Constraint-Handling in Evolutionary Algorithms

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Substantial part of this material is based on slides for tutorial 'Constraint-Handling Techniques used with Evolutionary Algorithms' presented at GECCO 2011 by Carlos A. Coello Coello and the technical report Carlos A. Coello Coello: A Survey of Constraint Handling Techniques used with Evolutionary Algorithms.

See http://dl.acm.org/citation.cfm?doid=2001858.2002130 and http://citeseer.ist.psu.edu/viewdoc/download?doi=10.1.1.43.9288&rep=rep1&type=pdf



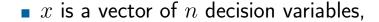
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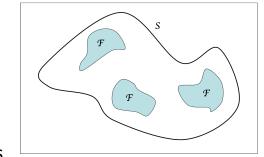
Motivation

The general nonlinear programming problem (NLP) can be formulated as solving the objective function

Minimize
$$f(x)$$
, $x = [x_1, ..., x_n]^T$; subject to $g_i(x) \le 0$, $i = 1, 2, ..., m$; $h_j(x) = 0$, $j = 1, 2, ..., p$;

where





- each x_i , $i=1,\ldots,n$ is bounded by lower and upper limits $x_i^{(L)} \leq x_i \leq x_i^{(U)}$, which define the search space \mathcal{S} ,
- ullet $\mathcal{F} \subseteq \mathcal{S}$ is the feasible region defined by m inequality and p equality constraints.

When solving NLP with EAs, equality constraints are usually transformed into inequality constraints of the form:

$$|h_j(x)| - \varepsilon \le 0$$

where ε is the tolerance allowed.

Standard Evolutionary Algorithm

```
Begin
  t=0
Initialize P(t)
Evaluate P(t)
while (not termination-condition) do
begin
  t=t+1
  Select P(t) from P(t-1)
  Recombine
  Evaluate P(t)
  end
End
```

Evolutionary Algorithms (EAs) have been found successful in solving a wide variety of optimization problems.

However, **EAs are unconstrained search techniques**. Therefore, it is necessary to incorporate constraints into components of the EA (i.e. the fitness function and genetic operators).

Taxonomy of Constraint-Handling Approaches

- Penalty functions.
- Special representations and operators.
- Repair algorithms.
- Multiobjective optimization techniques.

General Form of Fitness Function with Penalty Function

The idea of penalty functions is to transform a constrained optimization problem into unconstrained one by adding certain value to the objective function based on the amount of constraint violation present in the solution:

$$\psi(x) = f(x) + \sum_{i=1}^{m} r_i \times G_i + \sum_{j=1}^{p} r_j \times H_j$$

where $\psi(x)$ is the new objective function referred to as the **fitness function**, G_i and H_j are functions of the constraints violation $(g_i(x))$ and $h_j(x)$, and r_i and r_j are positive constants called **penalty coefficients** or penalty factors.

A common form of G_i :

$$G_i = \max(0, g_i(x))$$

A common form of H_i :

$$H_j = |h_j(x)|$$

or

$$H_j = \max(0, g_j(x)), \text{ for } g_j \equiv |h_j(x)| - \varepsilon \leq 0$$

Types of Penalty Functions used with EAs

Two kinds of penalty functions w.r.t. to the search strategy they imply:

- **Exterior** starting from an infeasible solution the search moves towards a feasible region.
- Interior the penalty term is chosen such that its value will be small at points far away from the constraint boundaries and will tend to infinity as the constraint boundaries are approached. Starting from a feasible solution, the subsequent points will always lie within the feasible region.

Constraint boundaries act as barriers preventing the search to leave the feasible region. This seems nice, but represents a severe drawback indeed.

Four categories of penalty functions based on the way its parameters are being determined:

- Death penalty.
- Static penalty.
- Dynamic penalty.
- Adaptive penalty.

Death Penalty

The **rejection of infeasible individuals** is probably the easiest way to handle constraints and it is also computationally efficient, because when a certain solution violates a constraint, it is rejected and generated again.

 The approach is to iterate, generating a new point at each iteration, until a feasible solution is found.

Thus, no further calculations are necessary to estimate the degree of infeasibility of such a solution.

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Criticism:

- Not advisable, except in the case of problems in which the proportion of feasible region in the whole search space is fairly large.
- No exploitation of the information from infeasible solutions.
- Search may "stagnate" in the presence of very small feasible regions.
- A variation that assigns a zero (or very bad) fitness to infeasible solutions may work better in practice.

Death Penalty: Relaxed

Fitness of an individual is determined using:

$$\begin{array}{rcl} f(x) & \text{if the solution is feasible} \\ fitness_i(x) &= & \\ K - \sum_{i=1}^s (\frac{K}{m}) & \text{otherwise} \end{array}$$

where

- s is the number of constraints satisfied, and
- K is a large constant.

If an individual is infeasible,

- its fitness is always worse than a fitness of any other feasible individual and
- its fitness is the same as the fitness of all the individuals that violate the same number of constraints.

Static Penalty

Approaches in which the **penalty coefficients do not depend** on the current generation number, they remain constant during the entire evolution.

The approach proposed in [Homaifar94] defines levels of violation of the constraints (and penalty coefficients associated to them):

$$fitness(x) = f(x) + \sum_{i=1}^{m} (R_{k,i} \times (max[0, g_i(x)])^2)$$

where $R_{k,i}$ are the penalty coefficients used, m is the total number of constraints, f(x) is the objective function, and $k=1,2,\ldots,l$, where l is the number of levels of violation defined by the user.

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Criticism:

■ The weakness of the method is the high number of parameters: for m constraints and l levels of violation for each, the method requires m(2l+1) parameters in total.

For m=5 and l=4, we need 45 parameters!!!

Clearly, the results are heavily parameter dependent.

 Presented method requires prior knowledge of the degree of constraint violation present in the problem (to define the levels of violation), which might not be easy to obtain in real-world applications.

- Penalty coefficients are difficult to generalize as they are, in general, problem-dependent.
- Anyway, it is not a good idea to keep the same penalty coefficient along the entire evolution.
 The population evolves, so why should the coefficients that bias the search direction be static?

Dynamic Penalty

Penalty functions in which the current generation number is involved in the computation of the corresponding penalty coefficients.

Typically, the penalty coefficients are defined in such a way that **they increase over time** pushing the search towards the feasible region.

The approach from [Joines94] evaluates individuals as follows:

$$fitness(x) = f(x) + (C \times t)^{\alpha} \times SVC(\beta, x)$$

where C, α and β are user-defined constants; recommended values are C=0.5 or 0.05, $\alpha=1$ or 2, and $\beta=1$ or 2.

 $SVC(\beta, x)$ is defined as:

$$SVC(\beta, x) = \sum_{i=1}^{m} G_i^{\beta}(x) + \sum_{j=1}^{p} H_j(x)$$

where $G_i(x)$ and $H_j(x)$ are functions of the constraints violation $(g_i(x))$ and $h_j(x)$.

Step-wise non-stationary penalty function increases the penalty proportionally to the generation number. The goal is to allow the GA to explore more of the search space before confining it to the feasible region.

Dynamic Penalty: Criticism

- It is difficult to derive good dynamic penalty functions in practice.

 The presented approach is sensitive to changes in values of α , β and C and there are no guidelines for choosing proper values for particular problem.
- If a bad penalty coefficient is chosen, the EA may converge to either non-optimal feasible solutions (if the penalty is too high) or to infeasible solutions (if the penalty is too low).
- No constraint normalization technique is used, thus certain constraint may undesirably dominate the whole penalty value.

Annealing Penalties

Variable Fitness GA – the fitness function of an individual is computed using

$$fitness_i(x) = \alpha(M, T) \cdot f(x)$$

where the objective function f(x) is to be maximized and

$$\alpha(M,T) = e^{-M/T}$$

is an attenuation factor based on the principle of simulated annealing that depends on

• temperature, T, a function of the running time, t, that tends to 0 as execution proceeds; It uses the following temperature updating schedule

$$T = O(1/\sqrt{t})$$

• extend of constraint violation, $M \geq 0$.

Remarks: Definition of the starting a final values of M strongly determine overall performance of the algorithm:

• When the GA begins, we want the penalty for constraint violation to be small $(\alpha \approx 1)$, in order to utilize infeasible states as needed;

Toward the end of execution we want α to be zero or nearly zero since infeasible solutions are unacceptable.

• The starting temperature should be of the same order of magnitude as the mean constraint violation \overline{M} .

The final temperature to be on the order of one one-hundredth of the mean constraint violation \overline{M} .

• Alternatively, we can compare the progression of the best observed fitness with the best observed feasible solution during a run. When these two match, it means the penalty imposition has driven the algorithm out of the infeasible region whereupon the temperature is sufficiently low. This is the desirable condition.

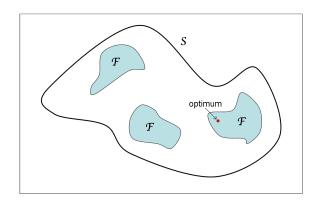
If runs are ending with a difference between the best fitness solution and best feasible solution, then the final temperature should be taken lower.

Adaptive Penalty: Motivation

Let's assume the penalty fitness function of the following form:

$$\psi(x) = f(x) + r_g \times \sum_{i=1}^{m+p} G_i(x)^2$$

Deciding on an optimal (or near-optimal) value r_g is a difficult optimization problem itself.



- If r_g is too small, an infeasible solution may not be penalized enough. Hence, infeasible solutions may be evolved by an EA.
- If r_g is too large, a feasible solution is very likely to be found, but could be of very poor quality. A large r_g discourages the exploration of infeasible regions.

This is inefficient for problems where feasible regions in the whole search space are disjoint and/or the constraint optimum lies close to the boundary of the feasible domain.

Reasonable exploration of infeasible regions may act as bridges connecting feasible regions.

How much exploration of infeasible regions $(r_q =?)$ is reasonable?

- It is problem dependent.
- ullet Even for the same problem, different stages of evol. search may require different r_g values.

Adaptive Penalty

- GA with non-linear penalty function.
- Adaptive Segregational Constraint Handling EA (ASCHEA).
- Stochastic Ranking.

GA with Non-linear Penalty Function

The method introduced in [Hadj-Alouane92] uses adaptive penalty function that takes a feedback from the search process

$$fitness_{i}(x) = f_{i}(x) + r_{g}(t) \sum_{i=1}^{m+p} G_{i}(x)^{2}$$

where $r_g(t)$ is updated every generation according to the following rule:

$$r_g(t+1) \ = \ \begin{array}{ll} (1/\beta_1) \cdot r_g(t), & \text{if case } \#1 \\ \beta_2 \cdot r_g(t), & \text{if case } \#2 \\ r_g(t), & \text{otherwise,} \end{array}$$

where

- case #1 the situation where the best individual in the last k generations was **always** feasible;
- case #2 the situation where the best individual in the last k generations was **never** feasible;
- $\beta_1, \beta_2 > 1$, and $\beta_1 \neq \beta_2$ (to avoid cycling).

Remarks:

■ The penalty component for the generation t+1 is decreased if the feasible region **was** effectively sampled within last k generations;

- The penalty component for the generation t+1 is increased if the feasible region was not effectively sampled within last k generations;
- It tries to avoid having either an all-feasible or an all-infeasible population.
- The problem is how to choose a proper generational gap (k) and the values of β_2 and β_1 .

Adaptive Segregational Constraint Handling EA - ASCHEA

The main idea in ASCHEA [Hamida00] is to maintain both feasible and infeasible individuals in the population, at least when it seems necessary.

It proposes adaptive mechanisms at the population level for constraint optimization based on three main components:

- 1. An adaptive penalty function takes care of the penalty coefficients according to the proportion of feasible individuals in the current population.
- 2. A constraint-driven mate selection used to mate feasible individuals with infeasible ones and thus explore the region around the boundary of the feasible domain.
- 3. A so-called segregational replacement strategy used to favor a given number of feasible individuals in the population.

ASCHEA: Adaptive Penalty

Let's assume the penalty function of the following form:

$$penal(x) = \alpha \sum_{i=1}^{m+p} G_i(x)$$

The penalty coefficient α is adapted based on the desired proportion of feasible solutions in the population τ_{target} and the current proportion at generation t τ_t :

if
$$(\tau_t > \tau_{target})$$
 $\alpha(t+1) = \alpha(t)/fact$
otherwise $\alpha(t+1) = \alpha(t) * fact$

where fact > 1 is a user-defined parameter, a recommended value is around 1.1.

A recommended value of τ_{target} is around 0.6.

ASCHEA: Constraint-driven Mate Selection

Selection mechanism chooses the mate of feasible individuals to be infeasible.

• Only applied when too few (w.r.t au_{target}) feasible individuals are present in the population.

More precisely, to select the mate x_2 for a first parent x_1 :

if $(0 < \tau_t < \tau_{target})$ and $(x_1 \text{ is feasible})$ select x_2 among infeasible solutions only otherwise select x_2 according to fitness

ASCHEA: Segregational Replacement

Deterministic replacement mechanism used in ES-like scheme that should further enhance the chances of survival of feasible individuals.

Assume a population of μ parents, from which λ offspring are generated. Depending on the replacement scheme

- ullet μ individuals out of λ offspring in case of the (μ, λ) -ES, or
- ullet μ individuals out of λ offspring plus μ parents in case of the $(\mu + \lambda)$ -ES

are selected to the new population in the following way:

- 1. First, feasible solutions are selected without replacement based on their fitness, until $\tau_{select} * \mu$ have been selected, or no more feasible solution is available.
- 2. The population is then filled in using standard deterministic selection on the remaining individuals, based on the penalized fitness.

Thus, a user-defined proportion of τ_{select} feasible solutions is considered superior to all infeasible solutions.

A recommended value of τ_{select} is around 0.3.

ASCHEA: Conclusions

- Feasibility elitism as soon as a feasible individual appears, it can only disappear from the population by being replaced by a better feasible solution, even if the penalty coefficient reaches very small value.
- Constraint-driven mate selection accelerates the movement toward the feasible region of infeasible individuals, and helps to explore the region close to the boundary of the feasible domain.
- Adaptability the penalty adaptation as well as the constraint-driven mate selection are activated based on the actual proportion of feasible solutions in the population.

Let's assume the penalty fitness function of the following form:

$$\psi(x) = f(x) + r_g \times \sum_{i=1}^{m+p} G_i(x)^2$$

For a given penalty coefficient $r_g > 0$, let the ranking of λ individuals be

$$\psi(x_1) \le \psi(x_2) \le \dots \le \psi(x_\lambda),\tag{1}$$

For any given adjacent pair i and i+1 in the ranked order

$$f_i + r_g G_i \le f_{i+1} + r_g G_{i+1}$$
 where $f_i = f(x_i)$ and $G_i = G(x_i)$ (2)

we define so called critical penalty coefficient

$$\check{r}_i = (f_{i+1} - f_i)/(G_i - G_{i+1}) \text{ for } G_i \neq G_{i+1}$$
(3)

For given choice of $r_g \ge 0$, there are three different cases which may give rise to inequality (2):

- 1. $f_i \leq f_{i+1}$ and $G_i \geq G_{i+1}$: Objective function plays a dominant role in determining the inequality and the value of r_g should be $0 < r_g < \check{r}_i$.
- 2. $f_i \ge f_{i+1}$ and $G_i < G_{i+1}$: Penalty function plays a dominant role in determining the inequality and the value of r_q should be $0 < \check{r}_i < r_q$.
- 3. $f_i < f_{i+1}$ and $G_i < G_{i+1}$: The comparison is nondominated and $\check{r}_i < 0$. Neither the objective nor the penalty function can determine the inequality by itself.

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Ex.:

$$f_i = 10$$
, $G_i = 7$
 $f_{i+1} = 20$ $G_{i+1} = 5$
 $\check{r}_i = (20 - 10)/(7 - 5) = 5 \implies 0 < r_g < 5$
 $r_g = 4$: $38 \le 40$ the inequality (2) holds
 $r_g = 6$: $52 \nleq 50$ the inequality (2) does not hold

- 2. $f_i \ge f_{i+1}$ and $G_i < G_{i+1}$: Penalty function plays a dominant role in determining the inequality and the value of r_g should be $0 < \check{r}_i < r_g$.
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$$f_i = 20$$
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 $f_{i+1} = 10$ $G_{i+1} = 7$
 $\check{r}_i = (10 - 20)/(5 - 7) = 5 \implies 5 < r_g$
 $r_g = 4$: $40 \nleq 38$ the inequality (2) does not hold
 $r_q = 6$: $50 \leq 52$ the inequality (2) holds

3. $f_i < f_{i+1}$ and $G_i < G_{i+1}$: The comparison is nondominated and $\check{r}_i < 0$. Neither the objective nor the penalty function can determine the inequality by itself.

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Ex.:

$$f_i=10$$
, $G_i=5$
$$f_{i+1}=20 \quad G_{i+1}=7$$

$$\check{r}_i=(20-10)/(5-7)=-5 \quad \Longrightarrow \quad \text{the inequality (2) holds for all } r_g>0$$

The value of r_g has no impact on the inequality (2) when nondominant and feasible individuals are compared.

The value of r_g is critical in the first two cases. It has to be within a certain range $\underline{r}_g < r_g < \overline{r}_g$

- 1. \underline{r}_g is the minimum critical penalty coefficient computed from adjacent individuals ranked only according to the objective function.
- 2. \overline{r}_g is the maximum critical penalty coefficient computed from adjacent individuals ranked only according to the penalty function.

Both bounds are problem dependent and may vary from generation to generation.

There are three categories of r_g values

- 1. $r_g < \underline{r}_q$: Underpenalization All comparisons are based only on the fitness function.
- 2. $r_g > \overline{r}_g$: Overpenalization All comparisons are based only on the penalty function.
- 3. $\underline{r}_g < r_g < \overline{r}_g$: All comparisons are based on a combination of objective and penalty functions. This is what a good constraint-handling technique should do to balance between preserving feasible individuals and rejecting infeasible ones.

But the optimal r_g is hard to determine.

Stochastic Ranking: Realization

This approach [Runarsson00] consists of an evolutionary algorithm that uses a penalty function and a rank-based selection.

- The idea of this approach is that the balance between objective and penalty functions is achieved directly and explicitly.
- It does not require explicit definition of the penalty coefficient r_g value. Instead, it requires a **user-defined parameter** P_f , **which determines the balance** between the objective function and the penalty function.

Rank-based selection

- The population is sorted using an algorithm similar to bubble-sort.
- Parameter P_f specifies a probability of using only the objective function for comparisons of infeasible solutions.
 - If both individuals are feasible then the probability of comparing them according to the objective function is 1. Otherwise, it is P_f .
 - The reminder of the comparisons are realized based on the sum of constraint violation.
 - Recommended range of P_f values is (0.4, 0.5)

The P_f introduces the stochastic component to the ranking process, so that some solutions may get a good rank even if they are infeasible.

Stochastic Ranking: Bubble-sort-like Procedure

```
I_j = j \ \forall \ j \in \{1, \ldots, \lambda\}
     for i = 1 to N do
          for j = 1 to \lambda - 1 do
 3
               sample u \in U(0,1)
 4
               if (\phi(I_j) = \phi(I_{j+1}) = 0) or (u < P_f) then
                  if (f(I_i) > f(I_{i+1})) then
 6
                     swap(I_i, I_{i+1})
                  fi
               else
9
10
                  if (\phi(I_j) > \phi(I_{j+1})) then
                     swap(I_i, I_{i+1})
11
12
                  fi
13
               fi
14
          od
15
          if no swap done break fi
```

©Runarsson, T. P. and Yao, X.: Stochastic Ranking for Constrained Evolutionary Optimization.

Stochastic Ranking: Conclusions

- Does not use any specialized variation operators.
- lacktriangle Does not require a priori knowledge about a problem since it does not use any penalty coefficient r_g in a penalty function.
- The approach is easy to implement.

Special Representation and Operators

Random keys – a representation for representing permutations, it is an efficient method for encoding ordering and scheduling problems.

ullet A random key vector of length l consists of l random values (keys) that are floating numbers between zero and one.

Ex.:
$$\overrightarrow{r}_5 = (0.17, 0.92, 0.63, 0.75, 0.29)$$

• Of importance for the interpretation of the random key vector is the position of the keys.
The positions of the keys in the vector are ordered according to their values in ascending or descending order which gives a permutation of l numbers.

Ex.: Random key vector \overrightarrow{r}_5 can be interpreted as the sequence $\overrightarrow{r}_5^s = 2 \to 4 \to 3 \to 5 \to 1$

Properties of the encoding:

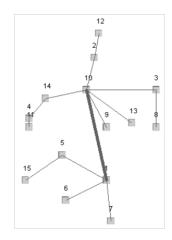
- A valid sequence \overrightarrow{r}_l^s can always be created from the random key vector as long as there are no two keys r_i and r_j with the same values.
- There are many possibilities for the construction of each particular sequence (permutation).
- Lokality of the random keys is high a small change in the genotype (the vector \overrightarrow{r}_l) leads to a small change in the phenotype (the sequence \overrightarrow{r}_l^s).
- When using EAs with random keys, all kinds of standard crossover and mutation operators can be used that always produce only valid solutions (i.e. the interpreted permutations).

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Random Keys for the Network Design Problems

Network Design Problem

- A tree network is defined as a connected graph with n nodes and n-1 links (there are no loops in the tree).
- Between any two nodes there exists exactly one possible path for the flow.
- The goal is to minimize the overall cost for constructing and maintaining the tree network that is calculated by summing-up the cost of all links.



Encoding tree networks with Network Random Keys (NetKeys) [Rothlauf02]

- The real-valued NetKeys are interpreted as the importance of the link. The higher the value of the allele, the higher the probability that the link is used for the tree.
- Every NetKey vector represents a valid network structure.

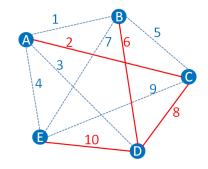
Random Keys for the Network Design Problems

Constructing the tree network from the NetKey vector

- 1. Let i=0, G be an empty graph with n nodes, and \overrightarrow{r}_l^s the sequence with length l=n(n-1)/2 that could be constructed from the NetKey vector \overrightarrow{r}_l . All possible links of G are numbered from 1 to l.
- 2. Let j be the number at the ith position of $\overrightarrow{r}_{l}^{s}$.
- 3. If the insertion of the link with number j in G would not create a cycle, then insert the link with number j in G.
- 4. Stop, if there are n-1 links in G.
- 5. Incrementi and continue with step 2.

Ex.:

position	1	2	3	4	5	6	7	8	9	10
NetKey	0.55	0.73	0.09	0.23	0.40	0.82	0.65	0.85	0.75	0.90
link	A-B	A-C	A-D	A-E	B-C	B-D	В-Е	C-D	C-E	D-E



Approaches based on Evolutionary Multiobjective Optimization

General form of multi-objective optimization problem

$$\begin{array}{ll} \mbox{Minimize/maximize} & f_m(x), & m = 1, 2, ..., M; \\ \mbox{subject to} & g_j(x) \geq 0, & j = 1, 2, ..., J; \\ & h_k(x) = 0, & k = 1, 2, ..., K; \\ & x_i^{(L)} \leq x_i \leq x_i^{(U)}, & i = 1, 2, ..., n. \end{array}$$

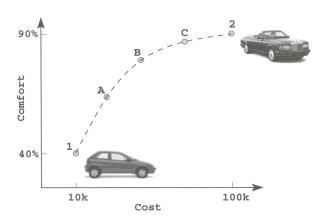
- x is a vector of n decision variables: $x = (x_1, x_2, ..., x_n)^T$;
- g_j , h_k are inequality and equality constraints, respectively.

Conflicting objectives

- A solution that is extreme with respect to one objective requires a compromise in other objectives.
- A sacrifice in one objective is related to the gain in other objective(s).

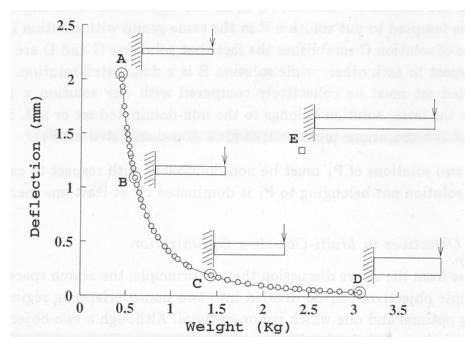
Motivation example: Buying a car

- two extreme hypothetical cars 1 and 2,
- cars with a trade-off between cost and comfort A,
 B, and C.



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Multiobjective Techniques: Using Pareto Schemes



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Pareto dominance: A solution $x^{(1)}$ is said to dominate the other solution $x^{(2)}$, $x^{(1)} \leq x^{(2)}$, if $x^{(1)}$ is no worse than $x^{(2)}$ in all objectives and $x^{(1)}$ is strictly better than $x^{(2)}$ in at least one objective.

Solutions A, B, C, D are non-dominated solutions (Pareto-optimal solutions)

Solution E is *dominated* by C and B (E is non-optimal).

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Approaches based on Evolutionary Multiobjective Optimization

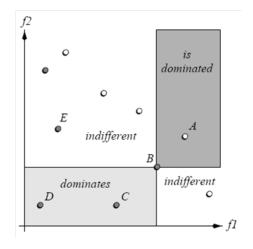
Two ways the NLP is transformed into a multiobjective optimization problem

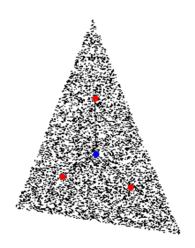
- NLP → Unconstrained Bi-objective Optimization: Transforms the NLP into an unconstrained bi-objective optimization problem with the objectives being (1) the original objective function and (2) the sum of constraint violation.
- NLP → Unconstrained Multiobjective optimization: Transforms the NLP into an unconstrained multiobjective optimization problem where the original objective function and each constraint are treated as separate objectives.

Bi-objective Optimization Techniques

[Zhou03] – uses a ranking procedure based on counting the number of individuals which are dominated for a given solution.

- Ties are solved by the sum of constraint violation.
- Simplex crossover operator used to generate a set of offspring where
 - the solution with the **highest Pareto strength** (the number of population members that are dominated by or equal to the solution with respect to the objective values, divided by the population size plus one) and
 - the solution with the lowest constraint violation
 are both selected to take part in the population.





Bi-objective Optimization Techniques

[Venkatraman05] – approach divided in two phases:

- 1. The population is ranked based only on the sum of constraint violation the goal is to find some feasible solutions.
- 2. Both objectives are taken into account.
 - Nondominated sorting is used to rank the population.
 - Niching scheme based on distance to the nearest neighbors is applied to promote a diversity of the population.

Disadvantage: The way the feasible region is approached is mostly at random because the quality is not considered in the first phase.

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Multiobjective Techniques: Using Non-Pareto Schemes

[Coello00] – MOP approach based on VEGA's idea, where the population is divided into sub-populations, and each sub-population focuses on optimization of one objective.

- m+1 sub-populations.
 - One sub-population handles the objective function of the problem and the individuals are selected based on the unconstrained objective function value.
 - Each of the m remaining sub-populations take one constraint as their fitness function.
 - The aim is that each of the sub-populations tries to reach the feasible region corresponding to one individual constraint.
 - By combining these sub-populations, the approach will reach the feasible region where all of the constraints are satisfied.
- The main drawback is that the number of sub-populations increases linearly with the number of constraints.

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Multiobjective Techniques: NPGA-based Approach

[Coello02] – based on the Niched-Pareto Genetic Algorithm that uses binary tournament selection based on Pareto non-dominance.

- Parameter S_r , which indicates the minimum number of individuals that will be selected through dominance-based tournament selection.
 - The remainder will be selected using a purely probabilistic approach. In other words, $(1 S_r)$ individuals in the population are probabilistically selected.
 - Tournament selection three possible situations when comparing two candidates
 - 1. Both are feasible. In this case, the candidate with a better fitness value wins.
 - 2. One is infeasible, and the other is feasible. The feasible candidates wins, regardless of its fitness function value.
 - 3 Both are infeasible
 - (a) Check both candidates whether they are dominated by ind. from the comparison set.
 - (b) If one is dominated by the comparison set, and the other is not dominated then the non-dominated candidate wins.
 - Otherwise, the candidate with the lowest amount of constraint violation wins, regardless of its fitness function value.
 - Probabilistic selection Each candidate has a probability of 50% of being selected.
- Robust, efficient and effective approach.

Multiobjective Techniques: Conclusions

- The most popular are the MOP approaches.
- The use of diversity mechanisms is found in most approaches.
- The use of explicit local search mechanisms is still scarce.

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