

PKR Lab-09 Solution

Algorithms

Algorithm 1: Multivariate Polynomial Division Algorithm

Input: $f, F = (f_1, \dots, f_s), \geq$ (monomial ordering)
Output: $(q_1, \dots, q_s), r$ such that $f = \sum_{i=1}^s q_i f_i + r$, $\text{LT}_{\geq}(r)$ is not divisible by any of $\text{LT}_{\geq}(f_i)$ or $r = 0$

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1  $q_1 \leftarrow \dots \leftarrow q_s \leftarrow r \leftarrow 0$ 
2  $p \leftarrow f$ 
3 while  $p \neq 0$  do
4    $i \leftarrow 1$ 
5   divisionoccured  $\leftarrow$  False
6   while  $i \leq s$  and not divisionoccured do
7     if  $\text{LT}_{\geq}(f_i)$  divides  $\text{LT}_{\geq}(p)$  then
8        $q_i \leftarrow q_i + \frac{\text{LT}_{\geq}(p)}{\text{LT}_{\geq}(f_i)}$ 
9        $p \leftarrow p - \frac{\text{LT}_{\geq}(p)}{\text{LT}_{\geq}(f_i)} f_i$ 
10      divisionoccured  $\leftarrow$  True
11     else
12        $i \leftarrow i + 1$ 
13   if not divisionoccured then
14      $r \leftarrow r + \text{LT}_{\geq}(p)$ 
15      $p \leftarrow p - \text{LT}_{\geq}(p)$ 
16 return  $(q_1, \dots, q_s), r$ 

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Tasks

Task 1. Consider a polynomial $f = y^5 + 2xy^2 + x$. Order the terms of f (from largest to smallest) according to the lexicographic monomial ordering for the ordering of variables $x > y$.

Solution: For $f = \sum_{\alpha \in \mathbb{Z}_{\geq 0}^n} a_{\alpha} \mathbf{x}^{\alpha}$ we define

$$\text{terms}(f) \stackrel{\text{def}}{=} \{a_{\alpha} \mathbf{x}^{\alpha} \mid a_{\alpha} \neq 0\}.$$

For f given in the task we have

$$\text{terms}(f) = \{y^5, 2xy^2, x\}.$$

We first compare y^5 with $2xy^2$. We have $y^5 = \mathbf{x}^{\alpha} = (x, y)^{(0,5)}$ and $xy^2 = \mathbf{x}^{\beta} = (x, y)^{(1,2)}$. Since the left most nonzero element of $\alpha - \beta = (-1, 3)$ is not positive and $\alpha \neq \beta$, then $\mathbf{x}^{\beta} \geq_{\text{lex}} \mathbf{x}^{\alpha}$, or

$$2xy^2 \geq_{\text{lex}} y^5.$$

Remark. In other words, in order to compare two monomials $\mathbf{x}^{\alpha}$ and $\mathbf{x}^{\beta}$ w.r.t. $\geq_{\text{lex}}$ we compare the exponents in individual variables starting from the largest variable towards the lowest variable: if $\alpha_j > \beta_j$, then $\mathbf{x}^{\alpha} \geq_{\text{lex}} \mathbf{x}^{\beta}$; if $\alpha_j < \beta_j$, then $\mathbf{x}^{\beta} \geq_{\text{lex}} \mathbf{x}^{\alpha}$; if $\alpha_j = \beta_j$, then we move to comparing the exponents α_{j+1} and β_{j+1} for the next variable.

Next we compare x with y^5 and x with $2xy^2$:

$$x \geq_{\text{lex}} y^5, 2xy^2 \geq_{\text{lex}} x.$$

Then

$$f = 2xy^2 + x + y^5$$

with the terms written in the decreasing order. □

Task 2. Consider a monomial ordering $x \geq_{\text{lex}} y$. Divide a polynomial $f = x^2 + xy + x + y^3 + 1$ by an ordered tuple of polynomials

$$F = (f_1, f_2, f_3) = (xy + x + 1, y^2 + 1, xy + 1)$$

using $\geq_{\text{lex}}$. Consider a permuted tuple $F' = (f_3, f_2, f_1)$ of F and divide f by F' using $\geq_{\text{lex}}$. Compare the two remainders.

Solution: We apply Algorithm 1. We represent f as a sum $p + \sum_i q_i \cdot f_i + r$ during every iteration of the **while** $p \neq 0$ **do** loop:

$$\begin{aligned} \underbrace{x^2 + xy + x + y^3 + 1}_f &= \underbrace{x^2 + xy + x + y^3 + 1}_p + \underbrace{0}_{q_1} \cdot \underbrace{(xy + x + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + 1)}_{f_3} + \underbrace{0}_r \\ &= \underbrace{xy + x + y^3 + 1}_p + \underbrace{0}_{q_1} \cdot \underbrace{(xy + x + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + 1)}_{f_3} + \underbrace{x^2}_r \\ &= \underbrace{y^3}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + x + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + 1)}_{f_3} + \underbrace{x^2}_r \\ &= \underbrace{-y}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + x + 1)}_{f_1} + \underbrace{y}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + 1)}_{f_3} + \underbrace{x^2}_r \\ &= \underbrace{0}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + x + 1)}_{f_1} + \underbrace{y}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + 1)}_{f_3} + \underbrace{x^2 - y}_r \end{aligned}$$

We now divide f by F' :

$$\begin{aligned} \underbrace{x^2 + xy + x + y^3 + 1}_f &= \underbrace{x^2 + xy + x + y^3 + 1}_p + \underbrace{0}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{0}_r \\ &= \underbrace{xy + x + y^3 + 1}_p + \underbrace{0}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{x^2}_r \\ &= \underbrace{x + y^3}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{x^2}_r \\ &= \underbrace{y^3}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{0}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{x^2 + x}_r \\ &= \underbrace{-y}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{y}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{x^2 + x}_r \\ &= \underbrace{0}_p + \underbrace{1}_{q_1} \cdot \underbrace{(xy + 1)}_{f_1} + \underbrace{y}_{q_2} \cdot \underbrace{(y^2 + 1)}_{f_2} + \underbrace{0}_{q_3} \cdot \underbrace{(xy + x + 1)}_{f_3} + \underbrace{x^2 + x - y}_r \end{aligned}$$

We notice that the remainder r of the division (as well as the quotient polynomials q_j and the number of iterations) depend on the order of the divisors f_j . $\square$