

Logical reasoning and programming, lab session 1

(September 22, 2025)

1.1 Decide which of the following formulae are tautologies:

- (a) $((p \rightarrow q) \rightarrow q) \rightarrow q$,
- (b) $((p \rightarrow q) \rightarrow p) \rightarrow p$,
- (c) $(p \rightarrow q) \vee (q \rightarrow p)$,
- (d) $((p \rightarrow q) \wedge q) \rightarrow p$,
- (e) $\neg p \rightarrow \neg(p \vee (p \wedge q))$.

Although formulae contain only two variables, try to find a better method than truth tables if possible.

1.2 Decide which of the following claims are true:

- (a) $\neg(\varphi \leftrightarrow \psi) \models \varphi \wedge \psi$,
- (b) $\neg(\varphi \leftrightarrow \psi) \models \varphi \vee \psi$,
- (c) $\neg(\varphi \leftrightarrow \psi) \models \varphi \rightarrow \psi$.

1.3 If $\varphi \rightarrow \psi \in \mathbf{TAUT}$ and $\psi \rightarrow \chi \in \mathbf{TAUT}$, then $\varphi \rightarrow \chi \in \mathbf{TAUT}$. Why? Does the claim still hold if we replace **TAUT** with **SAT** and why?

1.4 Show that $\models \varphi \leftrightarrow \psi$ iff $\varphi \equiv \psi$.

1.5 Are the following two C programs equivalent?

if (!a && !b) h ();	if (a) f ();
else if (!a) g ();	else if (b) g ();
else f ();	else h ();

Prove their equivalence formally, or provide a counterexample. Please, check your solution against these slides.

1.6 Let $\text{Cl}(\Gamma) = \{\varphi : \Gamma \models \varphi\}$. What is $\text{Cl}(\emptyset)$ equivalent to? Let Γ and Δ be sets of formulae. Check whether:

- (a) $\Gamma \subseteq \text{Cl}(\Gamma)$,
- (b) $\text{Cl}(\text{Cl}(\Gamma)) = \text{Cl}(\Gamma)$,
- (c) $\text{Cl}(\Gamma \cup \Delta) = \text{Cl}(\Gamma) \cup \text{Cl}(\Delta)$.

If the equality does not hold in (b) or (c), does at least one of the inclusions hold?

1.7 Recall that a Boolean function of n -variables is a function $f: \{0, 1\}^n \rightarrow \{0, 1\}$. Describe all functions of one variable. How many distinct Boolean functions of n variables exist? Try $n = 0, 1, 2, \dots$ first. Do you know, why are functions NAND ($\uparrow$) and NOR ($\downarrow$) interesting?¹

¹These connectives are also called Sheffer stroke and Peirce arrow, respectively. They are defined as $x \uparrow y := \neg(x \wedge y)$ and $x \downarrow y := \neg(x \vee y)$.

- 1.8** If we use standard rewriting rules for producing a CNF, then we usually conclude by some simplifications—remove duplicate clauses and literals. Why can we do that? Is it correct that there is no need for a variable to occur more than once in a clause?

- 1.9** Produce a formula in CNF which is equivalent to

$$\varphi = (a \rightarrow (c \wedge d)) \vee (b \rightarrow (c \wedge e)).$$

Then use the Tseytin transformation to produce a formula in CNF which is equisatisfiable to φ .

- 1.10** Use Boolean formula manipulation in PySAT or Limboole on the previous formula $\varphi = (a \rightarrow (c \wedge d)) \vee (b \rightarrow (c \wedge e))$.
- 1.11** Why is the formula produced by Tseytin transformation only equisatisfiable and not equivalent to the original formula?
- 1.12** Try solving SAT problems by hand in The SAT Game.