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**Question 1.**

Let  $X$  contain all real numbers from  $[0; 1]$  which can be represented using 256 bits. Let  $\mathcal{H} = X$ , and let the decision be given by  $H \in \mathcal{H}$  as

$$h(x) = 1 \text{ iff } x > H$$

Determine an  $m$  such that with probability at least 0.9,  $\text{err}(h) < 0.1$ , where  $h$  is an arbitrary hypothesis from  $\mathcal{H}$  consistent with  $m$  i.i.d. examples from  $X$ . Estimate it

- (a) without using any upper bounds seen in the lecture
- (b) using the upper bound with  $\ln |\mathcal{H}|$
- (c) using the upper bound with  $\text{VC}(\mathcal{H})$

**Answer:**

We have

$$\begin{aligned}\epsilon &= 0.1 \\ \delta &= 1 - 0.9 = 0.1 \\ |\mathcal{H}| &= 2^{256}\end{aligned}$$

- (a) For a fixed  $h$ , the probability that it is “bad” ( $\text{err}(h) > \epsilon$ ) and still consistent with  $m$  i.i.d. observations is at most  $(1 - \epsilon)^m = 0.9^m$ .

For an arbitrary  $h \in \mathcal{H}$ , we can bound the probability of at least one of them being “bad” by

$$\sum_{h \in \mathcal{H}} (1 - \epsilon)^m = |\mathcal{H}|(1 - \epsilon)^m = 2^{256}0.9^m$$

We want this probability to be smaller than  $\delta$ :

$$\begin{aligned}|\mathcal{H}|(1 - \epsilon)^m &< \delta \\ m &\geq \log_{1-\epsilon} \frac{\delta}{|\mathcal{H}|}\end{aligned}$$

i.e.,

$$m > \log_{0.9} \frac{0.1}{2^{256}} \approx 1707 \text{ examples (smallest such } m) \tag{1}$$

- (b)

$$\begin{aligned}m &> \frac{1}{\epsilon} \ln \frac{|\mathcal{H}|}{\delta} \\ m &> \frac{1}{0.1} \ln \frac{2^{256}}{0.1} \approx 1798 \text{ examples (smallest such } m)\end{aligned}$$

which is slightly greater than before because the upper bound  $(1 - \epsilon)^m < e^{-\epsilon m}$  ( $\epsilon > 0$ ) is used in the derivation of the formula.

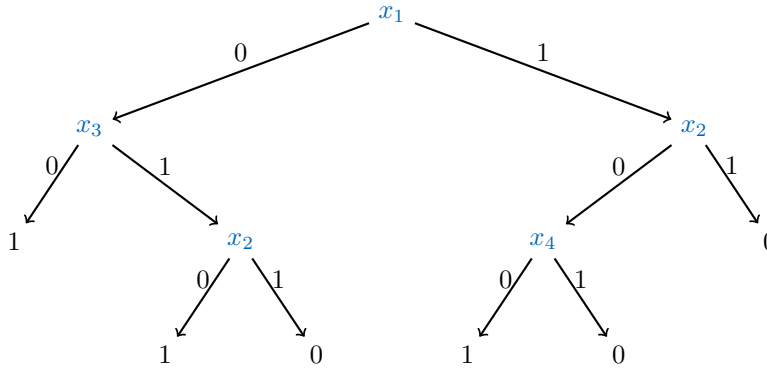
- (c)  $\text{VC}(\mathcal{H}) = 1$  because a single number from  $X$  can evidently be shattered (classified positively or negatively by hypotheses from  $\mathcal{H}$ ) but two different numbers from  $X$  cannot be shattered: the smaller cannot be made positive while the larger is negative.

$$m > \frac{8}{\epsilon} \left( \text{VC}(\mathcal{H}) \cdot \ln \frac{16}{\epsilon} + \ln \frac{2}{\delta} \right) \approx 646 \text{ examples}$$

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**Question 2.**

Consider the following decision tree:



- (a) Express the tree as a 3-DNF.
- (b) Express the tree as a 3-CNF.
- (c) How can we use (modify) the generalization algorithm to learn  $k$ -decision trees in the PAC learning model?

**Answer:**

- (a) A 3-DNF is a disjunction of minterms conjoining at most 3 literals.

We construct each minterm by following one path to a positive label. By disjoining all those paths, we get the final DNF tree representation.

$$(\neg x_1 \wedge \neg x_3) \vee (\neg x_1 \wedge x_3 \wedge \neg x_2) \vee (x_1 \wedge \neg x_2 \wedge \neg x_4)$$

- (b) A 3-CNF is a conjunction of maxterms (clauses) disjoining at most 3 literals.

We can use the fact that a negation of a DNF is a CNF. We can't negate the DNF constructed above directly, since then, we would consider paths going into the negative labels. However, we can construct a DNF going into the negative labels and negate that, giving us a CNF capturing the paths going into the positive labels.

$$\neg((\neg x_1 \wedge x_3 \wedge x_2) \vee (x_1 \wedge \neg x_2 \wedge x_4) \vee (x_1 \wedge x_2))$$

We could also construct the CNF directly. Consider each path to a negative label and “do everything in your power to avoid going down that path”. For example, considering the “path”  $\langle \neg x_1, x_3, x_2 \rangle$ , construct the clause  $(x_1 \vee \neg x_3 \vee \neg x_2)$ .

Regardless of the strategy used, the final 3-CNF reads as

$$(x_1 \vee \neg x_3 \vee \neg x_2) \wedge (\neg x_1 \vee x_2 \vee \neg x_4) \wedge (\neg x_1 \vee \neg x_2)$$

- (c) We can use the fact following from the exercises above that  $k$ -DT  $\subseteq k$ -CNF ( $k$ -DNF). Hence, we will use the generalization algorithm for learning  $k$ -CNFs ( $k$ -DNFs). Each clause (minterm) will be encoded by a new propositional variable and then we will simply be learning a conjunction (disjunction). There will be  $\sum_{i=1}^k \binom{n}{i} 2^i \leq \text{poly}(n)$  propositions, hence we will also learn efficiently. We won't be learning properly, though.

Using the adapted algorithm, we will be learning  $k$ -DTs in the mistake bound learning model, which implies that we will also be learning them in the PAC model.