x - 2y.



Source: [Wikipedia, 2013]

Usually written in a matrix form

maximize
$$c^T \cdot x$$
(1)subject to $A x \le o$ (2)

- (Mixed) Integer LP allows (some) variables to be discrete.
- LP with real values is in P class, ILP is NP-complete.

Linear programming in a nutshell

Solution of a ((M)I)LP

- One solution (a point in the polytope).
- No solution (the polytope is empty).
- Multiple solutions with equal objective function value.

Syntactical notes about fuzzyDL:

- $x \in \mathbb{R}$ will be real numbers.
- $y \in \mathbb{I}\mathbb{N}$ will be integer numbers.
- All values *x*, *y* will be bounded by [0, 1].

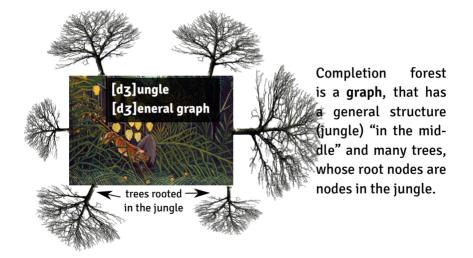
FuzzyDL algorithm overview

- Transforms ${\mathscr K}$ to the <code>negated-normal-form.1</code>
- Creates an witnessed interpretation of ${\mathcal K}$.
- During its working it creates
 - a completion forest and
 - a list of linear constraints \mathscr{C} .
- Linear constraints \mathscr{C} are solved using any mixed-integer-linear-programming solver.

Disclaimer: Not going beyond Ł-logic, no concrete data types.

¹Makes sure that the negation \neg appears only in front of concepts using: nnf($\neg \forall R \cdot C$) = $\exists R \cdot nnf(\neg C)$ and $nnf(\neg \exists R \cdot C) = \forall R \cdot nnf(\neg C)$.

Completion-forest informally



Completion-forest formally

The fuzzyDL algorithm starts with creating the "jungle". It contains all **individuals** (connected by an edge if they are linked by some relation).

Initialization

- Create a new vertex v_a for each **individual** a in the \mathcal{K} .
- Create an edge (v_a, v_b) for each role assertion between a and b.
- Add a label $\langle C, n \rangle$ to vertex *a* for each concept assertion $\langle a : C | n \rangle$.
- Add a label $\langle \mathbb{R}, n \rangle$ to edge (a, b)for each role assertion $\langle (a, b) : \mathbb{R} | n \rangle$.

Forest completion (1)

The reasoner applies each of the following rules sequentially:

- A If a vertex v is labeled $\langle C, l \rangle$, add $(x_{v:C} \ge l)$ into \mathscr{C} .
- \overline{A} If a vertex ν is labeled $\langle \neg C, l \rangle$, add $(x_{\nu:C} \leq i l)$ into \mathscr{C} .
- $\mathsf{R} \ \text{ If an edge } (\nu, w) \text{ is labeled } \langle \mathsf{R}, l \rangle \text{, add } (x_{(\nu, w): \mathsf{R}} \geq l) \text{ into } \mathscr{C}.$
- \perp If a vertex ν is labeled $\langle \perp, l \rangle$, add (l = o) into \mathscr{C} .

Forest completion (2)

□ If a vertex v is labeled $\langle C \sqcap D, l \rangle$, append labels $\langle C, x_1 \rangle$, $\langle D, x_2 \rangle$ to v and add the following constraints into \mathscr{C} (with fresh x_1, x_2, y):

$$y \le 1 - l$$

$$x_1 \le 1 - y$$

$$x_2 \le 1 - y$$

$$x_1 + x_2 = l + 1 - y$$

 $\begin{tabular}{ll} & \sqcup \begin{tabular}{ll} If a vertex v is labeled $\langle C \sqcup D$, l\rangle$, append labels $\langle C, x_1 \rangle$, $\langle C, x_2 \rangle$ to v and add $(x_1 + x_2 = l)$ into \varnotharpow (with fresh x_1, x_2, y)$. \end{tabular}$

Forest completion (3)

∀ If a vertex v is labeled $\langle \forall R \cdot C, l_1 \rangle$, an edge (v, w) is labeled $\langle R, l_2 \rangle$ and the rule has not been applied to this pair, then append the label $\langle C, x \rangle$ to w and add the following constraints into \mathscr{C} (with fresh x, y):

$$l_1+l_2-1\leq x\leq y\leq l_1+l_2$$

Forest completion: Example

Consider $\mathscr{K} = \{ \langle \exists \mathbb{R} \cdot \mathbb{C} \sqsubseteq \mathbb{D} | \mathbf{i} \rangle, \langle (a, b) : \mathbb{R} | \mathbf{0.7} \rangle, \langle b : \mathbb{C} | \mathbf{0.8} \rangle \}.$ Show that $glb(\mathscr{K}, a : \mathbb{D}) = \mathbf{0.5}.$

Termination (1)

Unless the rules are applied repeatedly, the algorithm (as explained so far) terminates.

For defining \exists rule, new nodes are added, which needs to refine the terminating condition.

Equivalence of labels

Two lists of labels $[\langle C_1, l_1 \rangle, ..., \langle C_n, l_n \rangle]$ and $[\langle C_1, l_1' \rangle, ..., \langle C_n, l_n' \rangle]$ are equivalent iff either

- l_i and l'_i are variables or
- *l_i* and *l'_i* are negated variables or
- l_i and l'_i are equal rationals.

Termination (2)

Directly blocked node

A node is directly blocked iff

- it is outside the "jungle" and
- none of its ancestors are blocked and
- it has an ancestor with equivalent labels.

Blocked node

A node is blocked iff either

- it is directly blocked or
- one of its predecessors is blocked.

Forest completion (4)

∃ If a vertex v is labeled $\langle \exists R \cdot C, l \rangle$ and it is not blocked, add a new vertex w and an edge (v, w), add labels $\langle C, x_2 \rangle$ to w, and $\langle R, x_1 \rangle$ to (v, w) and the following constraints into \mathscr{C} (with fresh x_1, x_2 and y):

$$y \le 1 - l$$

$$x_1 \le 1 - y$$

$$x_2 \le 1 - y$$

$$x_1 + x_2 = l + 1 - y$$

FuzzyDL: Overview

- The instance of MILP is created using constraints ${\mathscr C}$.
- In order to solve $glb(\mathcal{K}, \langle a : C \rangle)$, the objective function is set to minimize x in the MILP instance created for an augmented knowledge base $\mathcal{K} \cup \langle a : \neg C | \mathbf{1} x \rangle$.
- Similarly for $glb(\mathcal{K}, \langle a : C \sqsubseteq D \rangle)$ the augmented knowledge base is $\mathcal{K} \cup \langle a : C \sqcap \neg D | 1 - x \rangle$.
- ${\mathscr R}$ is inconsistent iff the MILP instance has no solution. Hence the $glb(\cdot,\cdot)$ is found if MILP instance has a solution.

Concrete data types

The domain $\Delta^{\mathscr{S}}$ is an unordered set. This is good for modelling cathegorical data: e.g. colors, people, ...

General idea: Extended interpretation

But we also need to include real numbers \mathbb{R} . The *fuzzy description logic with concrete datatypes* $SHF(\mathcal{D})$ uses "abstract objects" and "concrete objects":

$$\Delta^{\mathscr{I}} = \Delta^{\mathscr{I}}_{a} \cup \mathbb{R}$$

Concrete data types

- *Concrete individuals*, are interpreted as objects from **R**.
- Concrete concepts, are interpreted as subsets from IR.
- Concrete roles, are interpreted as subsets from ($\Delta_a^{\mathscr{I}} \times \mathbb{R}$).

All non-concrete notions are called *abstract*.

Concrete data types: New concepts

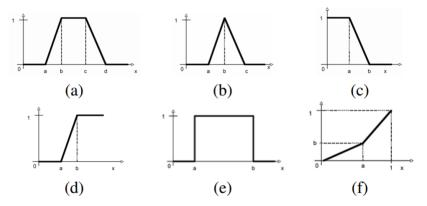


Fig. 1. (a) Trapezoidal function; (b) Triangular function; (c) *L*-function; (d) *R*-function; (e) Crisp interval; (f) Linear function.

Ex: Age of parents

(related adam bob parent) (related adam eve parent)

```
(define-fuzzy-concept around23 triangular(0,100, 18,23,26))
(define-fuzzy-concept moreTh17 right-shoulder(0,100, 13,21))
(instance bob (some age around23) 0.9)
(instance eve (some age moreTh17))
```

(define-fuzzy-concept young left-shoulder(0,100, 17,25))
(define-concept YoungPerson (some age young))

```
(min-instance? eve YoungPerson) (max-instance? eve YoungPerson)
(min-instance? bob YoungPerson) (max-instance? bob YoungPerson)
(min-instance? adam (all parent YoungPerson))
(max-instance? adam (all parent YoungPerson))
(min-instance? adam (some parent YoungPerson))
(max-instance? adam (some parent YoungPerson))
```

Ex: Age of parents

1. What are the bounds on α from $\langle eve : YoungPerson | \alpha \rangle$?

Start by drawing the concept around23, then construct an interpretation. How much freedom do you have when constructing the interpretation?

2. Let fuzzyDL reasoner give you both bounds on $\langle i : \text{YoungPerson} | \beta_i \rangle$ for $i \in \{eve, bob\}$.

How do you infer the bounds on $\langle adam :$ YoungPerson $|\gamma\rangle$?

Questions?! Ask, please.



Source: ragtagdoodles.deviantart.com

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