# Robot Motion Planning I. 

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Lecture 9
A4M36PAH - Planning and Games

## Robot Motio Planning I.

Introduction

Notation and Terminology

Sampling-based Motion Planning

Randomized Sampling-Based Methods

Optimal Motion Planners

## Literature

Robot Motion Planning, Jean-Claude Latombe, Kluwer Academic Publishers, Boston, MA, 1991.


Principles of Robot Motion: Theory, Algorithms, and Implementations, H. Choset, K. M. Lynch, S. Hutchinson, G. Kantor, W. Burgard, L. E. Kavraki and S. Thrun, MIT Press, Boston, 2005.

http://www.cs.cmu.edu/~biorobotics/book

Planning Algorithms, Steven M. LaValle, Cambridge University Press, May 29, 2006.
http://planning.cs.uiuc.edu

Robot Motion Planning and Control, Jean-Paul Laumond, Lectures Notes in Control and Information Sciences, 2009.
http://homepages.laas.fr/jpl/book.html

$\square$


## Robot Motion Planning - Motivational problem

■ How to transform high-level task specification (provided by humans) into a low-level description suitable for controlling the actuators?

To develop algorithms for such a transformation.
The motion planning algorithms provide transformations how to move a robot (object) considering all operational constraints.



## Piano Mover's Problem

## A classical motion planning problem

Having a CAD model of the piano, model of the environment, the problem is how to move the piano from one place to another without hitting anything.


Basic motion planning algorithms are focused primarily on rotations and translations.

- We need notion of model representations and formal definition of the problem.
- Moreover, we also need a context about the problem and realistic assumptions.

The plans have to be admissible and feasible.

## Robotic Planning Context

Mission Planning
Tasks and Actions Plans
symbol level


## Real Mobile Robots

In a real deployment, the problem is a more complex.

- The world is changing

■ Robots update the knowledge about the environment
localization, mapping and navigation
■ New decisions have to made

- A feedback from the environment Motion planning is a part of the mission replanning loop.


An example of robotic mission:
Multi-robot exploration of unknown environment
How to deal with real-world complexity?
Relaxing constraints and considering realistic assumptions.

## Notation

■ $\mathcal{W}$ - World model describes the robot workspace and its boundary determines the obstacles $\mathcal{O}_{i}$.

$$
2 D \text { world, } \mathcal{W}=\mathbb{R}^{2}
$$

■ A Robot is defined by its geometry, parameters (kinematics) and it is controllable by the motion plan.

- $\mathcal{C}$ - Configuration space ( $\mathcal{C}$-space)

A concept to describe possible configurations of the robot. The robot's configuration completely specify the robot location in $\mathcal{W}$ including specification of all degrees of freedom.
E.g., a robot with rigid body in a plane $\mathcal{C}=\{x, y, \varphi\}=\mathbb{R}^{2} \times S^{1}$.
$\square$ Let $\mathcal{A}$ be a subset of $\mathcal{W}$ occupied by the robot, $\mathcal{A}=\mathcal{A}(q)$.

- A subset of $\mathcal{C}$ occupied by obstacles is

$$
\mathcal{C}_{o b s}=\left\{q \in \mathcal{C}: \mathcal{A}(q) \cap \mathcal{O}_{i}, \forall i\right\}
$$

- Collision-free configurations are

$$
\mathcal{C}_{\text {free }}=\mathcal{C} \backslash \mathcal{C}_{\text {obs }}
$$

## Path / Motion Planning Problem

■ Path is a continuous mapping in $\mathcal{C}$-space such that

$$
\pi:[0,1] \rightarrow \mathcal{C}_{\text {free }}, \text { with } \pi(0)=q_{0}, \text { and } \pi(1)=q_{f}
$$

Only geometric considerations

- Trajectory is a path with explicate parametrization of time, e.g., accompanied by a description of the motion laws $(\gamma:[0,1] \rightarrow \mathcal{U}$, where $\mathcal{U}$ is robot's action space).

It includes dynamics.

$$
\left[T_{0}, T_{f}\right] \ni t \rightsquigarrow \tau \in[0,1]: q(t)=\pi(\tau) \in \mathcal{C}_{\text {free }}
$$

The planning problem is determination of the function $\pi(\cdot)$.
Additional requirements can be given:

- Smoothness of the path
- Kinodynamic constraints
E.g., considering friction forces

■ Optimality criterion shortest vs fastest (length vs curvature)

## Planning in $\mathcal{C}$-space

Robot motion planning robot for a disk robot with a radius $\rho$.


Motion planning problem in geometrical representation of $\mathcal{W}$


Motion planning problem in $\mathcal{C}$-space representation
$\mathcal{C}$-space has been obtained by enlarging obstacles by the disk $\mathcal{A}$ with the radius $\rho$.

By applying Minkowski sum: $\mathcal{O} \oplus \mathcal{A}=\{x+y \mid x \in \mathcal{O}, y \in \mathcal{A}\}$.

## Example of $\mathcal{C}_{\text {obs }}$ for a Robot with Rotation



A simple 2D obstacle $\rightarrow$ has a complicated $\mathcal{C}_{\text {obs }}$

- Deterministic algorithms exist

Requires exponential time in $\mathcal{C}$ dimension,
J. Canny, PAMI, 8(2):200-209, 1986

■ Explicit representation of $\mathcal{C}_{\text {free }}$ is impractical to compute.

## Representation of $\mathcal{C}$-space

How to deal with continuous representation of $\mathcal{C}$-space?

## Continuous Representation of $\mathcal{C}$-space

## Discretization

processing critical geometric events, (random) sampling roadmaps, cell decomposition, potential field

## Graph Search Techniques BFS, Gradient Search, A*

## Planning Methods - Overview

(selected approaches)

- Roadmap based methods

Create a connectivity graph of the free space.
■ Visibility graph
(complete but impractical)

- Cell decomposition
- Voronoi diagram

■ Discretization into a grid-based (or lattice-based) representation (resolution complete)

- Potential field methods
(complete only for a "navigation function", which is hard to compute in general)

Classic path planning algorithms

■ Randomized sampling-based methods
■ Creates a roadmap from connected random samples in $\mathcal{C}_{\text {free }}$

- Probabilistic roadmaps
samples are drawn from some distribution
■ Very successful in practice


## Visibility Graph

1. Compute visibility graph
2. Find the shortest path
E.g., by Dijkstra's algorithm


Found shortest path

Constructions of the visibility graph:
■ Naïve - all segments between $n$ vertices of the map $O\left(n^{3}\right)$

- Using rotation trees for a set of segments - $O\left(n^{2}\right)$
M. H. Overmars and E. Welzl, 1988


## Voronoi Diagram

1. Roadmap is Voronoi diagram that maximizes clearance from the obstacles
2. Start and goal positions are connected to the graph
3. Path is found using a graph search algorithm


Voronoi diagram


Path in graph


Found path

## Visibility Graph vs Voronoi Diagram

Visibility graph
■ Shortest path, but it is close to obstacles. We have to consider safety of the path.

An error in plan execution can
lead to a collision.

- Complicated in higher dimensions

Voronoi diagram

- It maximize clearance, which can provide conservative paths
- Small changes in obstacles can lead to large changes in the diagram
- Complicated in higher dimensions


A combination is called Visibility-Voronoi - R. Wein, J. P. van den Berg, D. Halperin, 2004

For higher dimensions we need other roadmaps.

## Cell Decomposition

1. Decompose free space into parts.

Any two points in a convex region can be directly connected by a segment.
2. Create an adjacency graph representing the connectivity of the free space.
3. Find a path in the graph.

## Trapezoidal decomposition



Centroids represent cells


Connect adjacency cells


Find path in the adjacency graph

Other decomposition (e.g., triangulation) are possible.

## Artificial Potential Field Method

- The idea is to create a function $f$ that will provide a direction towards the goal for any configuration of the robot.
■ Such a function is called navigation function and $-\nabla f(q)$ points to the goal.
- Create a potential field that will attract robot towards the goal $q_{f}$ while obstacles will generate repulsive potential repelling the robot away from the obstacles.

The navigation function is a sum of potentials.


Such a potential function can have several local minima.

## Avoiding Local Minima in Artificial Potential Field

■ Consider harmonic functions that have only one extremum

$$
\nabla^{2} f(q)=0
$$

■ Finite element method
Dirichlet and Neumann boundary conditions

J. Mačák, Master thesis, CTU, 2009

## Sampling-based Motion Planning

- Avoids explicit representation of the obstacles in $\mathcal{C}$-space
- A "black-box" function is used to evaluate a configuration $q$ is a collision free
(E.g., based on geometrical models and testing collisions of the models)
- It creates a discrete representation of $\mathcal{C}_{\text {free }}$

■ Configurations in $\mathcal{C}_{\text {free }}$ are sampled randomly and connected to a roadmap (probabilistic roadmap)

- Rather than full completeness they provides probabilistic completeness or resolution completeness

Probabilistic complete algorithms: with increasing number of samples an admissible solution would be found (if exists)

## Probabilistic Roadmaps

A discrete representation of the continuous $\mathcal{C}$-space generated by randomly sampled configurations in $\mathcal{C}_{\text {free }}$ that are connected into a graph.

■ Nodes of the graph represent admissible configuration of the robot.

■ Edges represent a feasible path (trajectory) between the particular configurations.

Probabilistic complete algorithms: with increasing number of samples an admissible solution would be found (if exists)


Having the graph, the final path (trajectory) is found by a graph search technique.

## Probabilistic Roadmap Strategies

## Multi-Query

■ Generate a single roadmap that is then used for planning queries several times.
■ An representative technique is Probabilistic RoadMap (PRM)

俥
Probabilistic Roadmaps for Path Planning in High Dimensional Configuration Spaces
Lydia E. Kavraki and Petr Svestka and Jean-Claude Latombe and Mark H. Overmars,
IEEE Transactions on Robotics and Automation, 12(4):566-580, 1996.

## Single-Query

■ For each planning problem constructs a new roadmap to characterize the subspace of $\mathcal{C}$-space that is relevant to the problem.

- Rapidly-exploring Random Tree - RRT
- Expansive-Space Tree - EST Hsu et al., 1997
■ Sampling-based Roadmap of Trees - SRT
(combination of multiple-query and single-query approaches)
Plaku et al., 2005


## Multi-Query Strategy

Build a roadmap (graph) representing the environment

1. Learning phase
1.1 Sample $n$ points in $\mathcal{C}_{\text {free }}$
1.2 Connect the random configurations using a local planner
2. Query phase
2.1 Connect start and goal configurations with the PRM
E.g., using a local planner
2.2 Use the graph search to find the path

Probabilistic Roadmaps for Path Planning in High Dimensional Configuration Spaces
Lydia E. Kavraki and Petr Svestka and Jean-Claude Latombe and Mark H.
Overmars,
IEEE Transactions on Robotics and Automation, 12(4):566-580, 1996.
First planner that demonstrates ability to solve general planning problems in more than 4-5 dimensions.

## PRM Construction

\#1 Given problem domain

\#4 Connected roadmap

\#2 Random configuration

\#5 Query configurations

\#3 Connecting samples

\#6 Final found path


## Practical PRM

■ Incremental construction
■ Connect nodes in a radius $\rho$

- Local planner tests collisions up to selected resolution $\delta$
- Path can be found by Dijkstra's algorithm


What are the properties of the PRM algorithm?

We need a couple of more formalism.

## Path Planning Problem Formulation

- Path planning problem is defined by a triplet

$$
\mathcal{P}=\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right),
$$

- $\mathcal{C}_{\text {free }}=\operatorname{cl}\left(\mathcal{C} \backslash \mathcal{C}_{\text {obs }}\right), \mathcal{C}=(0,1)^{d}$, for $d \in \mathbb{N}, d \geq 2$
- $q_{\text {init }} \in \mathcal{C}_{\text {free }}$ is the initial configuration (condition)
- $\mathcal{G}_{\text {goal }}$ is the goal region defined as an open subspace of $\mathcal{C}_{\text {free }}$
$■$ Function $\pi:[0,1] \rightarrow \mathbb{R}^{d}$ of bounded variation is called :
- path if it is continuous;
- collision-free path if it is path and $\pi(\tau) \in \mathcal{C}_{\text {free }}$ for $\tau \in[0,1]$;
- feasible if it is collision-free path, and $\pi(0)=q_{\text {init }}$ and $\pi(1) \in \mathrm{cl}\left(\mathcal{Q}_{\text {goal }}\right)$.
- A function $\pi$ with the total variation $\operatorname{TV}(\pi)<\infty$ is said to have bounded variation, where $\operatorname{TV}(\pi)$ is the total variation

$$
\operatorname{TV}(\pi)=\sup _{\left\{n \in \mathbb{N}, 0=\tau_{0}<\tau_{\mathbf{1}}<\ldots<\tau_{n}=s\right\}} \sum_{i=1}^{n}\left|\pi\left(\tau_{i}\right)-\pi\left(\tau_{i-1}\right)\right|
$$

- The total variation $\operatorname{TV}(\pi)$ is de facto a path length.


## Path Planning Problem

■ Feasible path planning:
For a path planning problem $\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$

- Find a feasible path $\pi:[0,1] \rightarrow \mathcal{C}_{\text {free }}$ such that $\pi(0)=q_{\text {init }}$ and $\pi(1) \in \mathrm{cl}\left(\mathcal{Q}_{\text {goal }}\right)$, if such path exists.
- Report failure if no such path exists.
- Optimal path planning:

The optimality problem ask for a feasible path with the minimum cost.
For $\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$ and a cost function $c: \Sigma \rightarrow \mathbb{R}_{\geq 0}$
■ Find a feasible path $\pi^{*}$ such that $c\left(\pi^{*}\right)=\min \{c(\pi): \pi$ is feasible $\}$.

- Report failure if no such path exists.

The cost function is assumed to be monotonic and bounded, i.e., there exists $k_{c}$ such that $c(\pi) \leq k_{c} \operatorname{TV}(\pi)$.

## Probabilistic Completeness 1/2

First, we need robustly feasible path planning problem
$\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$.
■ $q \in \mathcal{C}_{\text {free }}$ is $\delta$-interior state of $\mathcal{C}_{\text {free }}$ if the closed ball of radius $\delta$ centered at $q$ lies entirely inside $\mathcal{C}_{\text {free }}$.

$\square \delta$-interior of $\mathcal{C}_{\text {free }}$ is $\operatorname{int}_{\delta}\left(\mathcal{C}_{\text {free }}\right)=\left\{q \in \mathcal{C}_{\text {free }} \mid \mathcal{B}_{/, \delta} \subseteq \mathcal{C}_{\text {free }}\right\}$. A collection of all $\delta$-interior states.

- A collision free path $\pi$ has strong $\delta$-clearance, if $\pi$ lies entirely inside $\operatorname{int}_{\delta}\left(\mathcal{C}_{\text {free }}\right)$.
- $\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$ is robustly feasible if a solution exists and it is a feasible path with strong $\delta$-clearance, for $\delta>0$.


## Probabilistic Completeness 2/2

An algorithm $\mathcal{A L G}$ is probabilistically complete if, for any robustly feasible path planning problem $\mathcal{P}=\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$

$$
\lim _{n \rightarrow 0} \operatorname{Pr}(\mathcal{A L G} \text { returns a solution to } \mathcal{P})=1
$$

- It is a "relaxed' notion of completeness
- Applicable only to problems with a robust solution.


We need some space, where random configurations can be sampled

## Asymptotic Optimality 1/4

Asymptotic optimality relies on a notion of weak $\delta$-clearance
Notice, we use strong $\delta$-clearance for probabilistic completeness

- Function $\psi:[0,1] \rightarrow \mathcal{C}_{\text {free }}$ is called homotopy, if $\psi(0)=\pi_{1}$ and $\psi(1)=$ $\pi_{2}$ and $\psi(\tau)$ is collision-free path for all $\tau \in[0,1]$.
- A collision-free path $\pi_{1}$ is homotopic to $\pi_{2}$ if there exists homotopy function $\psi$.

A path homotopic to $\pi$ can be continuously transformed to $\pi$ through $\mathcal{C}_{\text {free }}$.

## Asymptotic Optimality 2/4

- A collision-free path $\pi:[0, s] \rightarrow \mathcal{C}_{\text {free }}$ has weak $\delta$-clearance if there exists a path $\pi^{\prime}$ that has strong $\delta$-clearance and homotopy $\psi$ with $\psi(0)=\pi, \psi(1)=\pi^{\prime}$, and for all $\alpha \in(0,1]$ there exists $\delta_{\alpha}>0$ such that $\psi(\alpha)$ has strong $\delta$-clearance.

Weak $\delta$-clearance does not require points along a path to be at least a distance $\delta$ away from obstacles.


- A path $\pi$ with a weak $\delta$-clearance
- $\pi^{\prime}$ lies in int ${ }_{\delta}\left(\mathcal{C}_{\text {free }}\right)$ and it is the same homotopy class as $\pi$


## Asymptotic Optimality 3/4

■ It is applicable with a robust optimal solution that can be obtained as a limit of robust (non-optimal) solutions.

- A collision-free path $\pi^{*}$ is robustly optimal solution if it has weak $\delta$-clearance and for any sequence of collision free paths $\left\{\pi_{n}\right\}_{n \in \mathbb{N}}$, $\pi_{n} \in \mathcal{C}_{\text {free }}$ such that $\lim _{n \rightarrow \infty} \pi_{n}=\pi^{*}$,

$$
\lim _{n \rightarrow \infty} c\left(\pi_{n}\right)=c\left(\pi^{*}\right) .
$$

There exists a path with strong $\delta$-clearance, and $\pi^{*}$ is homotopic to such path and $\pi^{*}$ is of the lower cost.
■ Weak $\delta$-clearance implies robustly feasible solution problem
(thus, probabilistic completeness)

## Asymptotic Optimality 4/4

An algorithm $\mathcal{A L G}$ is asymptotically optimal if, for any path planning problem $\mathcal{P}=\left(\mathcal{C}_{\text {free }}, q_{\text {init }}, \mathcal{Q}_{\text {goal }}\right)$ and cost function $c$ that admit a robust optimal solution with the finite cost $c^{*}$

$$
\operatorname{Pr}\left(\left\{\lim _{i \rightarrow \infty} Y_{i}^{\mathcal{A L G}}=c^{*}\right\}\right)=1
$$

- $Y_{i}^{\mathcal{A L G}}$ is the extended random variable corresponding to the minimumcost solution included in the graph returned by $\mathcal{A L G}$ at the end of iteration $i$.


## Properties of the PRM Algorithm

- Completeness for the standard PRM has not been provided when it was introduced

■ A simplified version of the PRM (called sPRM) has been mostly studied

■ sPRM is probabilistically complete
What are the differences between PRM and sPRM?

## PRM vs simplified PRM (sPRM)

## PRM

Input: $q_{\text {init }}$, number of samples $n$, radius $\rho$
Output: PRM - $G=(V, E)$

$$
\begin{aligned}
& \bar{V} \leftarrow \emptyset ; E \leftarrow \emptyset \text {; } \\
& \text { for } i=0, \ldots, n \text { do } \\
& q_{\text {rand }} \leftarrow \text { SampleFree; } \\
& U \leftarrow \operatorname{Near}\left(G=(V, E), q_{\text {rand }}, \rho\right) \text {; } \\
& V \leftarrow V \cup\left\{q_{\text {rand }}\right\} \text {; } \\
& \text { foreach } u \in U \text {, with increasing } \\
& \left\|u-q_{r}\right\| \text { do } \\
& \text { if } q_{\text {rand }} \text { and } u \text { are not in the } \\
& \text { same connected component of } \\
& G=(V, E) \text { then } \\
& \text { if CollisionFree }\left(q_{\text {rand }}, u\right) \\
& \text { then } \\
& E \leftarrow E \cup \\
& \left\{\left(q_{\text {rand }}, u\right),\left(u, q_{\text {rand }}\right)\right\} \text {; }
\end{aligned}
$$

## sPRM Algorithm

Input: $q_{\text {init }}$, number of samples $n$, radius $\rho$
Output: PRM - $G=(V, E)$
$\bar{V} \leftarrow\left\{q_{\text {init }}\right\} \cup$
$\left\{\text { SampleFree }_{i}\right\}_{i=1, \ldots, n-1} ; E \leftarrow \emptyset$;
foreach $v \in V$ do

$$
U \leftarrow \operatorname{Near}(G=(V, E), v, \rho) \backslash\{v\} ;
$$

foreach $u \in U$ do
if CollisionFree $(v, u)$ then
$L E \leftarrow E \cup\{(v, u),(u, v)\} ;$
return $G=(V, E)$;
There are several ways for the set $U$ of vertices to connect them

- $k$-nearest neighbors to $v$
- variable connection radius $\rho$ as a function of $n$


## PRM - Properties

■ sPRM (simplified PRM)

- Probabilistically complete and asymptotically optimal

■ Processing complexity $O\left(n^{2}\right)$

- Query complexity $O\left(n^{2}\right)$

■ Space complexity $O\left(n^{2}\right)$
■ Heuristics practically used are usually not probabilistic complete

- $k$-nearest sPRM is not probabilistically complete
- variable radius SPRM is not probabilistically complete

Based on analysis of Karaman and Frazzoli
PRM algorithm:

+ Has very simple implementation
+ Completeness (for sPRM)
- Differential constraints (car-like vehicles) are not straightforward


## Comments about Random Sampling 1/2

■ Different sampling strategies (distributions) may be applied


■ Notice, one of the main issue of the randomized sampling-based approaches is the narrow passage

- Several modifications of sampling based strategies have been proposed in the last decades


## Comments about Random Sampling 2/2

- A solution can be found using only a few samples.

Do you know the Oraculum? (from Alice in Wonderland)
■ Sampling strategies are important

- Near obstacles
- Narrow passages
- Grid-based
- Uniform sampling must be carefully considered.

James J. Kuffner, Effective Sampling and Distance Metrics for 3D Rigid Body Path Planning, ICRA, 2004.


Naïve sampling


Uniform sampling of $\mathrm{SO}(3)$ using Euler angles

## Rapidly Exploring Random Tree (RRT)

Single-Query algorithm
■ It incrementally builds a graph (tree) towards the goal area.
It does not guarantee precise path to the goal configuration.

1. Start with the initial configuration $q_{0}$, which is a root of the constructed graph (tree)
2. Generate a new random configuration $q_{\text {new }}$ in $\mathcal{C}_{\text {free }}$
3. Find the closest node $q_{\text {near }}$ to $q_{\text {new }}$ in the tree
E.g., using KD-tree implementation like ANN or FLANN libraries
4. Extend $q_{\text {near }}$ towards $q_{\text {new }}$

Extend the tree by a small step, but often a direct control $u \in \mathcal{U}$ that will move robot the position closest to $q_{\text {new }}$ is selected (applied for $\delta t$ ).
5. Go to Step 2, until the tree is within a sufficient distance from the goal configuration

Or terminates after dedicated running time.

## RRT Construction

\#1 new random configuration

$q_{\text {new }}$

## \#3 possible actions from $q_{\text {near }}$


\#2 the closest node

\#4 extended tree


## RRT Algorithm

- Motivation is a single query and control-based path finding
- It incrementally builds a graph (tree) towards the goal area.


## RRT Algorithm

Input: $q_{\text {init }}$, number of samples $n$
Output: Roadmap $G=(V, E)$

```
\(\bar{V} \leftarrow\left\{q_{i n i t}\right\} ; E \leftarrow \emptyset\);
for \(i=1, \ldots, n\) do
    \(q_{\text {rand }} \leftarrow\) SampleFree;
    \(q_{\text {nearest }} \leftarrow \operatorname{Nearest}\left(G=(V, E), q_{\text {rand }}\right)\);
    \(q_{\text {new }} \leftarrow \operatorname{Steer}\left(q_{\text {nearest }}, q_{\text {rand }}\right)\);
    if CollisionFree \(\left(q_{\text {nearest }}, q_{\text {new }}\right)\) then
            \(L V \leftarrow V \cup\left\{x_{\text {new }}\right\} ; E \leftarrow E \cup\left\{\left(x_{\text {nearest }}, x_{\text {new }}\right)\right\}\);
return \(G=(V, E)\);
```

Extend the tree by a small step, but often a direct control $u \in \mathcal{U}$ that will move robot to the position closest to $q_{\text {new }}$ is selected (applied for $d t$ ).

Rapidly-exploring random trees: A new tool for path planning S. M. LaValle,

Technical Report 98-11, Computer Science Dept., Iowa State University, 1998

## Properties of RRT Algorithms

■ Rapidly explores the space
$q_{\text {new }}$ will more likely be generated in large not yet covered parts.
■ Allows considering kinodynamic/dynamic constraints (during the expansion).

- Can provide trajectory or a sequence of direct control commands for robot controllers.

■ A collision detection test is usually used as a "black-box".
E.g., RAPID, Bullet libraries.

■ Similarly to PRM, RRT algorithms have poor performance in narrow passage problems.

- RRT algorithms provides feasible paths.

It can be relatively far from optimal solution, e.g., according to the length of the path.
■ Many variants of RRT have been proposed.

## RRT - Examples 1/2



Alpha puzzle benchmark


Bugtrap benchmark


Apply rotations to reach the goal


Courtesy of V. Vonásek and P. Vaněk

## RRT - Examples 2/2

■ Planning for a car-like robot


■ Planning on a 3D surface


- Planning with dynamics
(friction forces)


Courtesy of V. Vonásek and P. Vaněk

## Car-Like Robot

- Configuration

$$
\vec{x}=\left(\begin{array}{l}
x \\
y \\
\phi
\end{array}\right)
$$

position and orientation

- Controls

$$
\vec{u}=\binom{v}{\varphi}
$$

forward velocity, steering angle
■ System equation

$$
\begin{aligned}
\dot{x} & =v \cos \phi \\
\dot{y} & =v \sin \phi \\
\dot{\varphi} & =\frac{v}{L} \tan \varphi
\end{aligned}
$$

ICC


Kinematic constraints $\operatorname{dim}(\overrightarrow{\boldsymbol{u}})<\operatorname{dim}(\vec{x})$
Differential constraints on possible $\dot{q}$ :

$$
\dot{x} \sin (\phi)-\dot{y} \cos (\phi)=0
$$

## Control-Based Sampling

- Select a configuration $q$ from the tree $T$ of the current configurations
- Pick a control input $\overrightarrow{\boldsymbol{u}}=(v, \varphi)$ and integrate system (motion) equation over a short period

$$
\left(\begin{array}{c}
\Delta x \\
\Delta y \\
\Delta \varphi
\end{array}\right)=\int_{t}^{t+\Delta t}\left(\begin{array}{c}
v \cos \phi \\
v \sin \phi \\
\frac{v}{L} \tan \varphi
\end{array}\right) d t
$$



■ If the motion is collision-free, add the endpoint to the tree
E.g., considering $k$ configurations for $k \delta t=d t$.

## RRT and Quality of Solution

- RRT provides a feasible solution without quality guarantee

Despite of that, it is successfully used in many practical applications
■ In 2011, a systematical study of the asymptotic behaviour of randomized sampling-based planners has been published

It shows, that in some cases, they converge to a nonoptimal value with a probability 1.
Sampling-based algorithms for optimal motion planning Sertac Karaman, Emilio Frazzoli International Journal of Robotic Research, 30(7):846-894, 2011.

http://sertac.scripts.mit.edu/rrtstar

## RRT and Quality of Solution $1 / 2$

- Let $Y_{i}^{R R T}$ be the cost of the best path in the RRT at the end of iteration $i$.
- $Y_{i}^{R R T}$ converges to a random variable

$$
\lim _{i \rightarrow \infty} Y_{i}^{R R T}=Y_{\infty}^{R R T} .
$$

- The random variable $Y_{\infty}^{R R T}$ is sampled from a distribution with zero mass at the optimum, and

$$
\operatorname{Pr}\left[Y_{\infty}^{R R T}>c^{*}\right]=1 .
$$

Karaman and Frazzoli, 2011

- The best path in the RRT converges to a sub-optimal solution almost surely.


## RRT and Quality of Solution 2/2

■ RRT does not satify a necessary condition for the asymptotic optimality

■ For $0<R<\inf _{q \in \mathcal{Q}_{\text {goal }}}\left\|q-q_{\text {init }}\right\|$, the event $\left\{\lim _{n \rightarrow \infty} Y_{n}^{R T T}=c^{*}\right\}$ occurs only if the $k$-th branch of the RRT contains vertices outside the $R$-ball centered at $q_{\text {init }}$ for infinitely many $k$.

See Appendix B in Karaman\&Frazzoli, 2011

- It is required the root node will have infinitely many subtrees that extend at least a distance $\epsilon$ away from $q_{\text {init }}$

The sub-optimality is caused by disallowing new better paths to be discovered.

## Rapidly-exploring Random Graph (RRG)

## RRG Algorithm

Input: $q_{\text {init }}$, number of samples $n$
Output: $G=(V, E)$

```
\(V \leftarrow \emptyset ; E \leftarrow \emptyset\);
for \(i=0, \ldots, n\) do
    \(q_{\text {rand }} \leftarrow\) SampleFree;
    \(q_{\text {nearest }} \leftarrow \operatorname{Nearest}\left(G=(V, E), q_{\text {rand }}\right)\);
    \(q_{\text {new }} \leftarrow \operatorname{Steer}\left(q_{\text {nearest }}, q_{\text {rand }}\right)\);
    if CollisionFree \(\left(q_{\text {nearest }}, q_{\text {new }}\right)\) then
            \(\mathcal{Q}_{\text {near }} \leftarrow \operatorname{Near}(G=\)
                \(\left.(V, E), q_{\text {new }}, \min \left\{\gamma_{R R G}(\log (\operatorname{card}(V)) / \operatorname{card}(V))^{1 / d}, \eta\right\}\right)\);
                \(V \leftarrow V \cup\left\{q_{\text {new }}\right\} ; E \leftarrow E \cup\left\{\left(q_{\text {nearest }}, q_{\text {new }}\right),\left(q_{\text {new }}, q_{\text {nearest }}\right)\right\}\);
                foreach \(q_{\text {near }} \in \mathcal{Q}_{\text {near }}\) do
                if CollisionFree \(\left(q_{\text {near }}, q_{\text {new }}\right)\) then
                    \(E \leftarrow E \cup\left\{\left(q_{\text {rand }}, u\right),\left(u, q_{r a n d}\right)\right\} ;\)
return \(G=(V, E)\);
```

Proposed by Karaman and Frazzoli (2011). Theoretical results are related to properties of Random Geometric Graphs (RGG) introduced by Gilbert (1961) and further studied by Penrose (1999).

## RRG Expansions

■ At each iteration, RRG tries to connect new sample to the all vertices in the $r_{n}$ ball centered at it.

- The ball of radius

$$
r(\operatorname{card}(V))=\min \left\{\gamma_{R R G}\left(\frac{\log (\operatorname{card}(V))}{\operatorname{card}(V)}\right)^{1 / d}, \eta\right\}
$$

where

- $\eta$ is the constant of the local steering function
- $\gamma_{R R G}>\gamma_{R R G}^{*}=2(1+1 / d)^{1 / d}\left(\mu\left(\mathcal{C}_{\text {free }}\right) / \xi_{d}\right)^{1 / d}$
- $d$-dimension of the space;
- $\mu\left(\mathcal{C}_{\text {free }}\right)$ - Lebesgue measure of the obstacle-free space;
- $\xi_{d}$ - volume of the unit ball in $d$-dimensional Euclidean space.
- The connection radius decreases with $n$
- The rate of decay $\approx$ the average number of connections attempted is proportional to $\log (n)$


## RRG Properties

■ Probabilistically complete

- Asymptotically optimal

■ Complexity is $O(\log n)$

- Computational efficiency and optimality
- Attempt connection to $\Theta(\log n)$ nodes at each iteration;
- Reduce volume of the "connection" ball as $\log (n) / n$;
- Increase the number of connections as $\log (n)$.


## Other Variants of the Optimal Motion Planning

■ PRM* - it follows standard PRM algorithm where connections are attempted between roadmap vertices that are within connection radius $r$ as a function of $n$

$$
r(n)=\gamma_{P R M}(\log (n) / n)^{1 / d}
$$

■ RRT* - a modification of the RRG, where cycles are avoided A tree version of the $R R G$

- A tree roadmap allows to consider non-holonomic dynamics and kinodynamic constraints.
■ It is basically RRG with "rerouting" the tree when a better path is discovered.


## Example of Solution 1/2



RRT, $n=250$


RRT*, $n=250$


RRT, $n=500$


RRT*, $n=500$


RRT, $n=2500$


RRT*, $n=2500$


RRT, $n=10000$


RRT*, $n=10000$

Karaman \& Frazzoli, 2011

## Example of Solution 2/2



RRT, $n=20000$


RRT*, $n=20000$

## Overview of Randomized Sampling-based Algorithms

## Algorithm

## Probabilistic Asymptotic

Completeness Optimality

| sPRM | $\checkmark$ | $x$ |
| :---: | :---: | :---: |
| k-nearest sPRM | $x$ | $x$ |
| RRT | $\checkmark$ | $x$ |
| RRG | $\checkmark$ | $\checkmark$ |
| PRM* | $\checkmark$ | $\checkmark$ |
| RRT* | $\checkmark$ | $\checkmark$ |

Notice, k-nearest variants of RRG, PRM*, and RRT* are complete and optimal as well.

## Summary

- Introduction to motion planning

■ Overview of sampling-based planning methods

- Basic roadmap methods
- Visibility graph
- Voronoi diagram
- Cell decomposition
- Artificial potential field method

■ Randomized Sampling-based Methods and their properties (PRM, sPRM, RRT)
■ Optimal Motion Planners (RRG, PRM*, RRT*)

