# Uncertainty, Chance, and Utilities 

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Deterministic opponent $\rightarrow$ stochastic environment


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## Deterministic opponent $\rightarrow$ stochastic environment


$b_{1}, b_{2}, b_{3}$ - stochastic branches, uncertain outcomes of $a_{1}$ action.
CHANCE nodes are "virtual", $b_{1}, b_{2}, b_{3}$ are not actions!

Why? Actions may fail, ...


Video: Slipping robot. Vision for Robotics and Autonomous Systems, http://cyber.felk.cvut.cz/vras, https://youtu.be/kvEEHNyCHMs

Why? Action costs not deterministic, ..., getting to work
A At home
tram bike car

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Random variable: Function mapping situation on rails to values $T\left(r_{i}\right)=t_{i}$ :
$t_{1}=T\left(r_{1}\right)=3$ mins (free rails)
$t_{2}=T\left(r_{2}\right)=12$ mins (accident)
$t_{3}=T\left(r_{3}\right)=8 \mathrm{mins}$ (congestion)

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MAX/MIN depends on what the $t_{i}$ options and terminal numbers mean. The goal may be to get to work as fast as possible.

## Chance nodes values



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- Average case, not the worst case.
- Calculate expected utilities...


## Chance nodes values



- Average case, not the worst case.
- Calculate expected utilities...
- i.e. take weighted average (expectation) of successors

Random variables, probability distribution, ...

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A few reminders from laws of probability, Probabilities:

- always non-negative,
- sum over all possible outcomes is equal to 1 .


## Expectations, ...

How long does it take to go to work by tram?

- Depends on the random variable $T$ with possible values $t_{1}, t_{2}, t_{3}$ (corresponding to situation on rails).
- What is the expectation of the time?


## Expectations, ...

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Using values $t_{1}, t_{2}, t_{3}$ of random variable $T$ :

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E(T)=P\left(t_{1}\right) t_{1}+P\left(t_{2}\right) t_{2}+P\left(t_{3}\right) t_{3}
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$$
E(T)=P\left(t_{1}\right) t_{1}+P\left(t_{2}\right) t_{2}+P\left(t_{3}\right) t_{3}
$$

Or, using random outcomes $r_{1}, r_{2}, r_{3}$ :

$$
E(T)=P\left(r_{1}\right) T\left(r_{1}\right)+P\left(r_{2}\right) T\left(r_{2}\right)+P\left(r_{3}\right) T\left(r_{3}\right)
$$

Expected value of a discrete r.v.: Weighted average

## Expectimax

function Expectimax(state) return a value
if IS-TERMINAL(state): return UTILITY(state)
if state (next agent) is MAX: return MAX-VALUE(state)
if state (next agent) is CHANCE: return EXP-VALUE(state)
end function
function MAX-VALUE(state) return value $v$
$v \leftarrow-\infty$
for $a$ in ACTIONS(state) do
$v \leftarrow \max (v, \operatorname{EXPECTIMAX}(\operatorname{RESULT}($ state,$a)))$
end for
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end function
function EXP-VALUE(state) return value $v$
$v \leftarrow 0$
for all $r \in$ random outcomes do

$$
v \leftarrow v+P(r) \operatorname{EXPECTIMAX}(\operatorname{RESULT}(\text { state }, r))
$$

end for
end function

## How about the Reversi game?

- Is there any space for randomness?
- Is the opponent really greedy and clever enough?


## How about the Reversi game?

- Is there any space for randomness?
- Is the opponent really greedy and clever enough?
- Hope for chance when there is adversarial world - Dangerous optimism
- Assuming worst case even if it is not likely - Dangerous pessimism


## Games with chance and strategy



Random variable: Throwing two dice

Do we care which die comes first?

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A $1 / 24$
B $1 / 36$
C $1 / 18$
D $1 / 6$

[^0]
## Random variable: Throwing two dice

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## Mixing MAX, CHANCE, and MIN nodes



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Extra random agent that moves after each MAX and MIN agent

$$
\operatorname{EXPECTIMINIMAX}(s)=
$$

$$
=\{
$$



Extra random agent that moves after each MAX and MIN agent

$$
\operatorname{EXPECTIMINIMAX}(s)=
$$

$$
= \begin{cases}\operatorname{UTILITY}(s, \operatorname{MAX}) & \text { if } \operatorname{IS-TERMINAL}(s) \\ & \end{cases}
$$



Extra random agent that moves after each MAX and MIN agent

$$
\operatorname{EXPECTIMINIMAX}(s)=
$$

$$
=\left\{\begin{array}{lll}
\operatorname{UTiLITY}(s, \operatorname{mAx}) & \text { if } & \operatorname{IS-TERMINAL}(s) \\
\max _{a} \operatorname{EXPECTIMINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } & \operatorname{TO-PLAY}(s)=\operatorname{mAX} \\
& &
\end{array}\right.
$$



Extra random agent that moves after each MAX and MIN agent

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\min _{a} \operatorname{EXPECTImINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } & \operatorname{TO-PLAY}(s)=\text { MIN }
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\min _{a} \operatorname{EXPECTIMINIMAX}(\operatorname{RESULT}(s, a)) & \text { if } & \text { TO-PLAY }(s)=\text { MIN } \\
\sum_{r} P(r) \text { EXPECTIMINIMAX }(\operatorname{RESULT}(s, r)) & \text { if } & \text { TO-PLAY }(s)=\text { CHANCE }
\end{array}\right.
$$

Mixing chance into $\min / m a x$ tree. How big is the tree going to be?


- $b$ branching factor
- m maximum depth
- $n$ number of distinct rolls

What is the time complexity of EXPECTIMINIMAX?

A $O\left(b^{m n}\right)$
B $O\left(b^{m} n\right)$
C $O\left(b^{m} n^{b}\right)$
D $O\left(b^{m} n^{m}\right)$

Evaluation function MAX

MIN


## Evaluation function



- Left: $a_{1}$ is the best. Right: $a_{2}$ is the best. Ordering of the (terminal) leaves is the same.


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- Scale matters! Not only ordering.


## Evaluation function

## MAX

CHANCE

MIN


- Left: $a_{1}$ is the best. Right: $a_{2}$ is the best. Ordering of the (terminal) leaves is the same.
- Scale matters! Not only ordering.
- Can we prune the tree? ( $\alpha, \beta$ like?)

Pruning expectiminimax tree


## Pruning expectiminimax tree

MAX


## TERMINAL

- Bounds on terminal utilities needed. Terminal values from -2 to 2 .
- Monte Carlo simulation for evaluation of a position (state).


## Where to prune the Expectimax tree

- Assume terminal nodes bounded to -2 to 2 , inclusive
- Going from left to right.
- Which branches can be pruned out?



Assume terminal nodes bounded to -2 to 2 , inclusive. Going from left to right.

## Multi-player games

to move


## Multi-player games

to move
A

B


- Utility tuples
- Each player maximizes its own
- Coalitions, cooperations, competitions may be dynamic

Uncertainty recap (enough games, back to the robots/agents)


- Uncertain outcome of an action.

Uncertainty recap (enough games, back to the robots/agents)


- Uncertain outcome of an action.
- Robot/Agent may not know the current state!

Uncertain outcome of an action


Uncertain, partially observable environment

- Current state $s$ may be unknown, observations e
- Uncertain outcome, random variable RESULT(a)
- Probability of outcome $s^{\prime}$ given $\mathbf{e}$ is

$$
P\left(\operatorname{RESULT}(a)=s^{\prime} \mid a, \mathbf{e}\right)
$$

- Utility function $U(s)$ corresponds to agent preferences.
- Expected utility of an action a given e:

$$
E U(a \mid \mathbf{e})=\sum_{s^{\prime}} P\left(\operatorname{RESULT}(a)=s^{\prime} \mid a, \mathbf{e}\right) U\left(s^{\prime}\right)
$$



## Rational agent

Agent's expected utility of an action a given e:

$$
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What should a rational agent do?

## Rational agent

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Is it then all solved? Do we know all what we need?

## Rational agent

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What should a rational agent do?
Is it then all solved? Do we know all what we need?

- $P\left(\operatorname{Result}(a)=s^{\prime} \mid a, \mathbf{e}\right)$
- $U\left(s^{\prime}\right)$


## Utilities



- Where do utilities come from?
- Does averaging make sense?
- Do they exist?
- What if our preferences can't be described by utilities?


## Agent/Robot Preferences

- Prizes $A, B$
- Lottery: uncertain prizes $L=[p, A ;(1-p), B]$


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Preference, indifference, ...

- Robot prefers $A$ over $B: A \succ B$
- Robot has no preferences: $A \sim B$
- in between: $A \succsim B$


## Rational preferences

- Transitivity: $(A \succ B) \wedge(B \succ C) \Rightarrow(A \succ C)$
- Completeness: $(A \succ B) \vee(B \succ A) \vee(A \sim B)$
- Continuity: $(A \succ B \succ C) \Rightarrow \exists p[p, A ; 1-p, C] \sim B$
- Substituability: $A \sim B \Rightarrow[p, A ; 1-p, C] \sim[p, B ; 1-p C]$. The same for $\succ$ and $\sim$.
- Monotonocity: $A \succ B \Rightarrow(p>q) \Leftrightarrow[p, A ; 1-p, B] \succ[q, A ; 1-q, B]$. Agent must prefer a lottery with higher chance to win.
- Decomposability, compressing compound lotteries into one:
$[p, A ; 1-p,[q, B ; 1-q, C]] \sim[p, A ;(1-p) q, B ;(1-p)(1-q), C]$


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Axioms of utility theory.
Motivation: if agent/robot violates an axiom $\Rightarrow$ irrational agent/robot.

## Transitivity and decomposability

Goods $A, B, C$ and (nontransitive) preferences of an (irrational) agent $A \succ B \succ C \succ A$.

is equivalent to

(a)
(b)

## Maximum expected utility principle

Given the rational preferences (constraints), there exists a real valued function $u$ such that:

$$
\begin{aligned}
& u(A)>u(B) \quad \Leftrightarrow \quad A \succ B \\
& u(A)=u(B) \quad \Leftrightarrow \quad A \sim B
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Expected utility of a Lotery $L$ (outcomes $s_{i}$ with probabilities $p_{i}$ ):

$$
L\left(\left[p_{1}, S_{1} ; \cdots ; p_{n}, S_{n}\right]\right)=\sum_{i} p_{i} u\left(S_{i}\right)
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Proof in [3].

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Proof in [3].
Is a utility $u$ function unique?

Human utilities


## Utility of money

You triumphed in a TV show!
a) Take $\$ 1,000,000 \ldots$ or
b) Flip a coin and loose all or win $\$ 2,500,000$

## Utility of money: human psychology vs. hard data



## Utility of money: human psychology vs. hard data



## References I

Some figures from [2], Chapters 5, 16. Human utilities are discussed in [1]. This lecture has been also greatly inspired by the 7th lecture of CS 188 at http://ai.berkeley.edu as it convenietly bridges the world of deterministic search and sequential decisions in uncertain worlds.
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## References II

[3] John von Neumann and Oskar Morgenstern.
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[^0]:    ${ }^{1}$ Source of dice images: https://flyclipart.com/dice-clipart-tool-rolling-dice-clipart-248574

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