

# Problem solving by search II

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# Outline

- ▶ Graph search
- ▶ Heuristics (how to search faster)
- ▶ Greedy
- ▶  $A^*$ . A-star search.

# A Maze, what could possibly go wrong?

	0	1	2	3	4	
0	0.00	0.00	0.00	0.00	0.00	0
1	0.00	0.00	0.00	0.00	0.00	1
2	0.00	0.00	0.00	0.00	0.00	2
3	0.00	0.00	0.00	0.00	0.00	3
4	0.00	0.00	0.00	0.00	0.00	4
	0	1	2	3	4	

<https://youtu.be/WKSoedfRZQ4>

# Tree search the maze

**function** TREE\_SEARCH(env) **return** a solution or failure

initialize the frontier

**while** frontier **do**

node  $\leftarrow$  frontier.pop()

**if** goal in node **then return** node

**end if**

child\_nodes  $\leftarrow$  env.expand(node.state)

Add child\_nodes to frontier

**end while**

**end function**

	0	1	2	3	4	
0	0.00	0.00	0.00	0.00	0.00	0
1	0.00	0.00	0.00	0.00	0.00	1
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3	0.00	0.00	0.00	0.00	0.00	3
4	0.00	0.00	0.00	0.00	0.00	4
	0	1	2	3	4	

## A graph search

**function** GRAPH\_SEARCH(env) **return** a solution or failure

  init **frontier** by the start state

*initialize the explored set to be empty*

**while** frontier **do**

    node ← frontier.pop()

*add node.state to explored*

**if** goal in node **then return** node

**end if**

    child\_nodes ← env.expand(node.state)

**for all** child\_nodes **do**

*if child\_node.state not in explored then*

        add nodes to frontier

**end if**

**end for**

**end while**

**end function**



Do not forget: node is not the same as state!

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▷ What about frontier?



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**Do not forget:** `node` is not the same as `state`!

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▷ What about frontier?



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## The BFS graph search

**function** BFS\_GRAPH\_SEARCH(env) **return** a solution or failure

node  $\leftarrow$  env.observe()

frontier  $\leftarrow$  FIFOqueue(node)

explored  $\leftarrow$  set()

**while** frontier not empty **do**

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explored.add(node.state)

▷ Add state, not node!

child\_nodes  $\leftarrow$  env.expand(node.state)

**for all** child\_nodes **do**

if child\_node.state not in explored and not in frontier **then**

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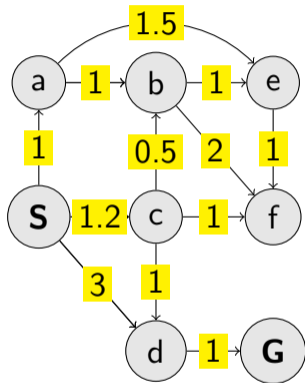
**end if**

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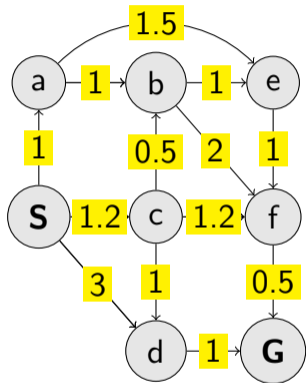
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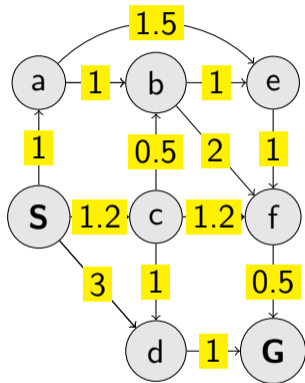
## Checking frontier in uniform costs graph search?



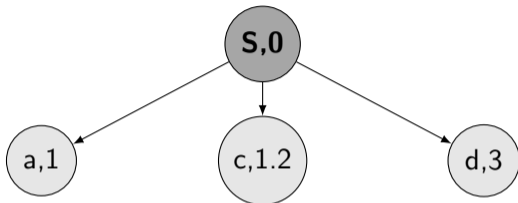
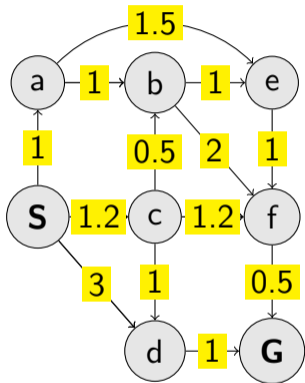
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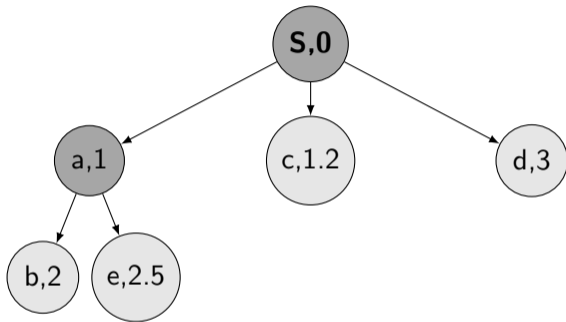
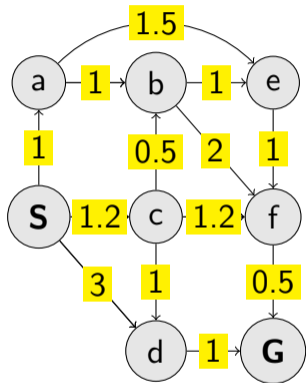
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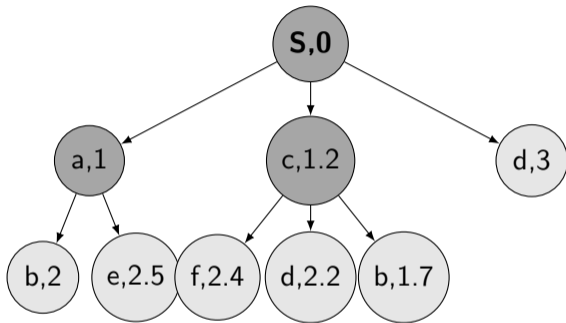
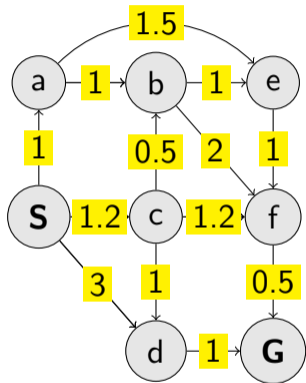
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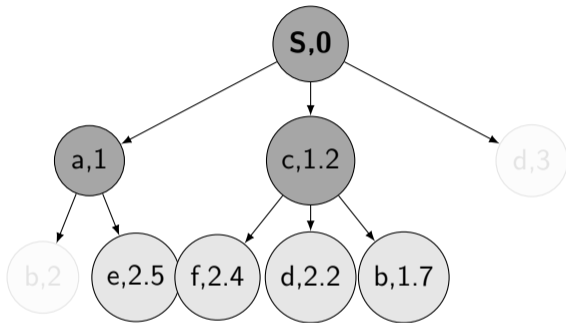
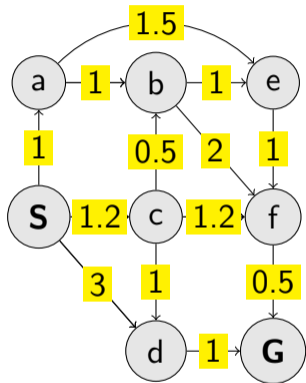


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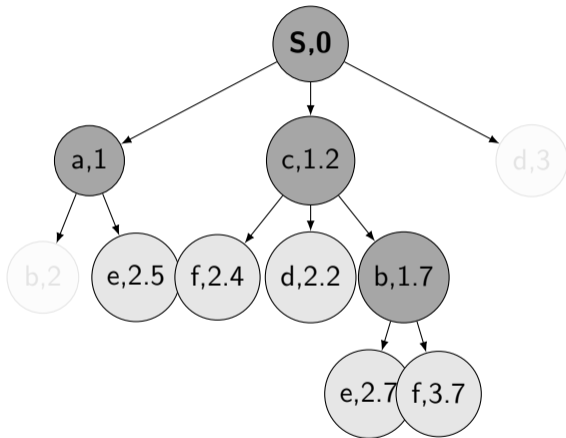
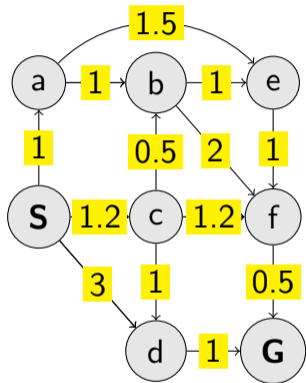




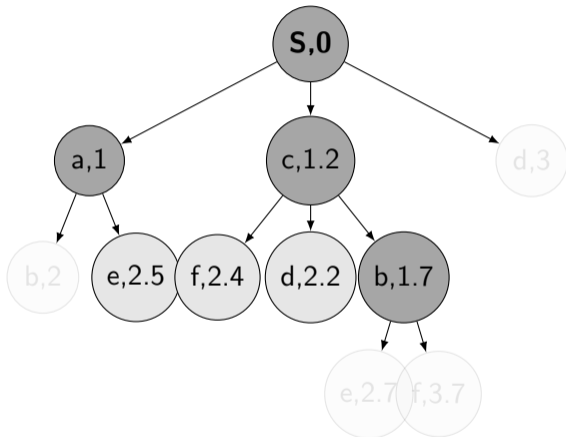
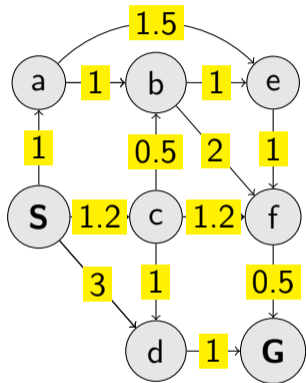
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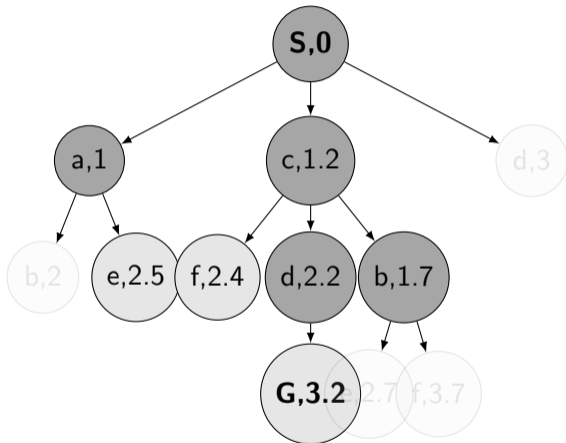
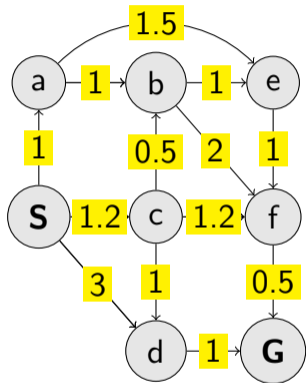
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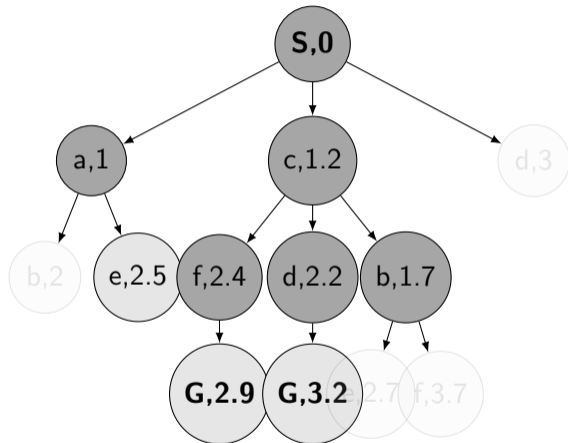
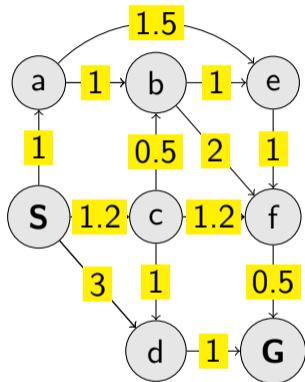
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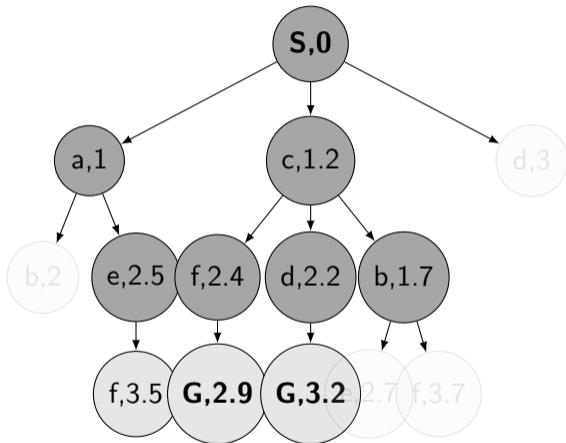
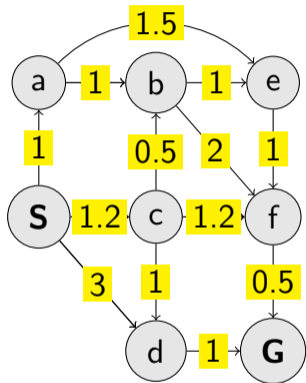
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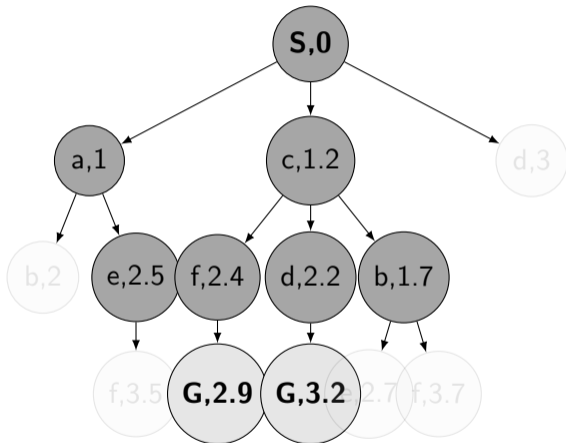
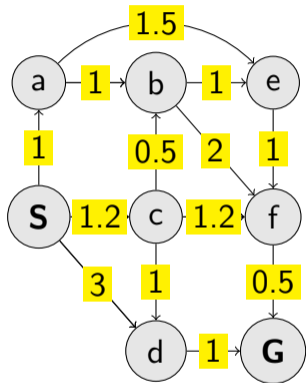
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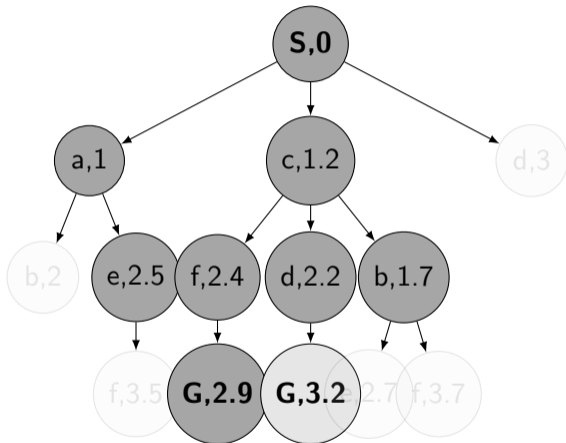
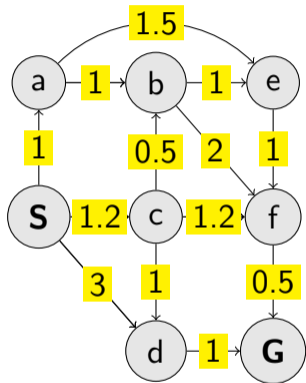
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## The UCS graph search

**function** UCS\_GRAPH\_SEARCH(env) **return** a solution or failure

node ← env.observe()

frontier ← priority\_queue(node)

explored ← set()

**while** frontier not empty **do**

node ← frontier.pop()

**if** node contains Goal **then return** node

**end if**

explored.add(node.state)

child\_nodes ← env.expand(node.state)

**for all** child\_nodes **do**

**if** child\_node.state not in explored and not in frontier **then**

frontier.insert(child\_node)

**else if** child\_node.state in frontier with higher cost **then**

replace that node with the child\_node

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▷ path\_cost for ordering

▷ check here!

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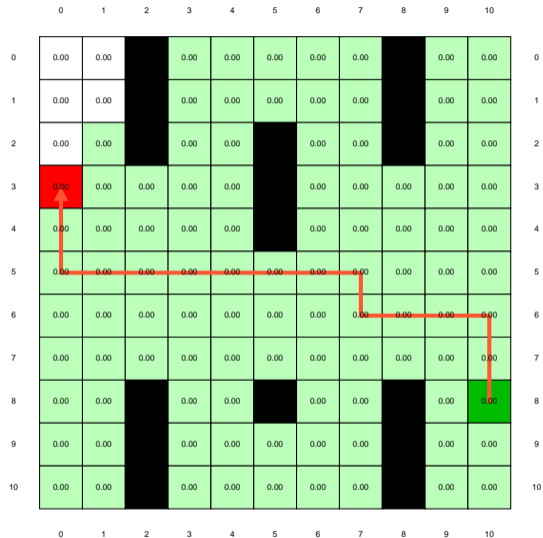
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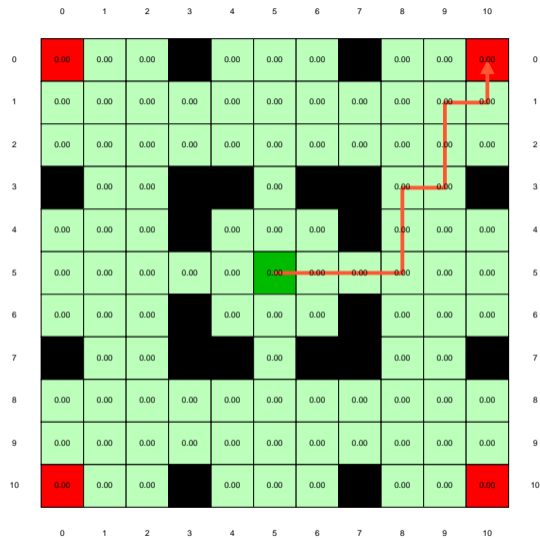
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# Few examples of search strategies so far



Run the demos.

# What is wrong with UCS and other strategies?



Run the demo, or see <https://youtu.be/TT5MY8xCgAg>

## Node selection, take argmin $f(n)$

Selecting next node to expand/visit (pop operation):

$$\text{node} \leftarrow \underset{n \in \text{frontier}}{\text{argmin}} f(n)$$

What is  $f(n)$  for DFS, BFS, and UCS?

- ▶ DFS:  $f(n) = n.depth$
- ▶ BFS:  $f(n) = -n.depth$
- ▶ UCS:  $f(n) = n.path\_cost$

The good: (one) frontier as a priority queue

(i.e., priority queue will work universally. Still, stack (LIFO) and queue (FIFO) are (conceptually) the perfect data structures for DFS and BFS, respectively.)

The bad: All the  $f(n)$  correspond to the cost from  $n$  to the start - only backward cost; cost-to-come (to  $n$ ).



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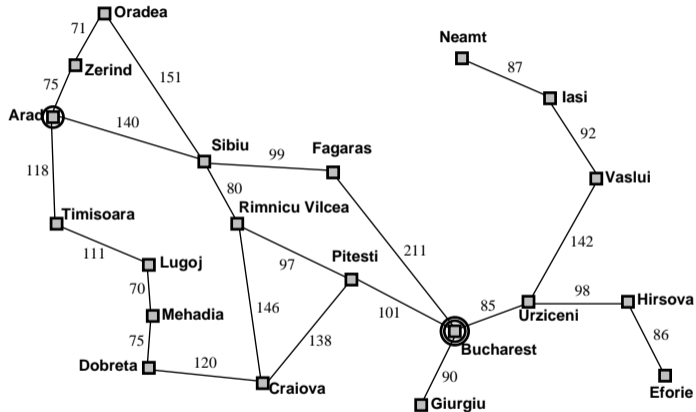
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The bad: All the  $f(n)$  correspond to the cost from  $n$  to the start - only **backward** cost; **cost-to-come** (to  $n$ ).

## How far are we from the goal **cost-to-go** ? – Heuristics

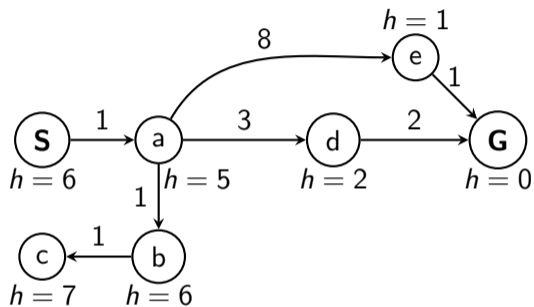
- ▶ A function that estimates how close a *state* is to the goal.
- ▶ Designed for a particular problem.
- ▶  $h(n.state)$  – it is function of the state (attribute of node)
- ▶ It is often shortened as  $h(n)$  – heuristic value of node  $n$ .

# Example of heuristics



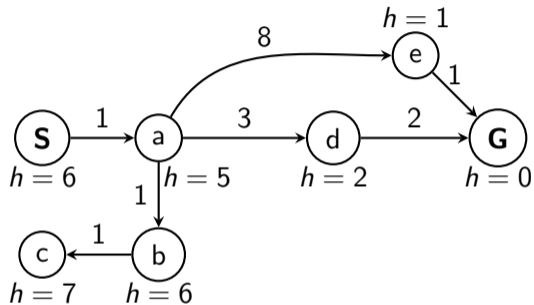
Arad	366	Mehadia	241
Bucharest	0	Neamt	234
Craiova	160	Oradea	380
Drobeta	242	Pitesti	100
Eforie	161	Rimnicu Vilcea	193
Fagaras	176	Sibiu	253
Giurgiu	77	Timisoara	329
Hirsova	151	Urziceni	80
Iasi	226	Vaslui	199
Lugoj	244	Zerind	374

Greedy, take the  $n^* = \operatorname{argmin}_{n \in \text{frontier}} h(n)$



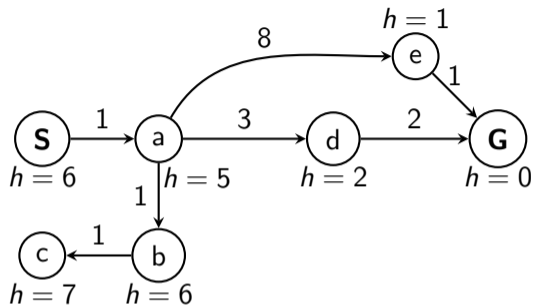
What is wrong (and nice) with the Greedy?

Greedy, take the  $n^* = \operatorname{argmin}_{n \in \text{frontier}} h(n)$



What is wrong (and nice) with the Greedy?

## A\* combines UCS and Greedy



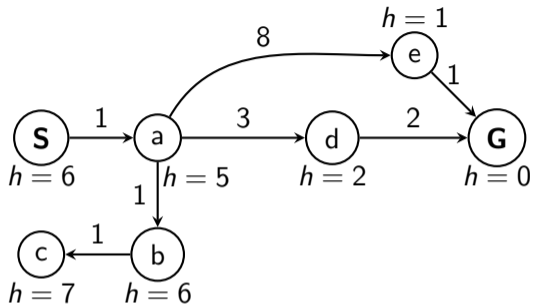
UCS orders by backward (path) cost  $g(n)$

Greedy uses heuristics (goal proximity)  $h(n)$

A\* orders nodes by:  $f(n) = g(n) + h(n)$



## A\* combines UCS and Greedy

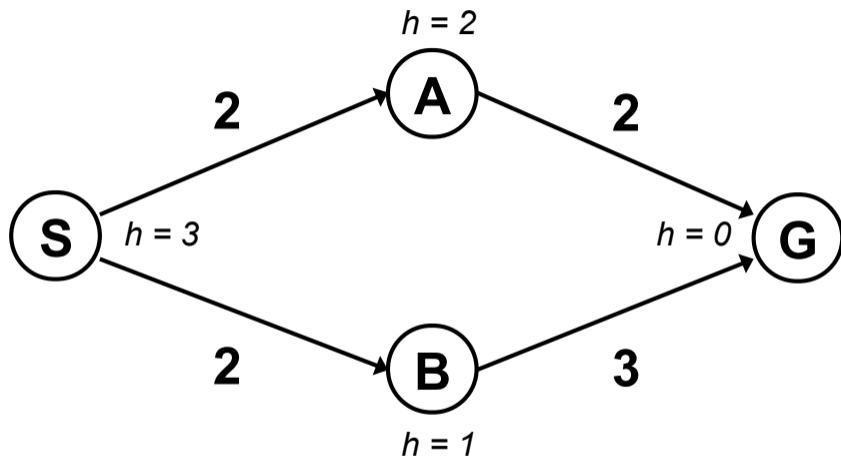


UCS orders by backward (path) cost  $g(n)$

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## When to stop A\*?

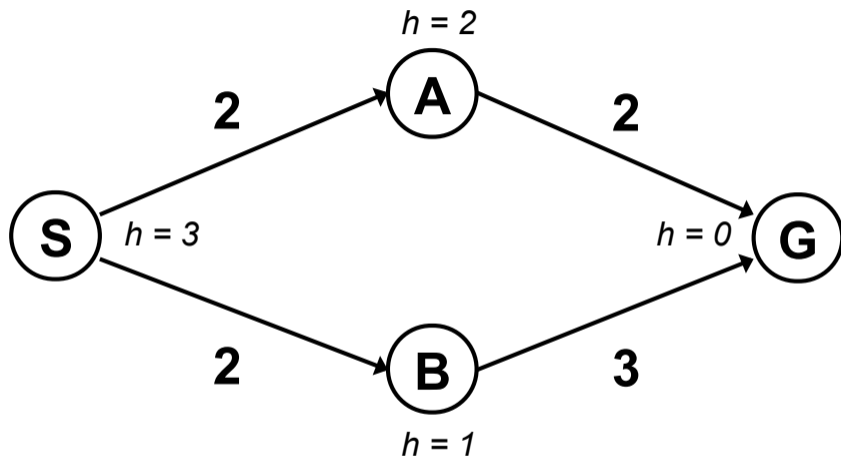


1

When popping G from frontier.

<sup>1</sup>Graph example: Dan Klein and Pieter Abbeel

## When to stop A\*?

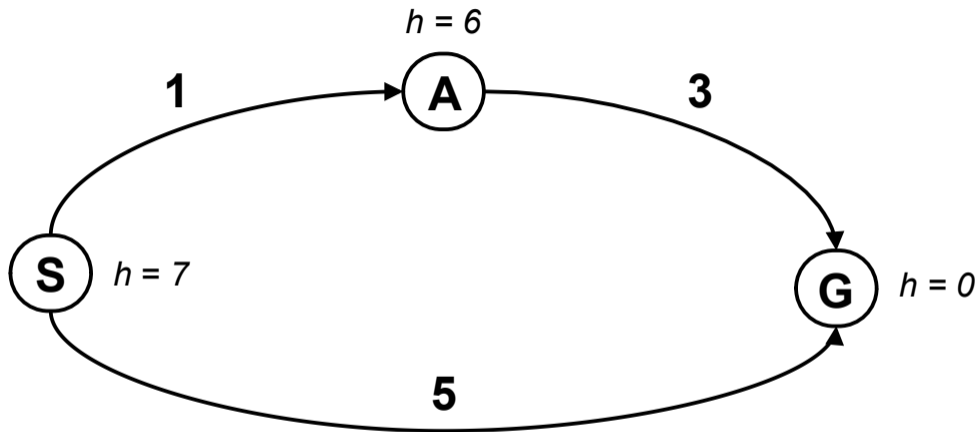


1

When popping G from frontier.

<sup>1</sup>Graph example: Dan Klein and Pieter Abbeel

Is  $A^*$  optimal?

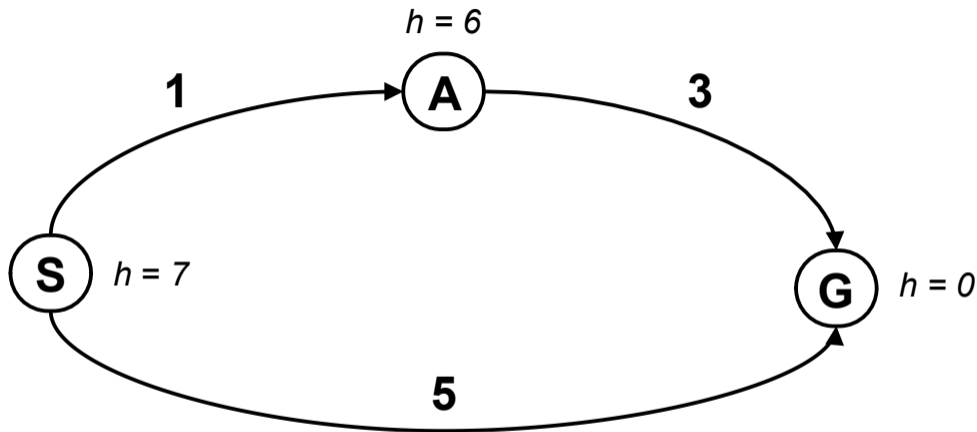


2

What is the problem?

<sup>2</sup>Graph example: Dan Klein and Pieter Abbeel

Is  $A^*$  optimal?



2

What is the problem?

---

<sup>2</sup>Graph example: Dan Klein and Pieter Abbeel

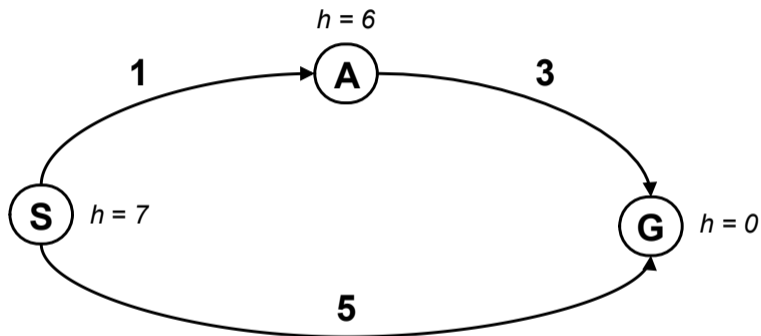
What is the right  $h(A)$ ?

A:  $0 \leq h(A) \leq 4$

B:  $h(A) \leq 3$

C:  $0 \leq h(A) \leq 3$

D:  $0 \leq h(A)$



## Admissible heuristics

A heuristic function  $h$  is admissible if:

$$\begin{aligned}h(n) &\leq \text{cost}(n.\text{state}, \text{Goal}_{\text{nearest}}) \\h(\text{Goal}) &= 0\end{aligned}$$

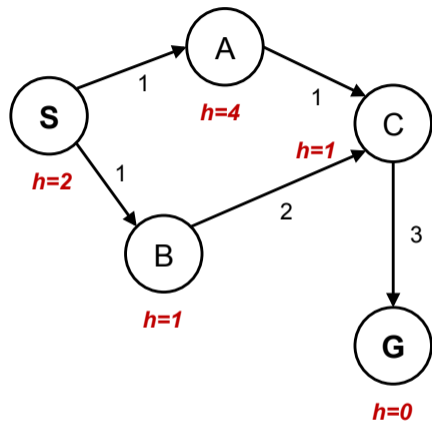
## Optimality of $A^*$ tree search

$A^*$  is optimal if  $h(n)$  is admissible.



# A\* graph search

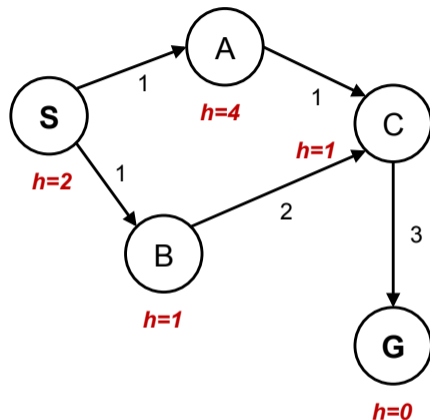
```
function GRAPH_SEARCH(env)
  frontier.insert(startnode)
  explored ← set()
  while frontier do
    node ← frontier.pop()
    if goal in node then return node
    end if
    child_nodes ← env.expand(node.state)
    explored.add(node.state)
    for all child_nodes do
      if child_node.state not in explored then
        frontier.insert(child_node)
      end if
    end for
  end while
end function
```



What went wrong?

# A\* graph search

```
function GRAPH_SEARCH(env)
  frontier.insert(startnode)
  explored ← set()
  while frontier do
    node ← frontier.pop()
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    child_nodes ← env.expand(node.state)
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    for all child_nodes do
      if child_node.state not in explored then
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    end for
  end while
end function
```



What went wrong?

## What would be the proper $h(A)$ ?

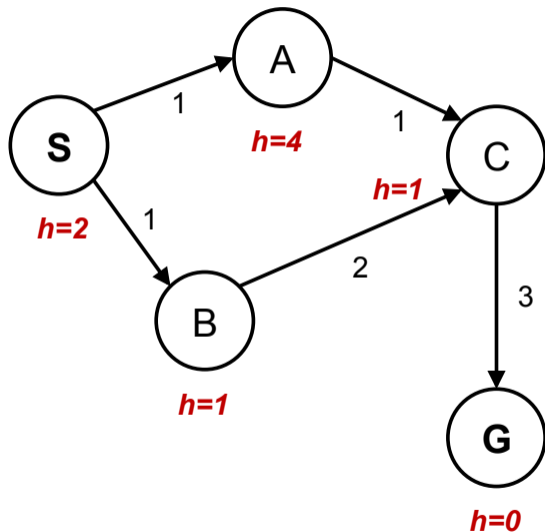
Consider other  $h(s)$  fixed.

A:  $h(A) = 1$

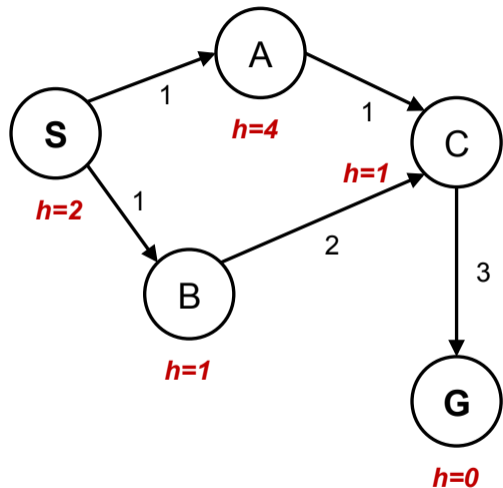
B:  $h(A) = 2$

C:  $1 \leq h(A) \leq 2$

D:  $0 \leq h(A) \leq 1$



## Consistent heuristics



Admissible  $h$ :

$$h(A) \leq \text{true cost } A \rightarrow G$$

Consistent  $h$ :

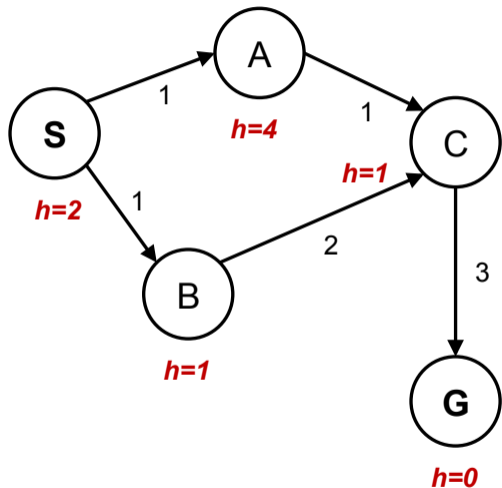
$$h(A) - h(C) \leq \text{true cost } A \rightarrow C$$

in general:

$$h(n) - h(s) \leq \text{true cost } n \rightarrow s \text{ for any pair: node } n \text{ and its successor } s$$

$f(n) = g(n) + h(n)$  along a path never decreases!

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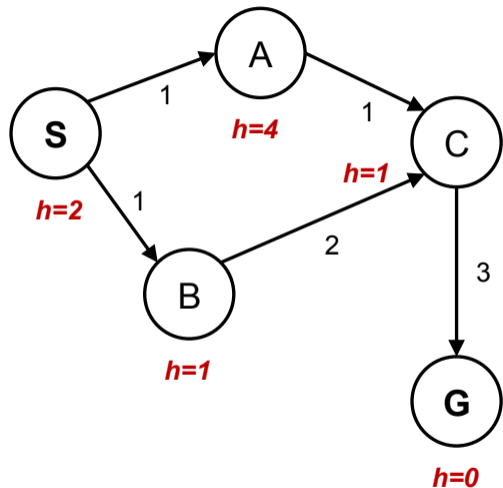
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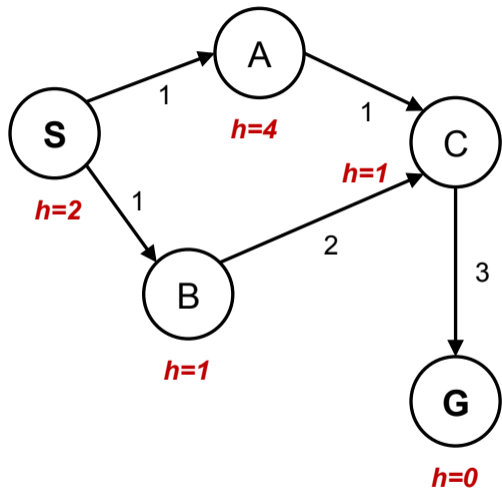
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$f(n) = g(n) + h(n)$  along a path never decreases!

## Optimality of $A^*$

- ▶ admissible  $h$  for tree search
- ▶ consistent  $h$  for graph search
- ▶ Are all consistent heuristics also admissible?

$$h(A) - h(C) \leq \text{cost}(A \rightarrow C)$$



## Optimality of $A^*$

- ▶ admissible  $h$  for tree search
- ▶ consistent  $h$  for graph search
- ▶ Are all consistent heuristics also admissible?  
 $h(A) - h(C) \leq \text{cost}(A \rightarrow C)$

# Summary

- ▶ Graph vs Tree search – avoid repeating expansions
- ▶ Effectiveness – adding heuristic estimates of cost-to-go
- ▶ Not all heuristics are equally good (admissibility, consistence, informativeness)

## References, further reading

Some figures from [2]. Chapter 2 in [1] provides a compact/dense intro into search algorithms. (State space) Search algorithms are ubiquitous, explanations in many (text)books about Algorithms.

Nice online course from UC Berkeley (CS 188 Intro to AI):

[http://ai.berkeley.edu/lecture\\_videos.html](http://ai.berkeley.edu/lecture_videos.html) Lecture: Informed Search.

[1] Steven M. LaValle.

*Planning Algorithms.*

Cambridge, 1st edition, 2006.

Online version available at: <http://planning.cs.uiuc.edu>.

[2] Stuart Russell and Peter Norvig.

*Artificial Intelligence: A Modern Approach.*

Prentice Hall, 3rd edition, 2010.

<http://aima.cs.berkeley.edu/>.