

## ALG 04

### **Queue**

**Operations Enqueue, Dequeue, Front, Empty....**

**Cyklic queue implementation**

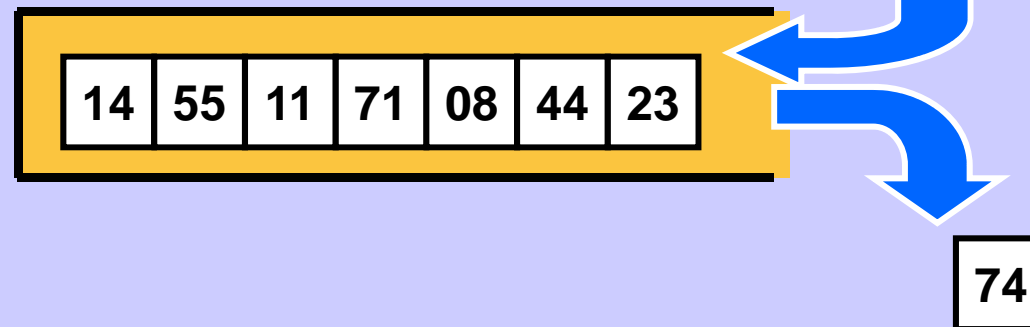
**Breadth-first search (BFS) in a tree**

# Stack

Elements are stored at the stack top before they are processed.

Stack bottom

Stack top



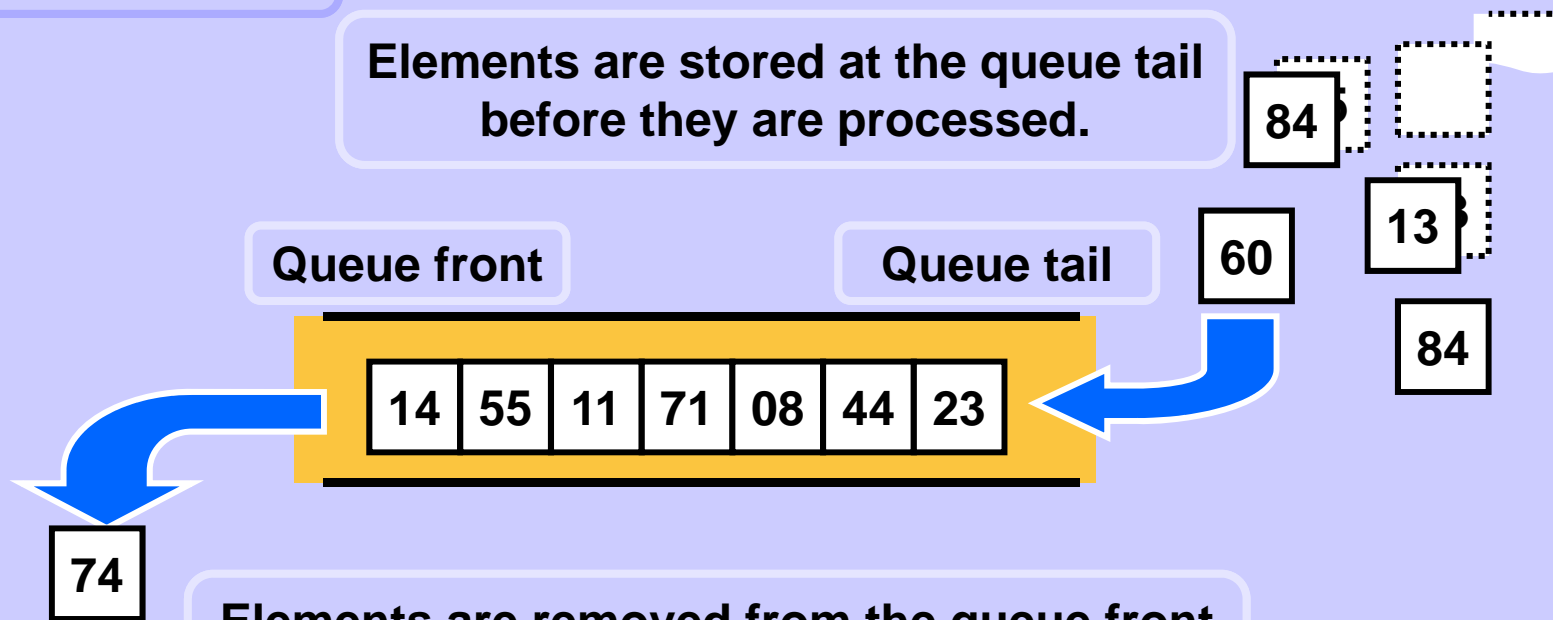
Elements are removed from the stack top and then they are processed.

## Operation names

Put at the top	Push
Remove from the top	Pop
Read the top	Top
Is the stack empty?	Empty

## Queue

Elements are stored at the queue tail before they are processed.



Elements are removed from the queue front and then they are processed.

### Operation names

Insert at the tail

Enqueue / InsertLast / Push ...

Remove from the front

Dequeue / delFront / Pop ...

Read the front elem

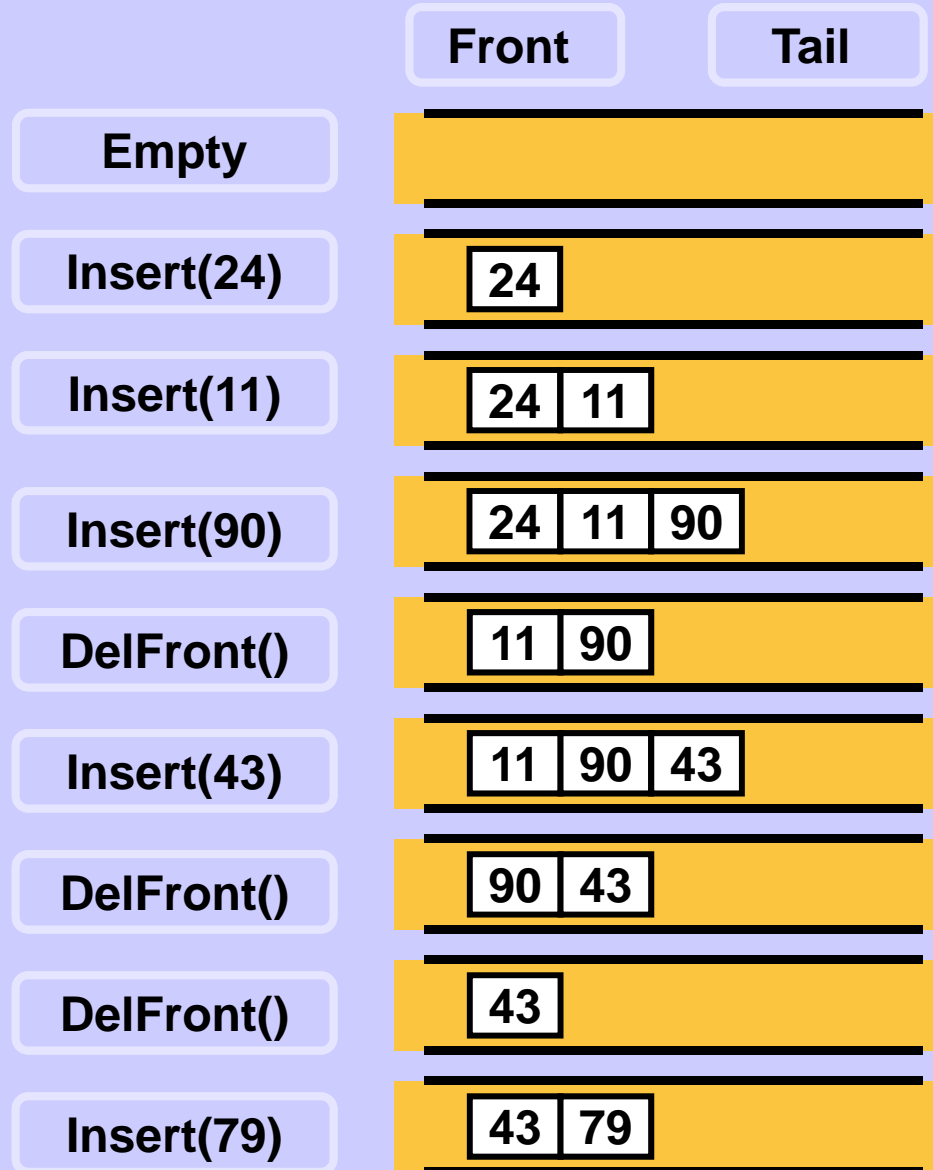
Front / Peek ...

Is the queue empty?

Empty

# Queue

Easy example  
of a queue  
life cycle.



## Cyclic queue implementation in an array

An empty queue in a fixed length array

Insert 24, 11, 90, 43, 70.

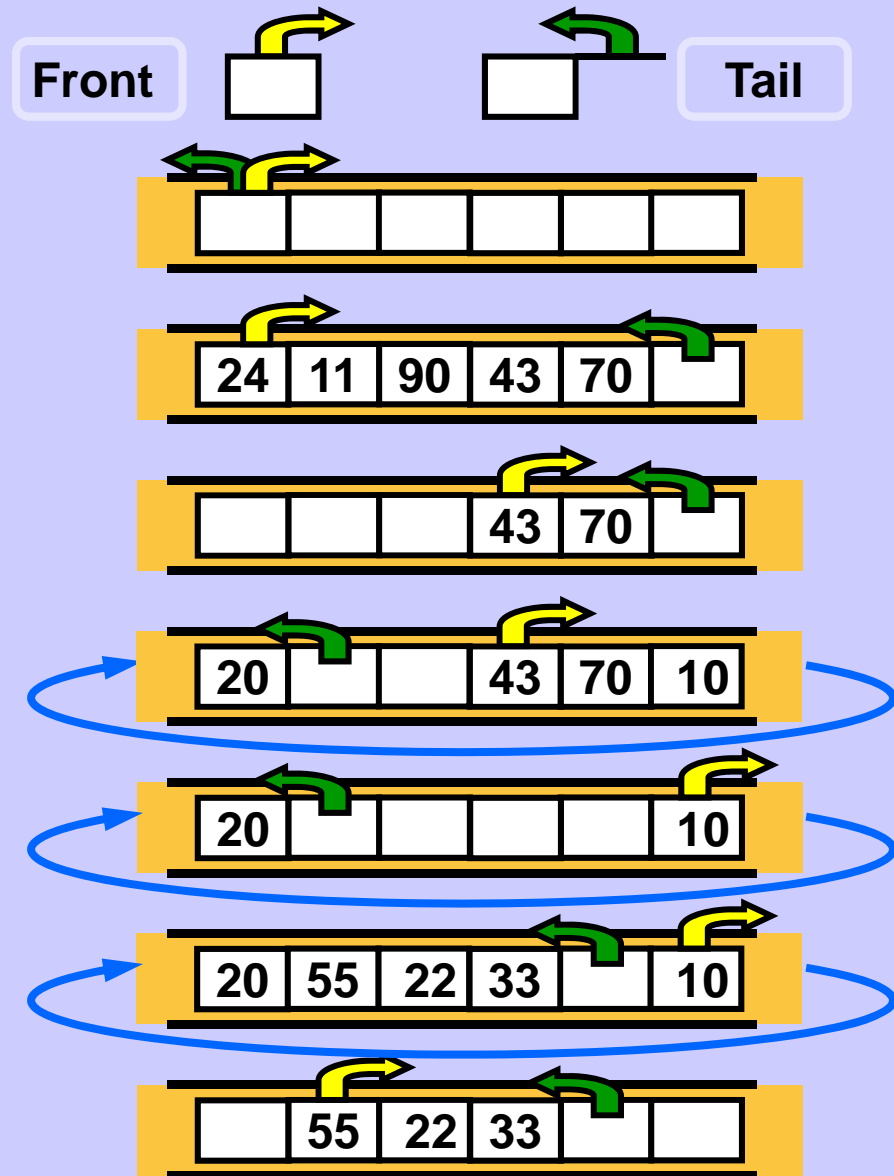
DelFront, DelFront, DelFront .

Insert 10, 20.

DelFront, DelFront .

Insert 55, 22, 33.

DelFront, DelFront .



## Cyclic queue implementation in an array

Tail index points to the first free position behind the last queue element.  
 Front index points to the first position occupied by a queue element.  
 When both indices point to the same position the queue is empty.

```

class Queue {
    Node q [];
    int size;
    int front;
    int tail;

    //constructor:
    Queue(int qsize) {
        size = qsize;
        q = new Node[size];
        front = 0;
        tail = 0;
    }

    boolean Empty() {
        return (tail==front);
    }

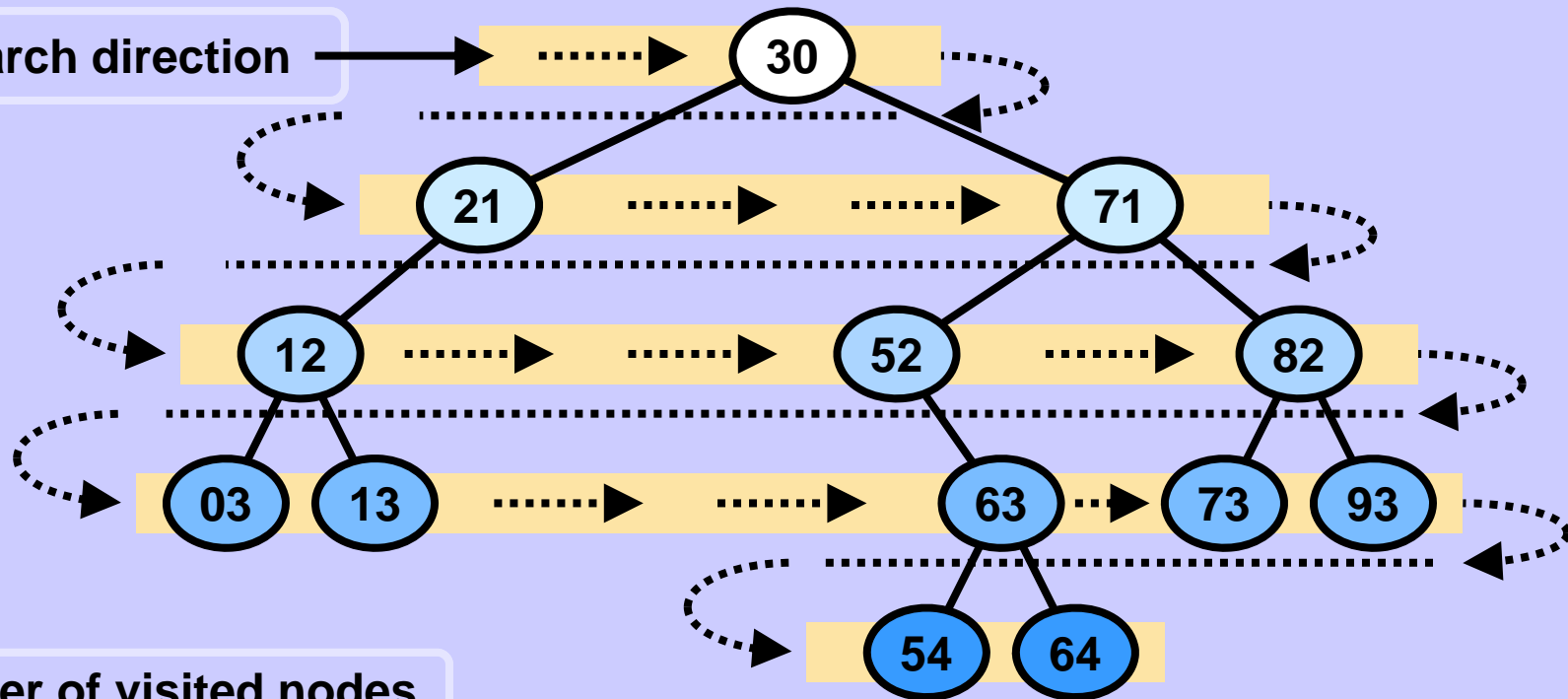
    void Enqueue(Node node) {
        if ((tail+1 == front) ||
            (tail-front == size-1))
            ... // queue full, fix it

        q[tail++] = node;
        if (tail==size) tail=0;
    }

    Node Dequeue() {
        Node n = q[front++];
        if (front==size) front=0;
        return n;
    }
} // end of Queue
  
```

## Breadth-first search (BFS) of a tree

Search direction



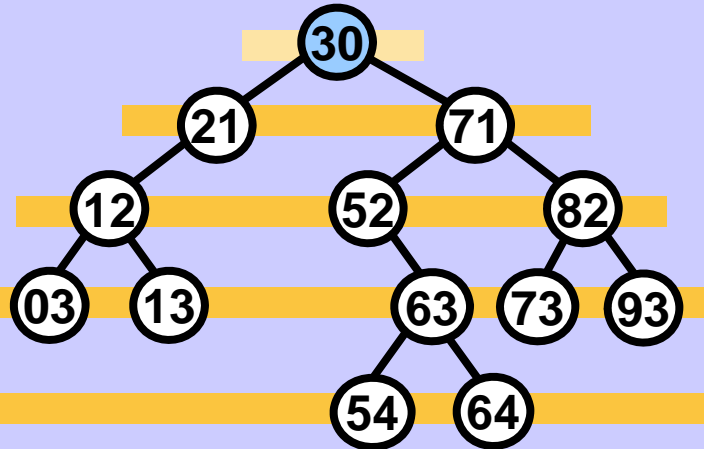
Order of visited nodes

30 21 71 12 52 82 03 13 63 73 93 54 64

Nor the tree structure nor the recursion support this approach directly.

## Breadth-first search (BFS) of a tree

### Initialization



Output

Create an empty queue.



Enqueue the tree root.



Front

Tail

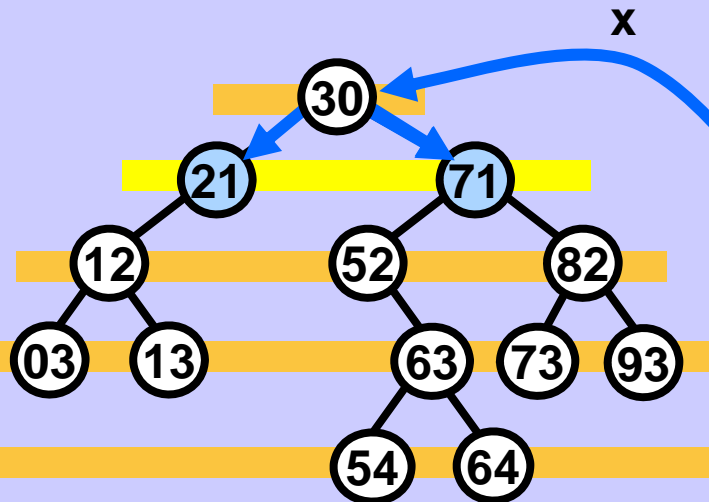
### Main loop

While the queue is not empty do:

1. Remove the first element from the queue and process it.
2. Enqueue the children of removed element.



# Breadth-first search (BFS) of a tree

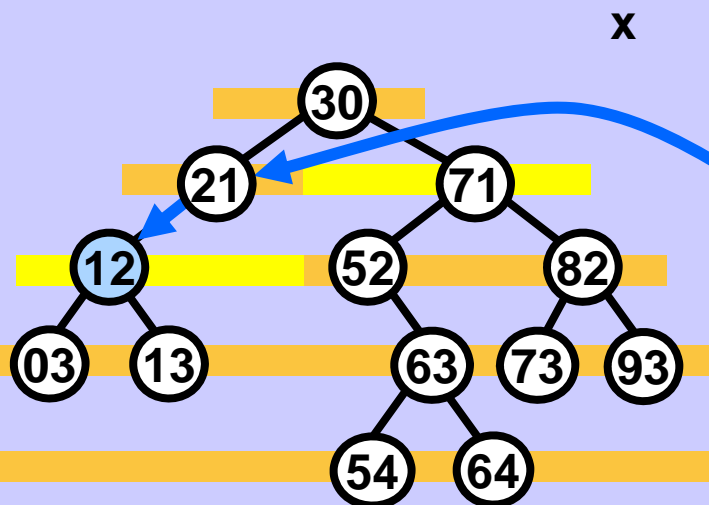
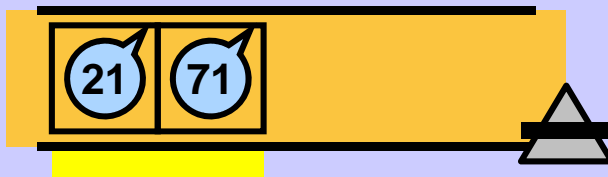


Output 30

1. `x = Dequeue(), print (x.key).`



2. `Enqueue(x.left), Enqueue(x.right). *`

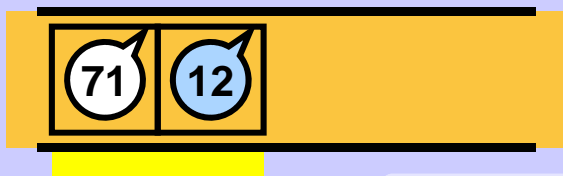


Output 30 21

1. `x = Dequeue(), print (x.key).`

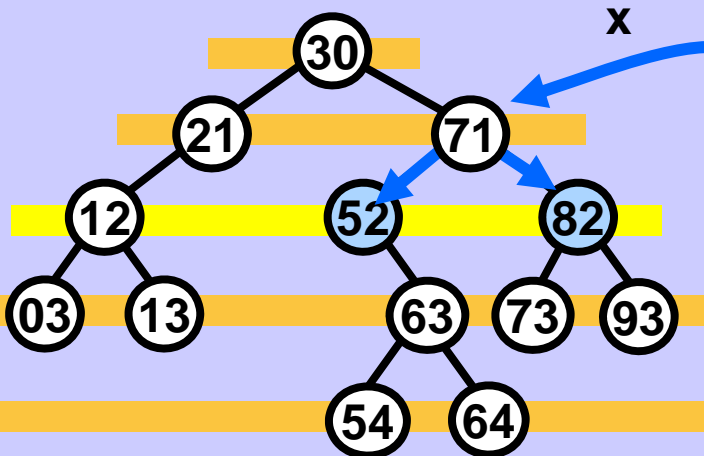


2. `Enqueue(x.left), Enqueue(x.right). *`



\*) if exists

# Breadth-first search (BFS) of a tree

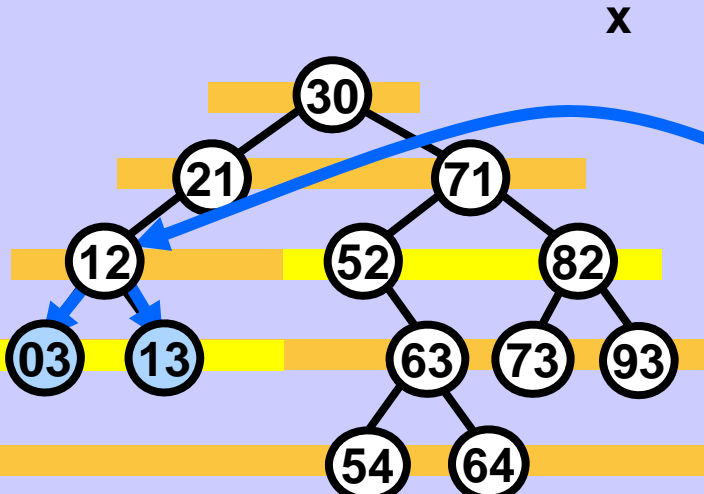
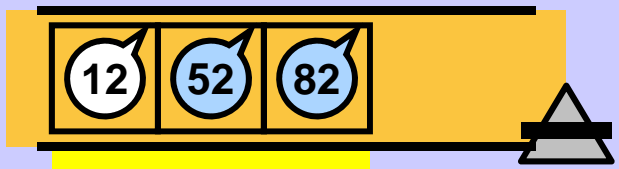


Output 30 21 71

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$



2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



Output 30 21 71 12

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

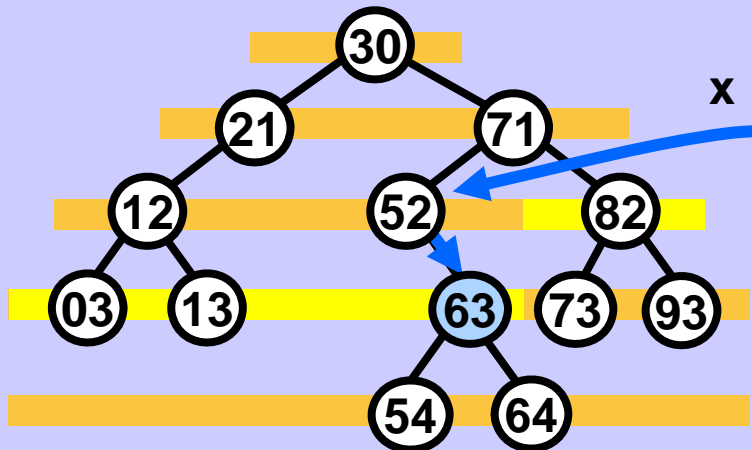


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



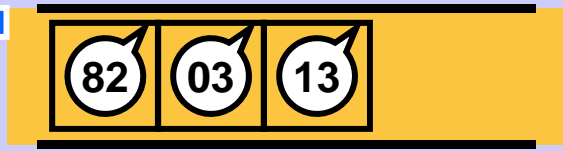
\*) if exists

# Breadth-first search (BFS) of a tree

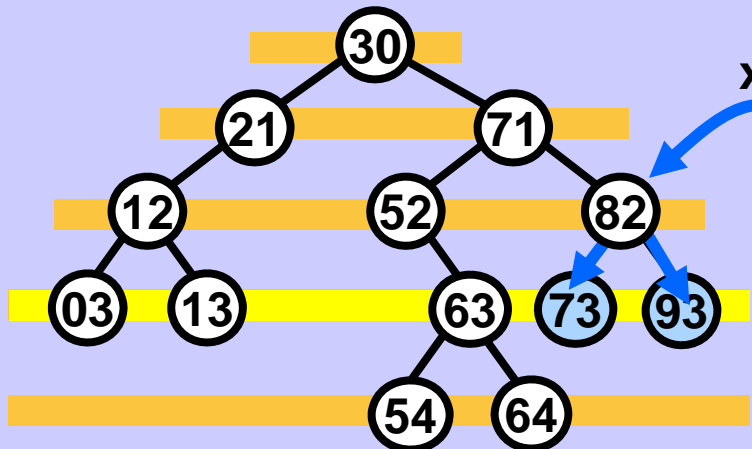


Output 30 21 71 12 52

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

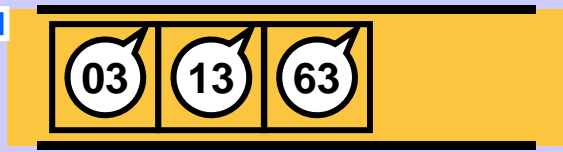


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$

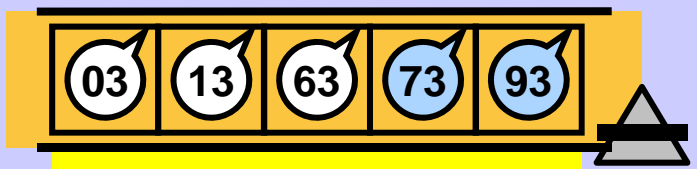


Output 30 21 71 12 52 82

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

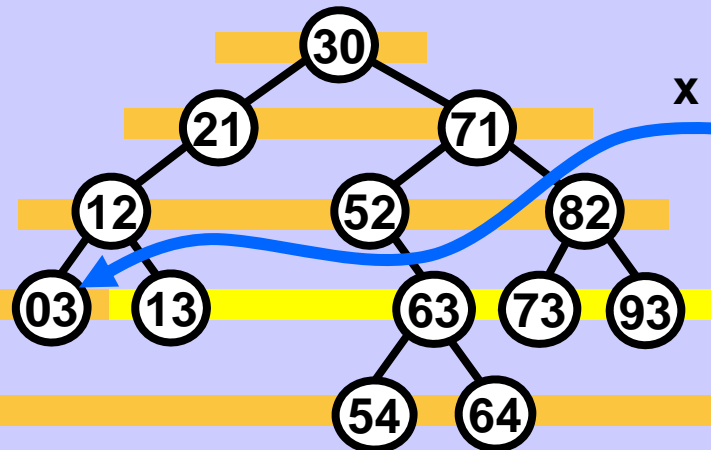


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



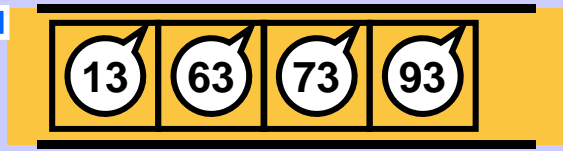
\*) if exists

# Breadth-first search (BFS) of a tree

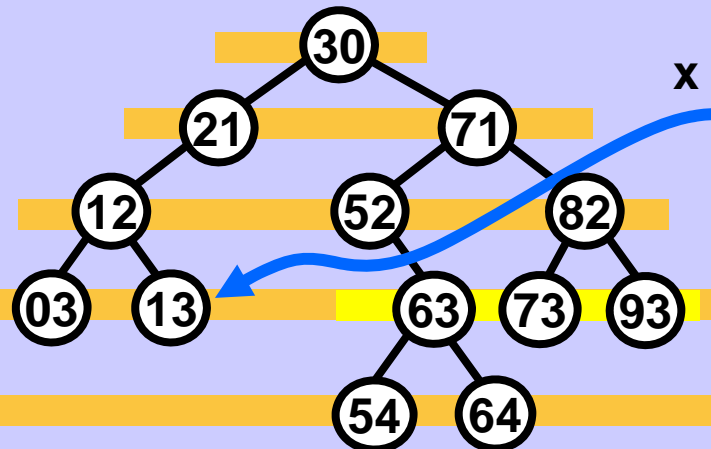


Output 30 21 71 12 52 82 03

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

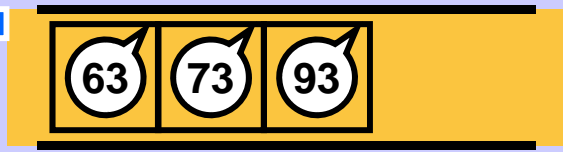


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



Output 30 21 71 12 52 82 03 13

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

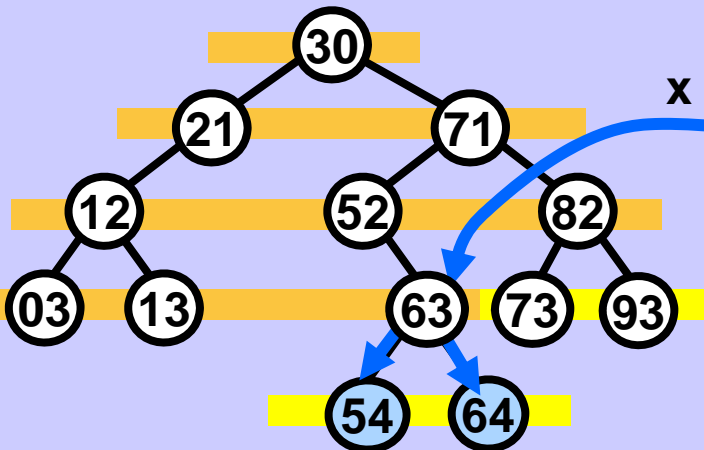


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



\*) if exists

# Breadth-first search (BFS) of a tree

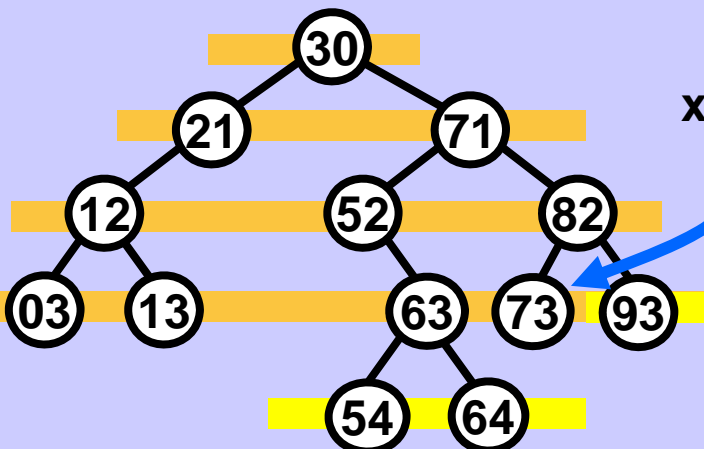


Output 30 21 71 12 52 82 03 13 63

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

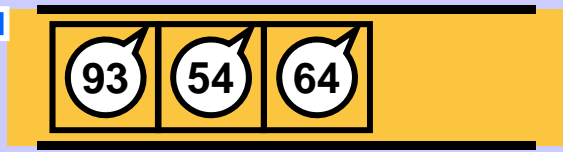


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



Output 30 21 71 12 52 82 03 13 63 73

1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$

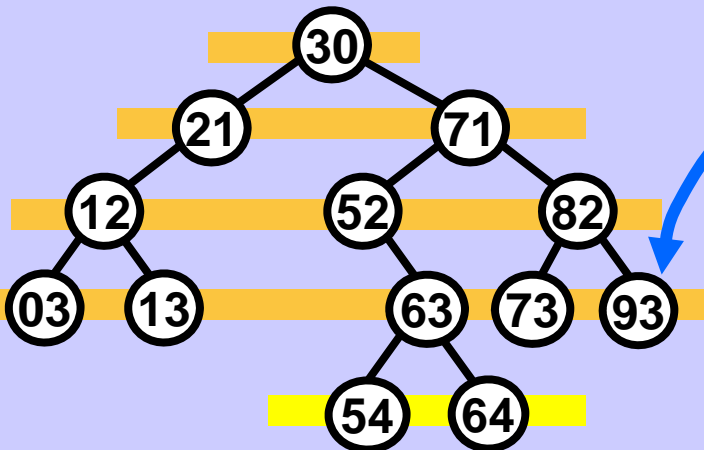


2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



\*) if exists

# Breadth-first search (BFS) of a tree



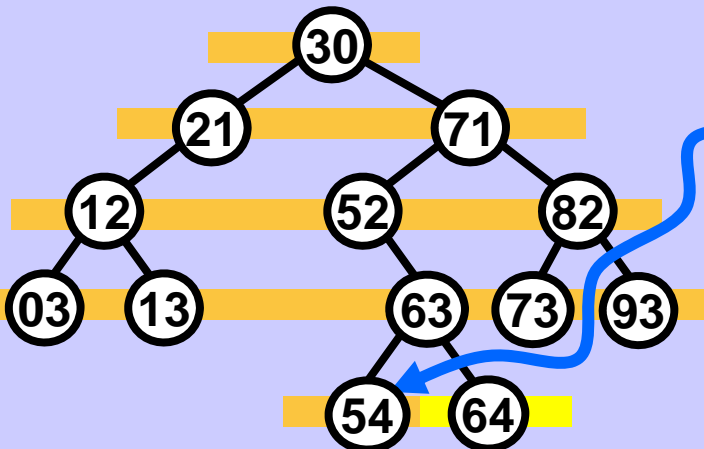
1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$



2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



Output 30 21 71 12 52 82 03 13 63 73 93



1.  $x = \text{Dequeue}(), \text{print}(x.\text{key}).$



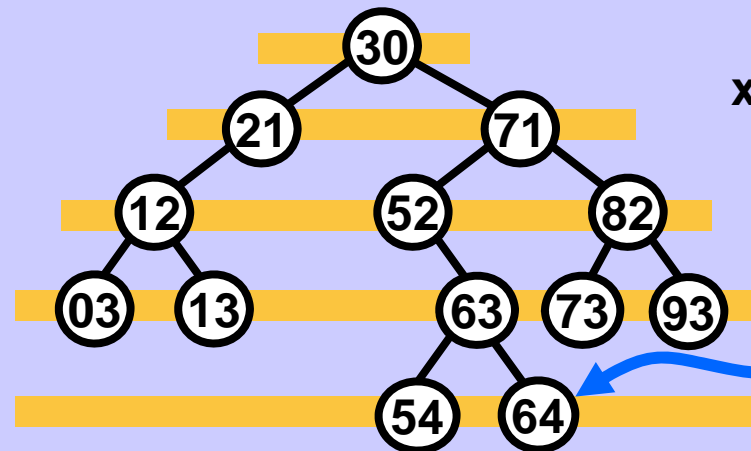
2.  $\text{Enqueue}(x.\text{left}), \text{Enqueue}(x.\text{right}). *$



Output 30 21 71 12 52 82 03 13 63 73 93 54

\*) if exists

## Breadth-first search (BFS) of a tree



1. `x = Dequeue(), print (x.key).`

2. `Enqueue(x.left), Enqueue(x.right). *`

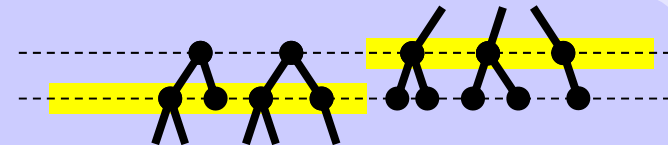
Output

30 21 71 12 52 82 03 13 63 73 93 54 64

\*) if exists.

The queue is empty,  
BFS is complete.

An unempty **queue** always contains exactly  
-- some (or all) nodes of one level and  
-- all children of those nodes of this level which have already left the queue.



Sometimes the queue contains just nodes of one level. See above:



## Breadth-first search (BFS) of a tree

```
void binaryTreeBFS (Node node) {  
    if (node == null) return;  
    Queue q = new Queue();           // init  
    q.Enqueue(node);                 // root into queue  
    while (!q.Empty()) {  
        node = q.Dequeue();  
        print(node.key);             // process node  
        if (node.left != null) q.Enqueue(node.left);  
        if (node.right != null) q.Enqueue(node.right);  
    }  
}
```