

ALG 07

Selection sort (Select sort)

Insertion sort (Insert sort)

Bubble sort deprecated

Quicksort

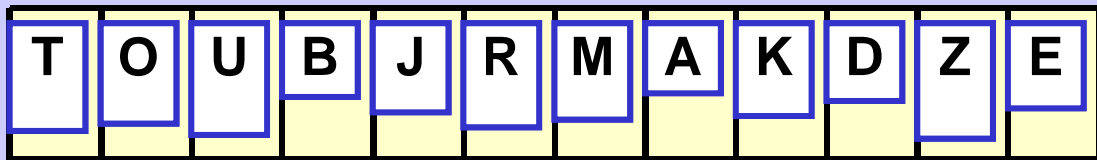
Sort stability

Selection sort

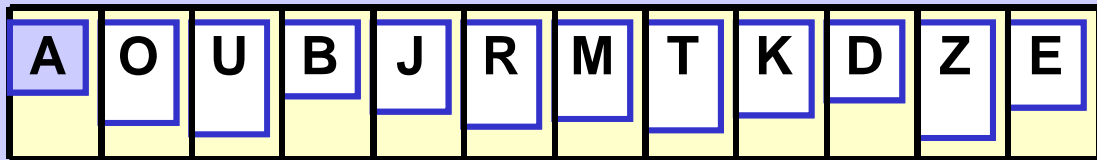
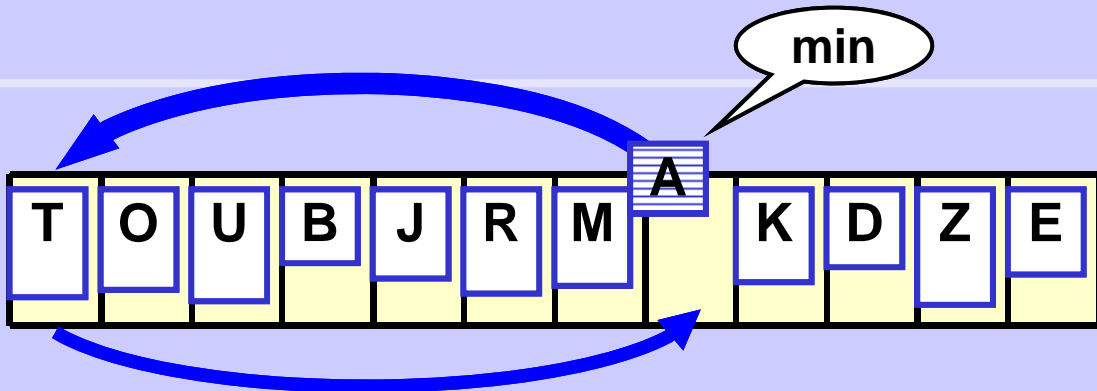
Unsorted

Sorted

Start

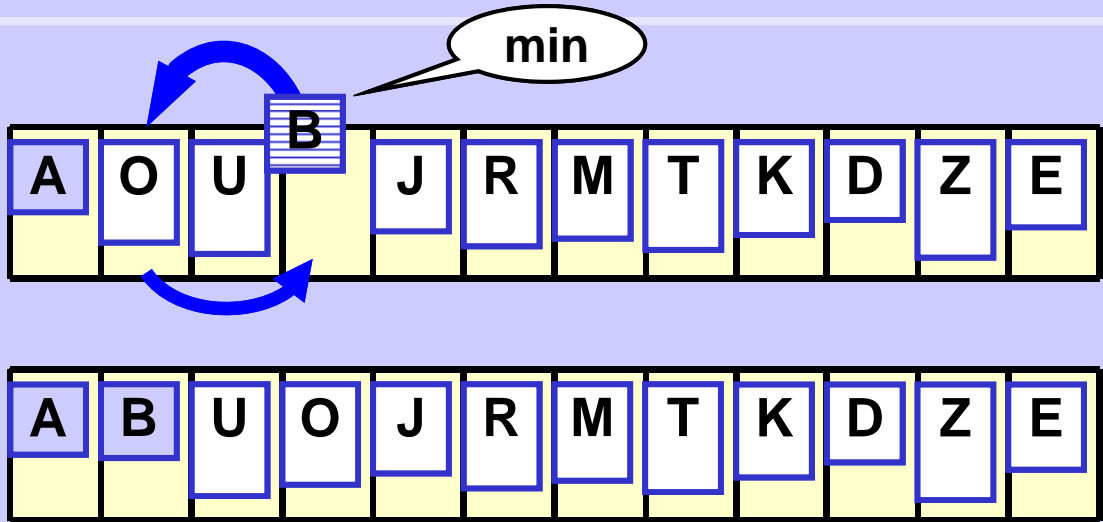


Step1

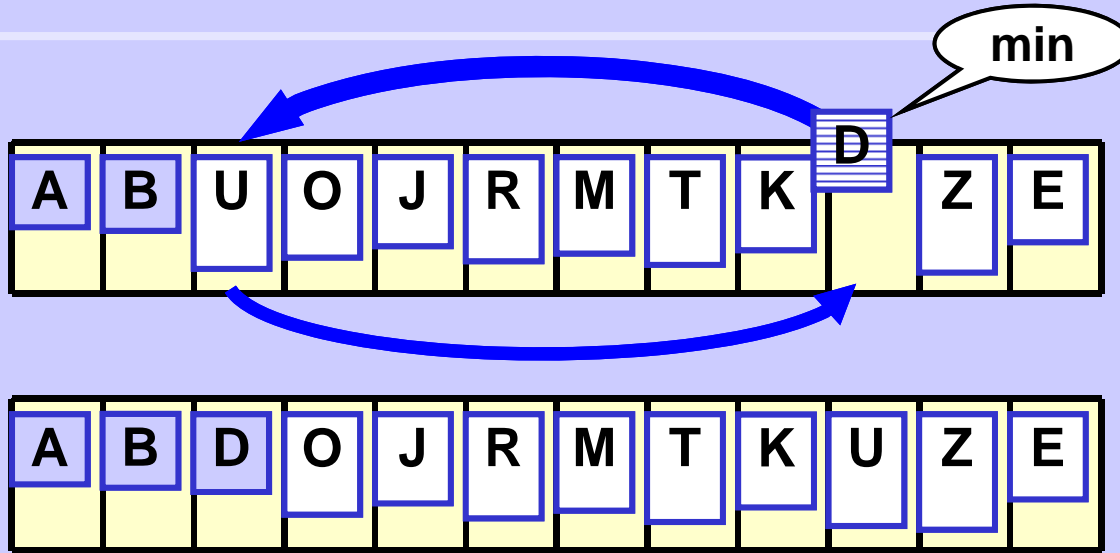


Selection sort

Step 2



Step 3

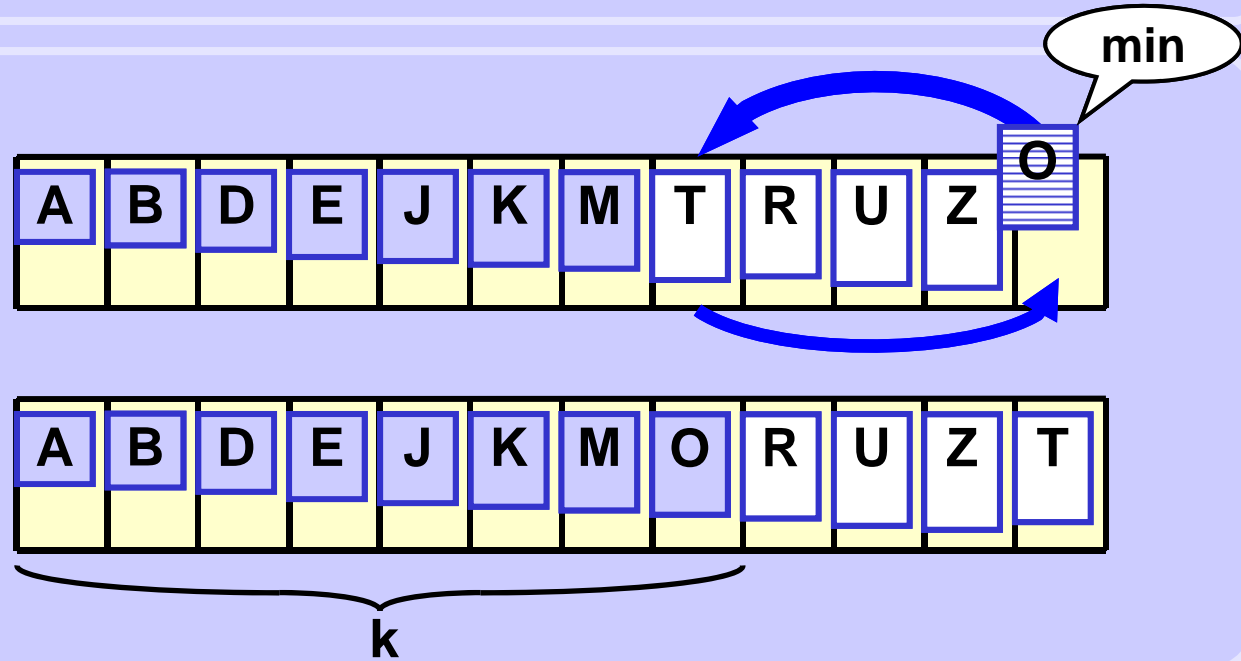


Selection sort

...

...

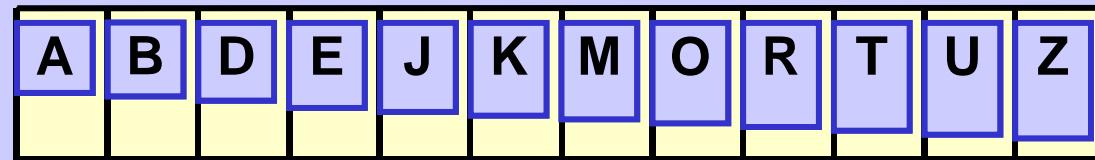
Step k



...

...

All sorted

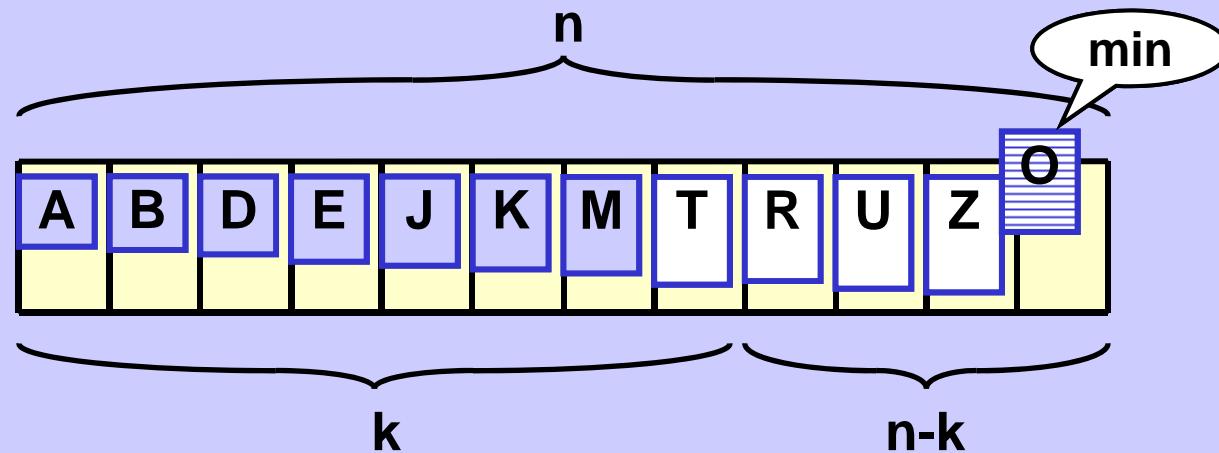


Selection sort

```
def selectSort( arr ):
    n = len( arr )
    for i in range( n-1 ):
        jmin = i
        # select minimum
        for j in range( i+1, n ):
            if arr[j] < arr[jmin]:
                jmin = j
        # put minimum to its place
        swap( arr, i, jmin )
```

Selection sort

Step k



Select minimum



.....

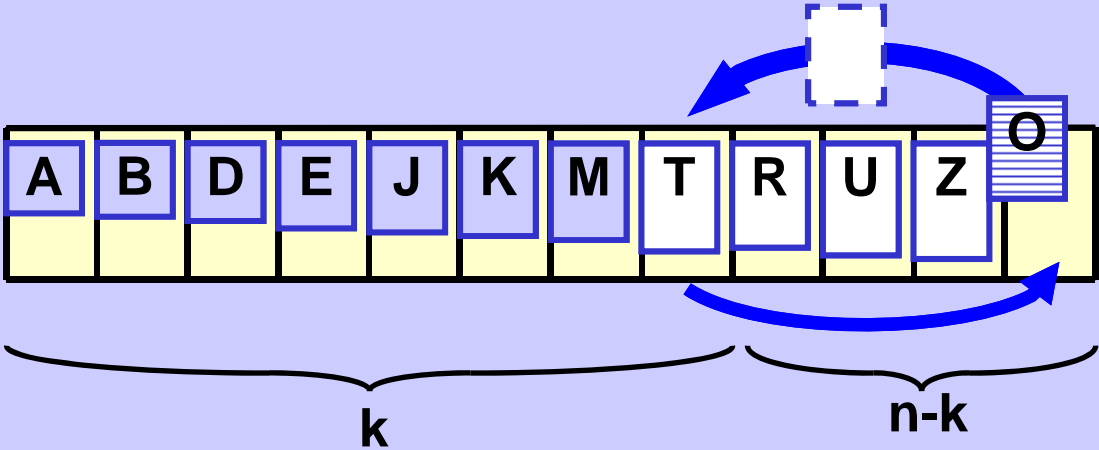
(n-k) tests

Tests
total

$$\sum_{k=1}^{n-1} (n-k) = \sum_{k=1}^{n-1} n - \sum_{k=1}^{n-1} k = n(n-1) - \frac{n(n-1)}{2} = \frac{1}{2}(n^2 - n)$$

Selection sort

Step k



Moves 3

Moves total

$$\sum_{k=1}^{n-1} 3 = 3(n-1)$$

Selection sort**Resume****Tests
total**

$$\frac{1}{2}(n^2 - n) = \Theta(n^2)$$

**Moves
total**

$$3(n-1) = \Theta(n)$$

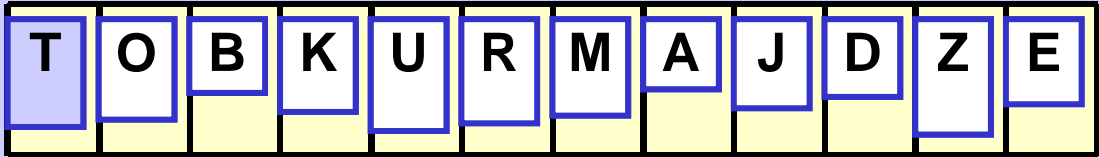
**Operations
total**

$$\frac{1}{2}(n^2 - n) + 3(n-1) = \Theta(n^2)$$

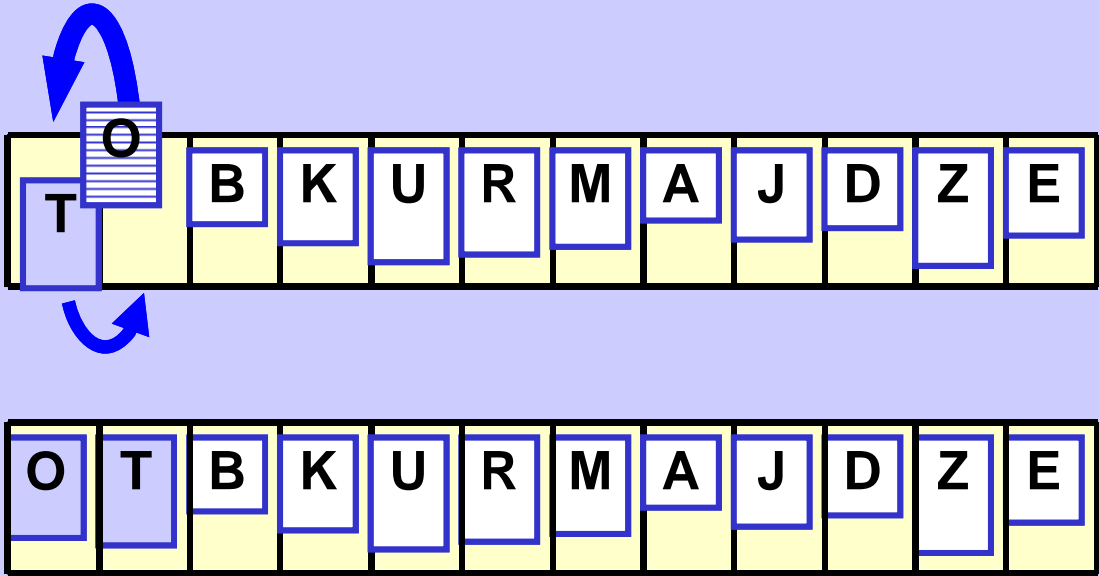
Asymptotic complexity of Selection Sort is $\Theta(n^2)$

Insertion sort

Start

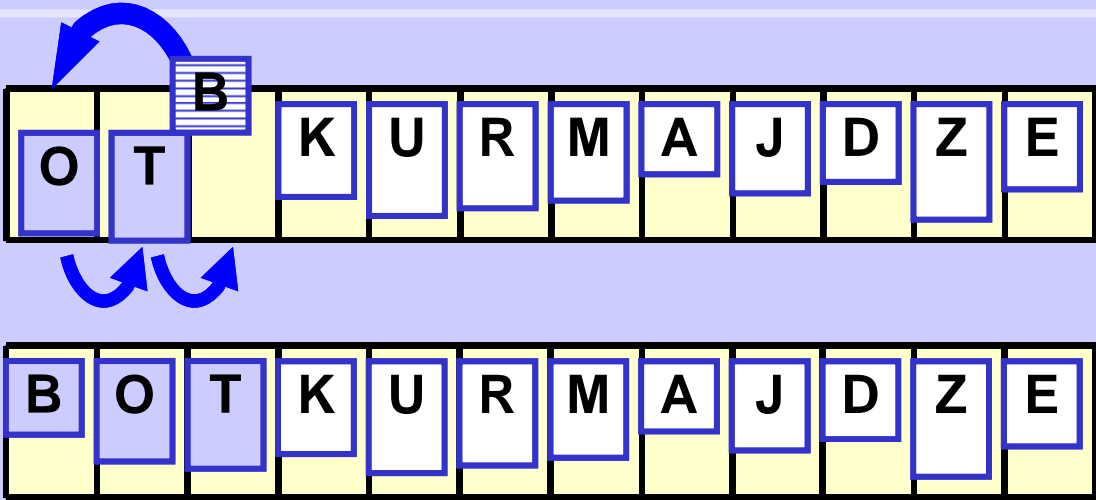


Step 1

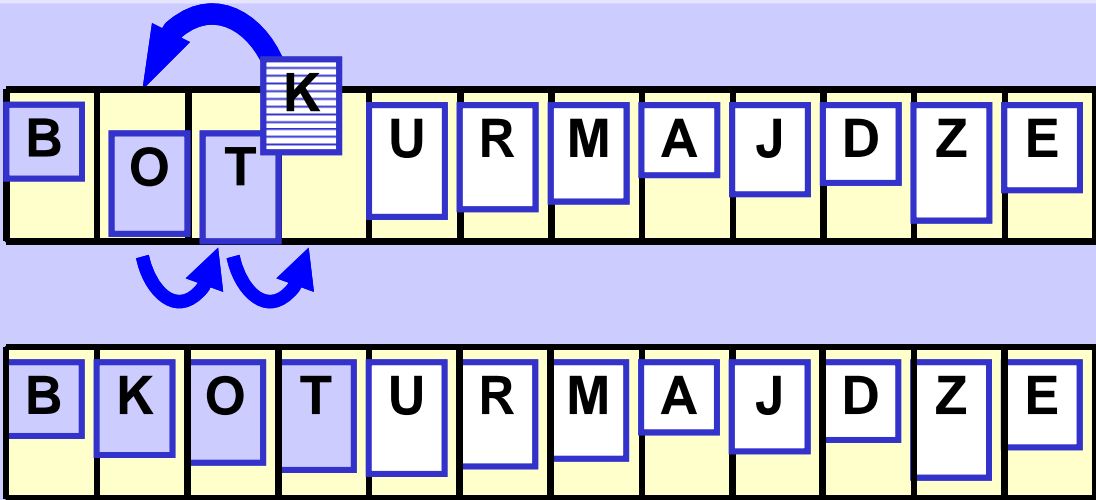


Insertion sort

Step 2



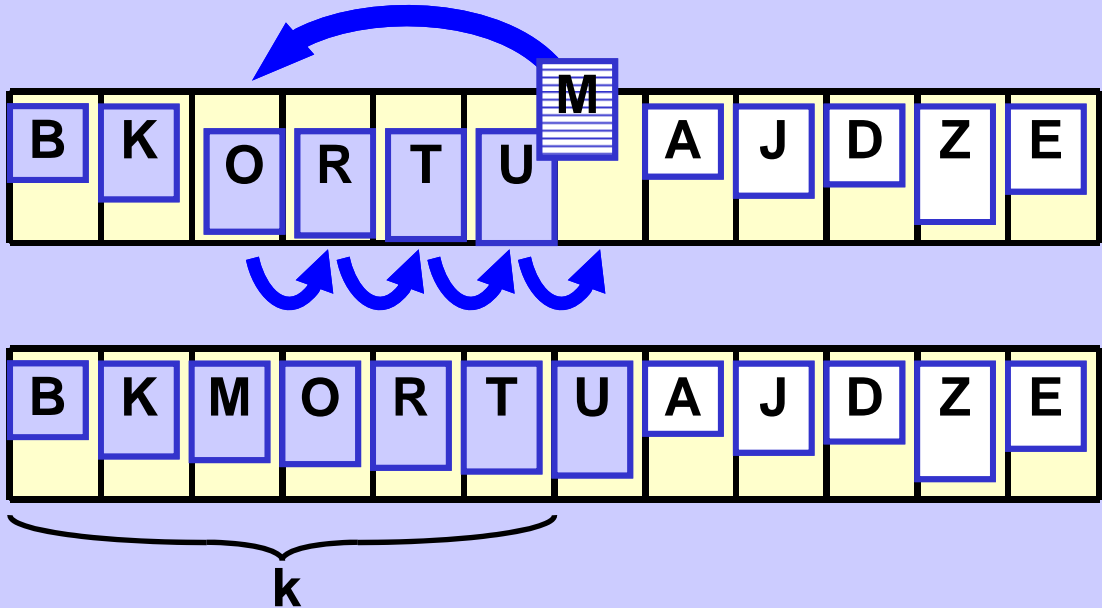
Step 3



Insertion sort

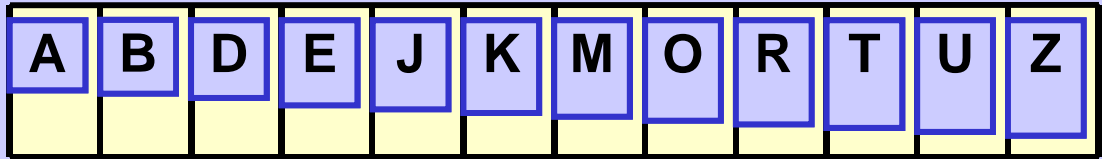
...

Step k



...

All sorted

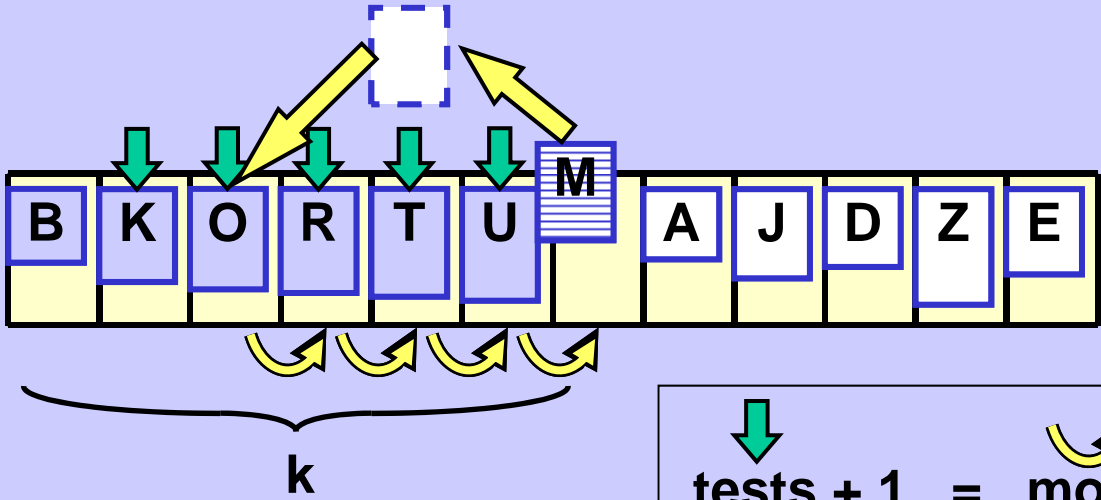





Insertion sort

```
def insertSort( arr ):
    n = len( arr )
    for i in range( 1, n ):
        # find & make place for arr[i]
        insVal = arr[i]
        j = i-1
        while j >= 0 and arr[j] > insVal:
            arr[j+1] = arr[j]
            j -= 1;
        # insert arr[i] to the correct place
        arr[j+1] = insVal
```

Insertion sort

Step k



tests + 1 = moves

tests	1	best case
	k	worst case
	$(k+1)/2$	“average” case

moves	2	best case
	k+1	worst case
	$(k+3)/2$	“average” case

Insertion sort

Resume

Tests
total

$n - 1$	$= \Theta(n)$	best case
$(n^2 - n)/2$	$= \Theta(n^2)$	worst case
$(n^2 + n - 2)/4$	$= \Theta(n^2)$	“average” case

Moves
total

$2n - 2$	$= \Theta(n)$	best case
$(n^2 + n - 2)/2$	$= \Theta(n^2)$	worst case
$(n^2 + 5n - 6)/4$	$= \Theta(n^2)$	“average” case

Asymptotic complexity of Insertion sort is $O(n^2)$ (!!)

Bubble sort

Start

T	O	B	U	J	R	M	A	K	D	Z	E
---	---	---	---	---	---	---	---	---	---	---	---

Phase 1

T	O	B	U	J	R	M	A	K	D	Z	E
---	---	---	---	---	---	---	---	---	---	---	---

Blue arrows indicate a swap between 'T' and 'O'.

O	T	B	U	J	R	M	A	K	D	Z	E
---	---	---	---	---	---	---	---	---	---	---	---

Blue arrows indicate a swap between 'T' and 'B'.

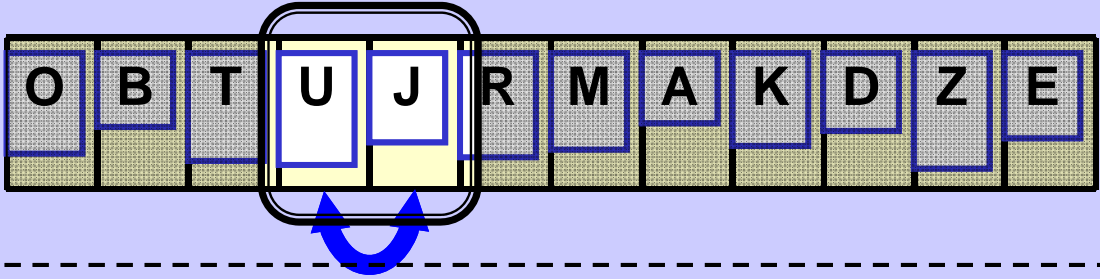
O	B	T	U	J	R	M	A	K	D	Z	E
---	---	---	---	---	---	---	---	---	---	---	---

Blue arrows indicate a swap between 'T' and 'U'.

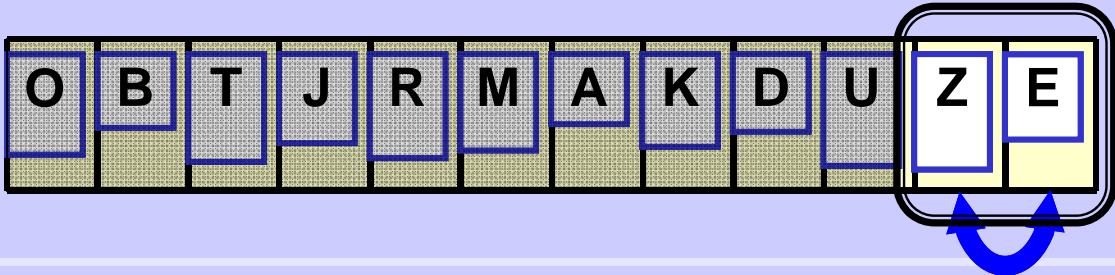


Bubble sort

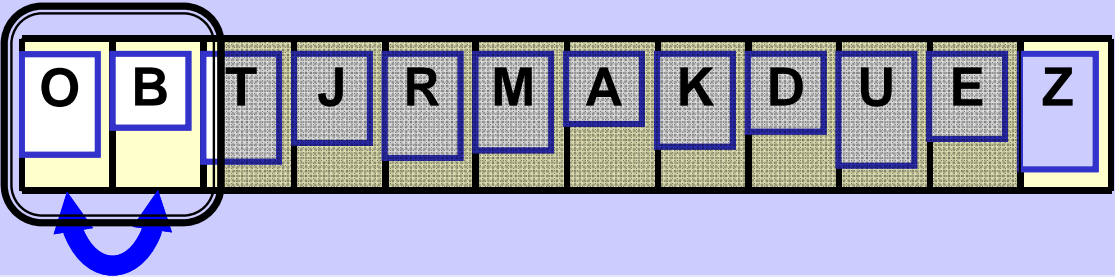
Phase 1



... etc ...

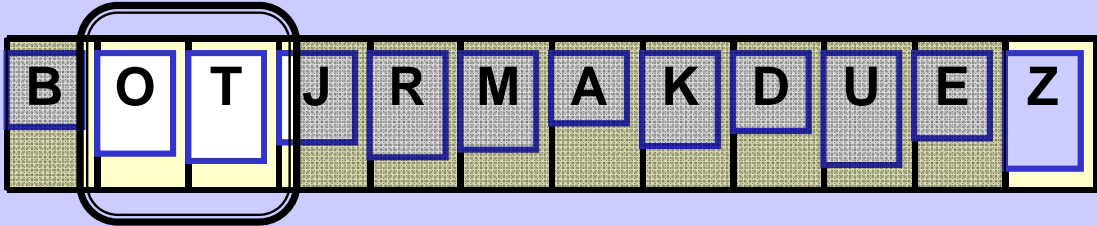


Phase 2

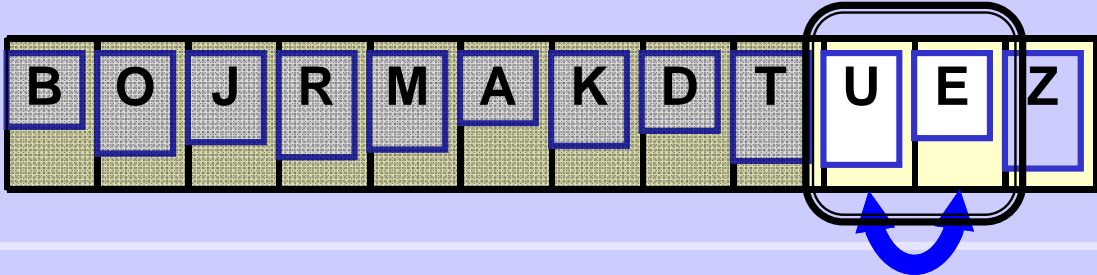


Bubble sort

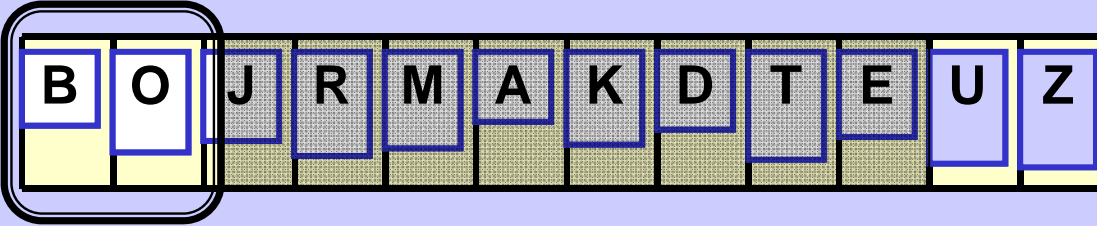
Phase 2



...etc...

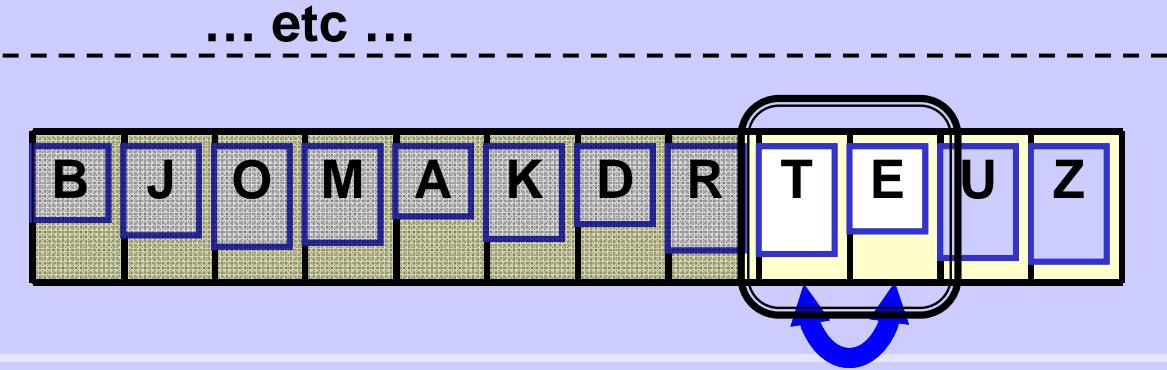


Phase 3



Bubble sort

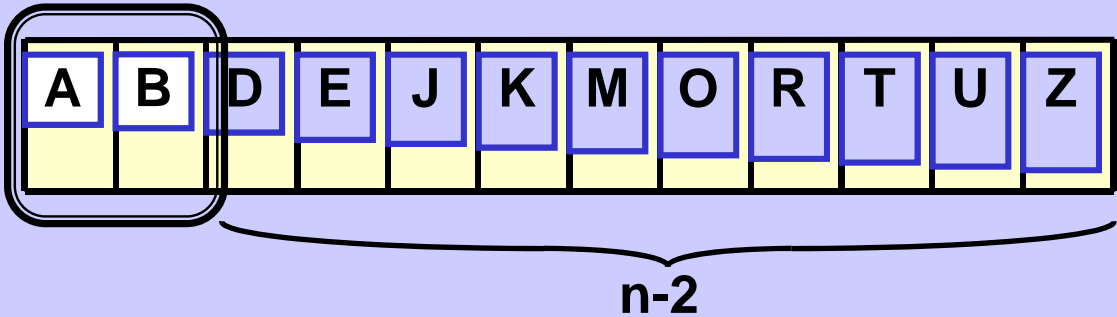
Phase 3



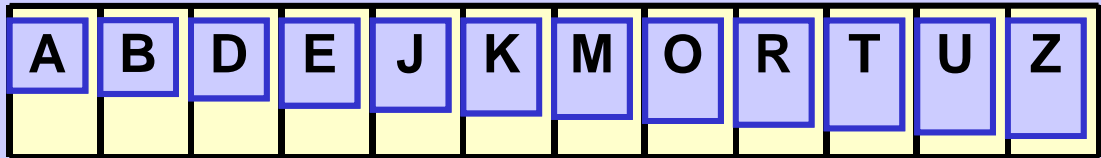
...

...

Phase n-1



All sorted



Bubble sort

```
#decreasing lastPos
for lastPos in range( len(arr)-1, -1, -1 ):
    for j in range( lastPos ):
        if arr[j] > arr[j+1]: swap( arr, j, j+1 )
```

Resume

Tests
total

$$(n-1) + (n-2) + \dots + 2 + 1 = \frac{1}{2}(n^2 - n) = \Theta(n^2)$$

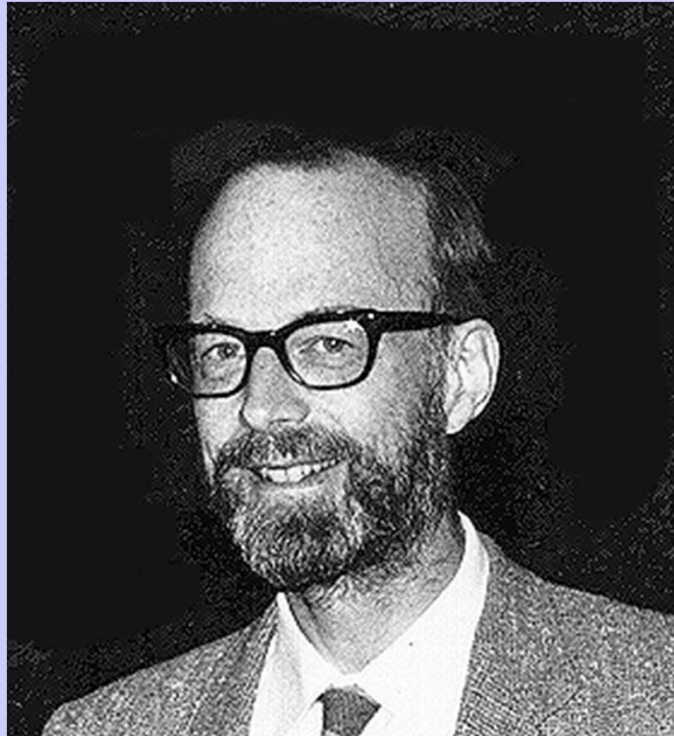
Moves
total

$$0 = \Theta(1) \quad \text{best case}$$

$$\frac{1}{2}(n^2 - n) = \Theta(n^2) \quad \text{worst case}$$

Asymptotic complexity of Bubble sort is $\Theta(n^2)$

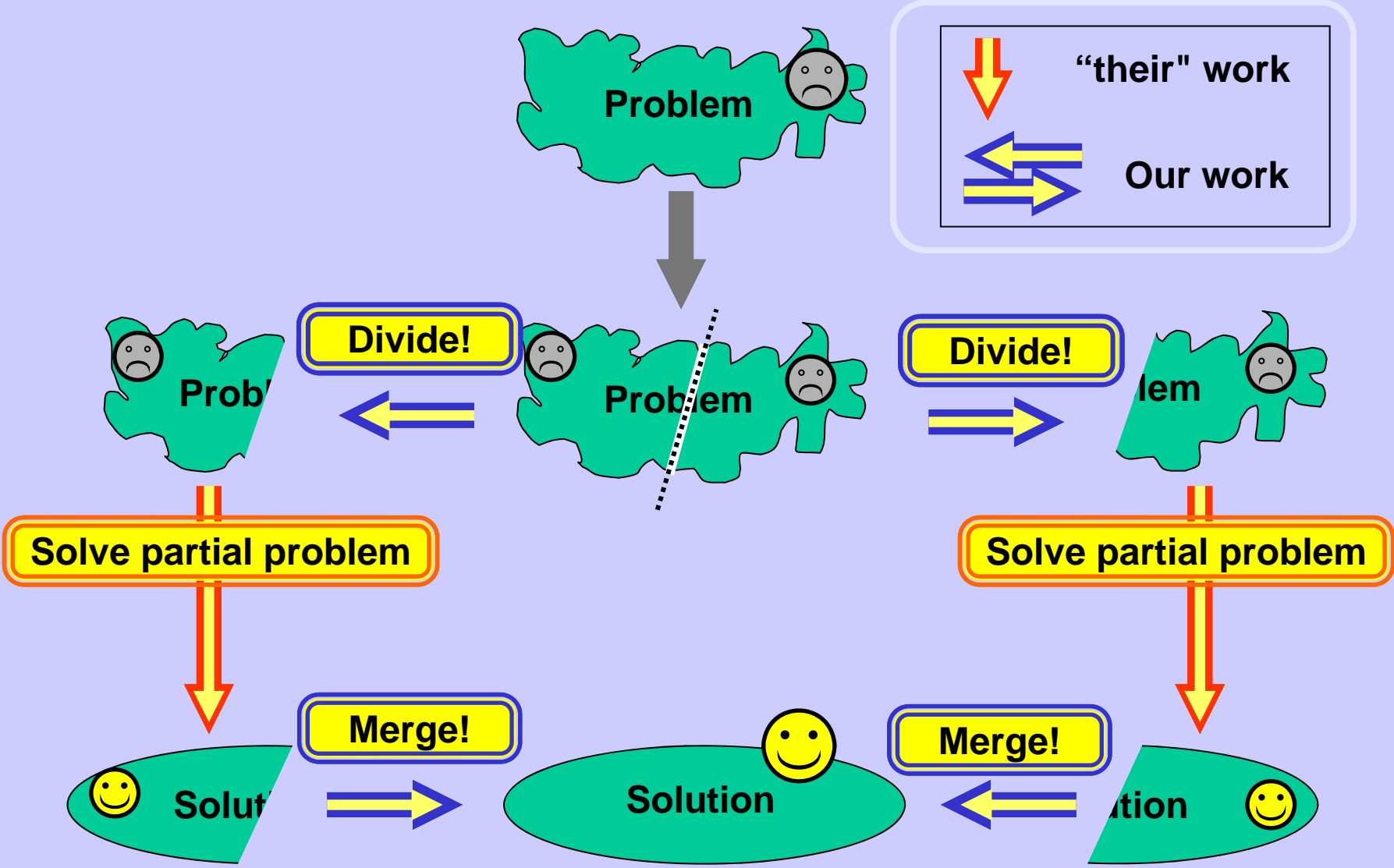
Quicksort



Sir Charles Antony Richard Hoare

C. A. R. Hoare: Quicksort. Computer Journal, Vol. 5, 1, 10-15 (1962)

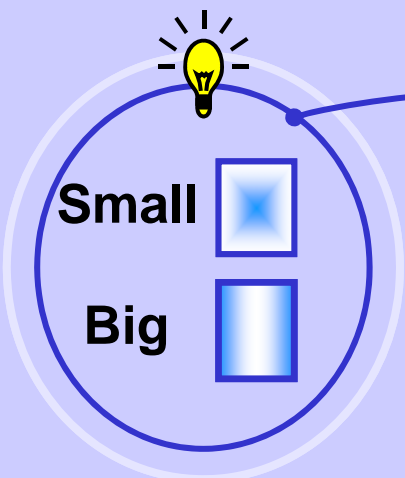
Divide and conquer! Divide et impera!



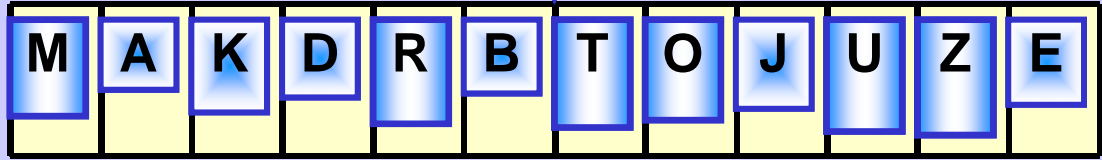
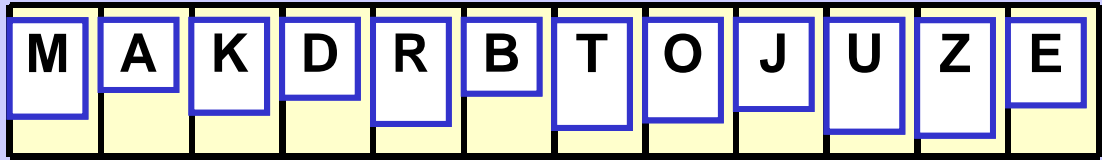
Quicksort

The idea

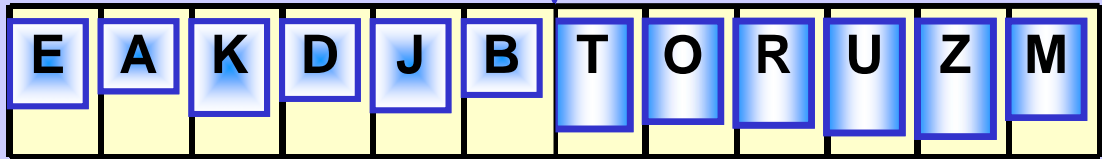
Start



Divide & Conquer!

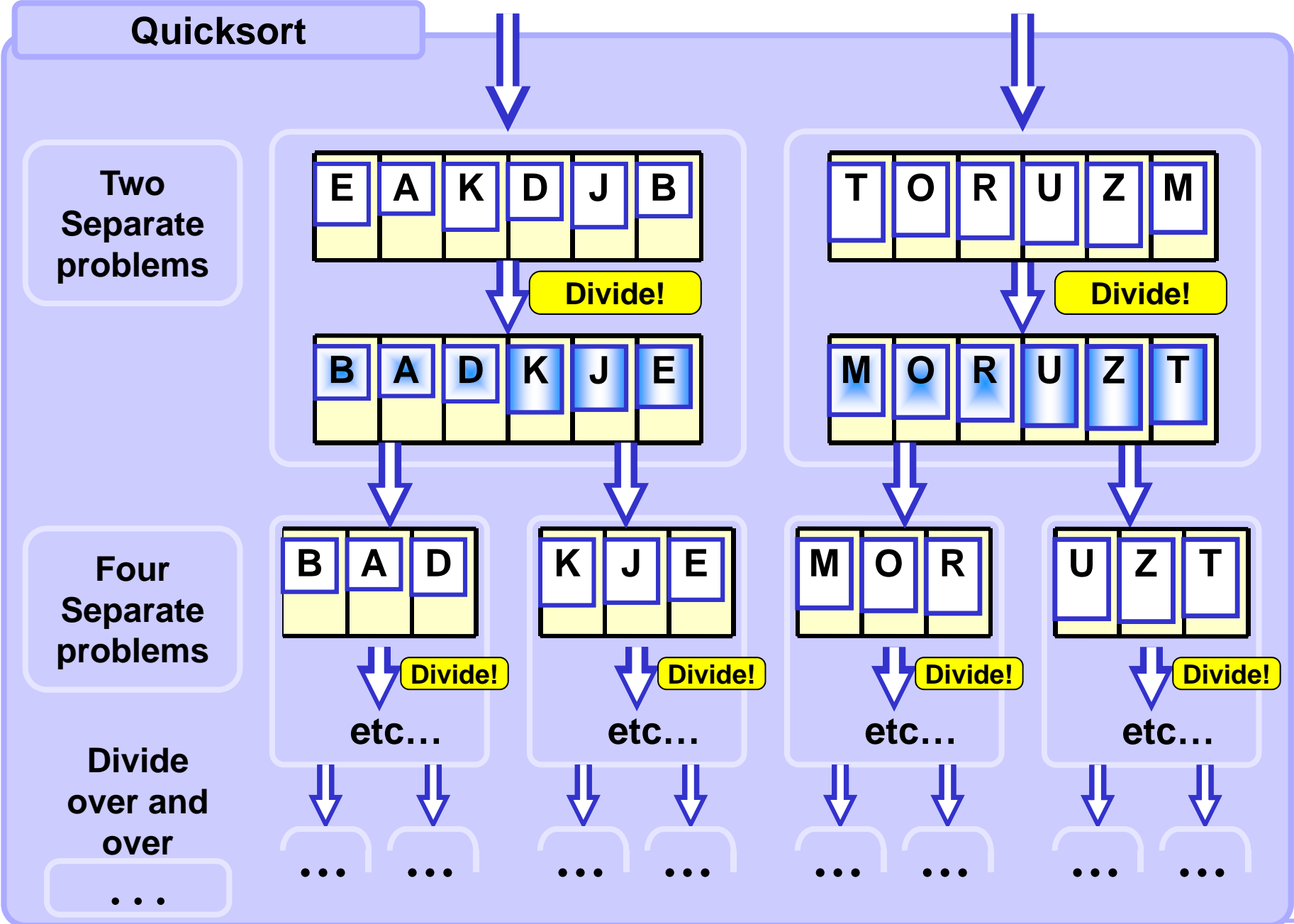


Divide!

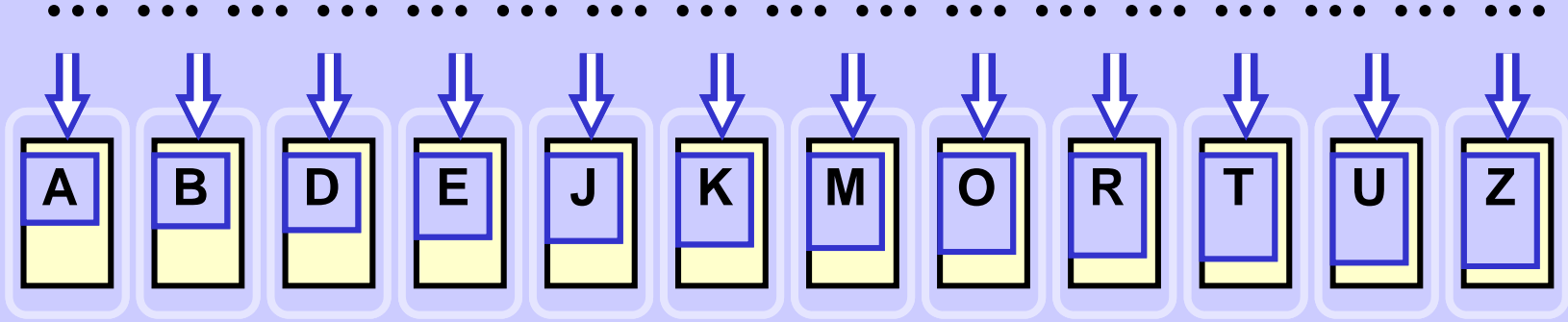


Small

Big



Quicksort

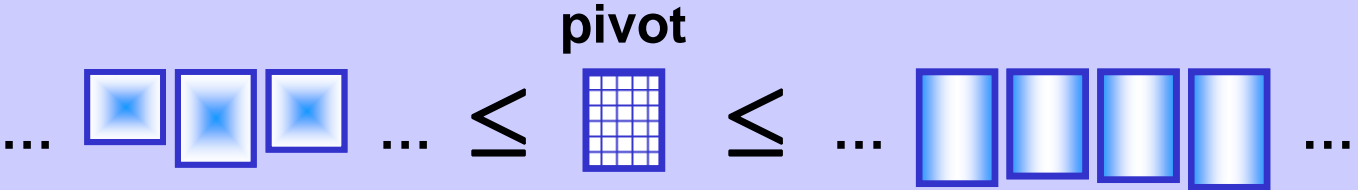


Conquered!

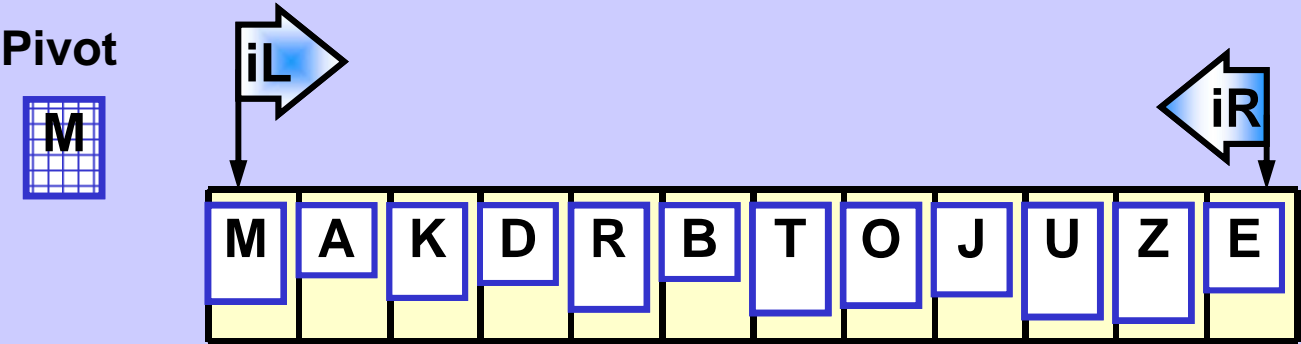
Quicksort

Dividing

Pivot



Init

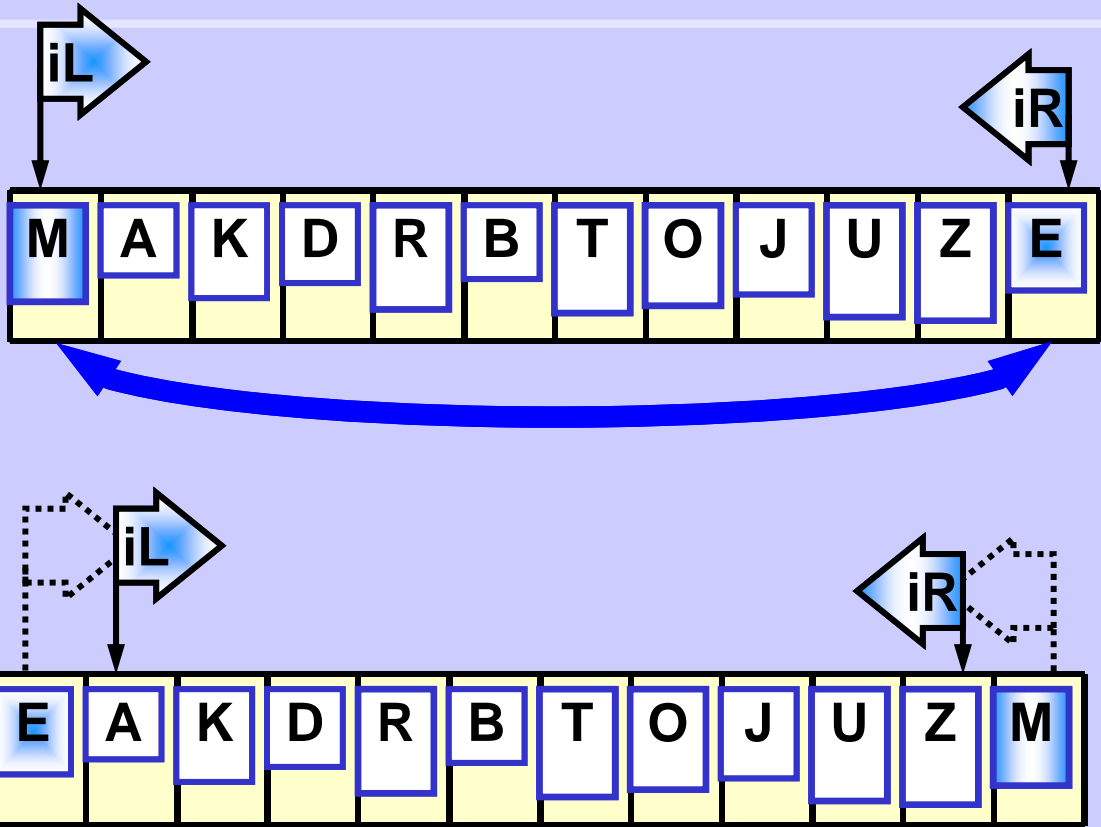


Quicksort

Dividing

Step 1

Pivot

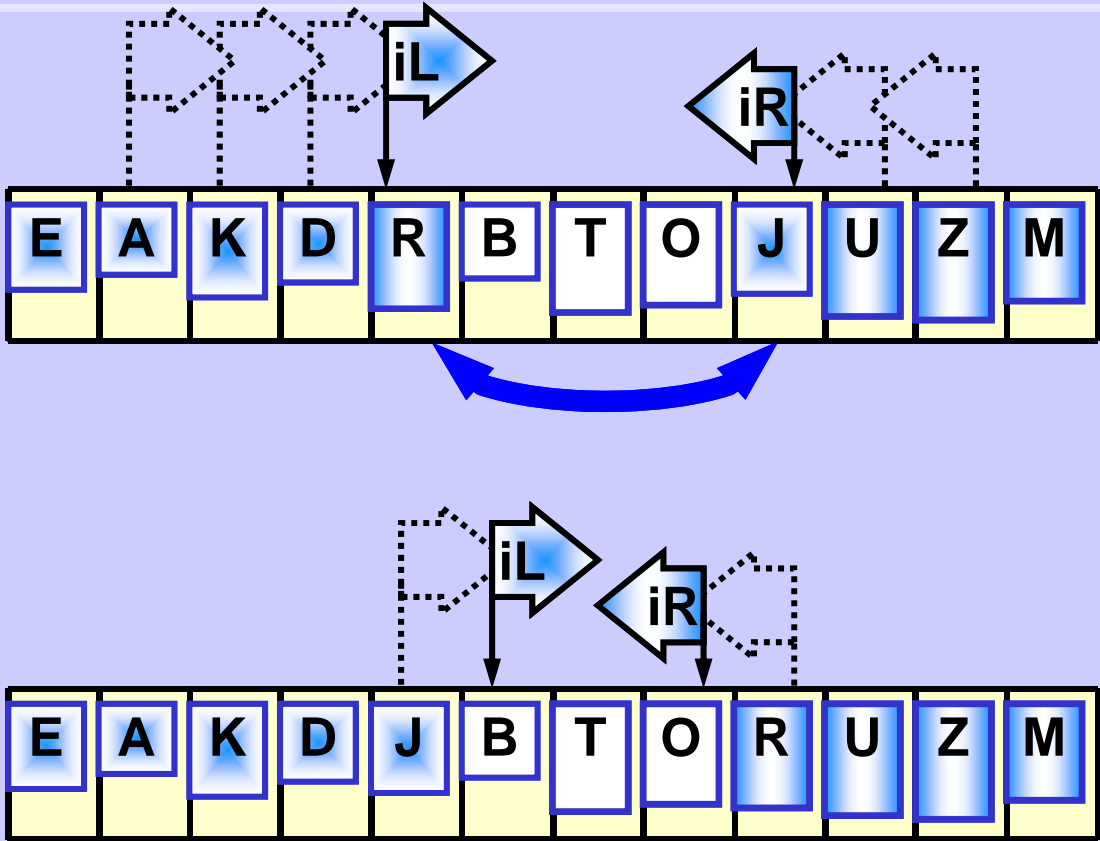


Quicksort

Dividing

Step 2

Pivot

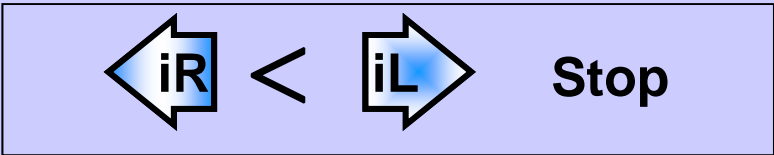
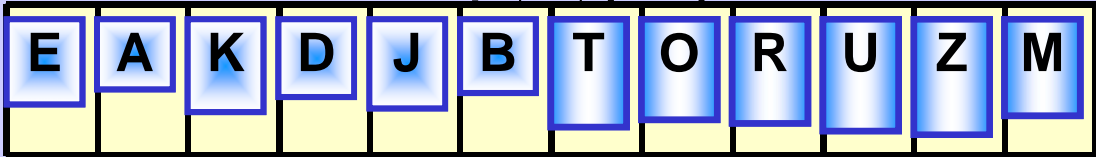


Quicksort

Dividing

Step 3

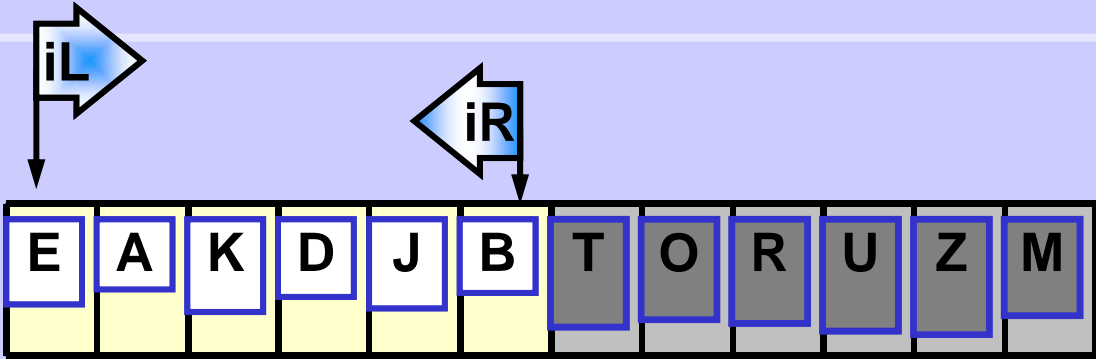
Pivot



Divide!

Init

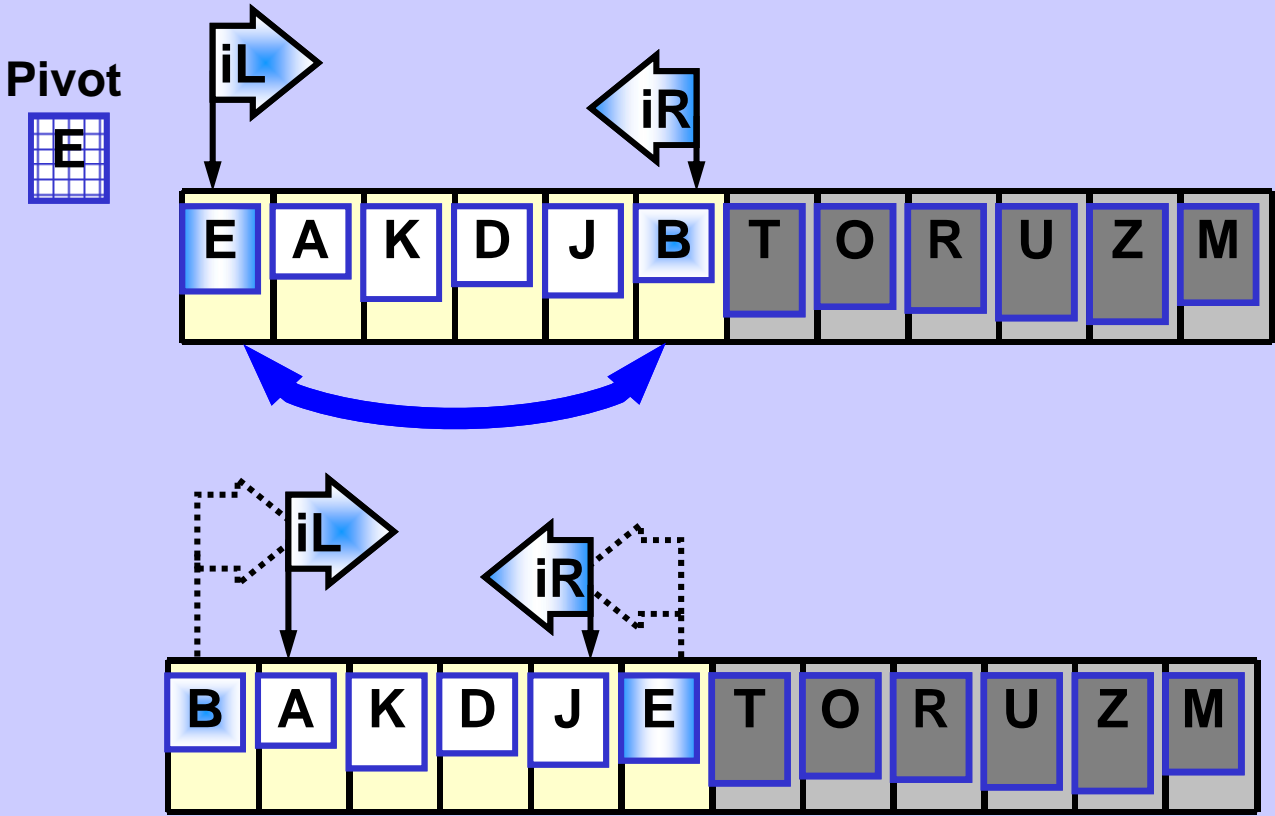
Pivot



Quicksort

Dividing

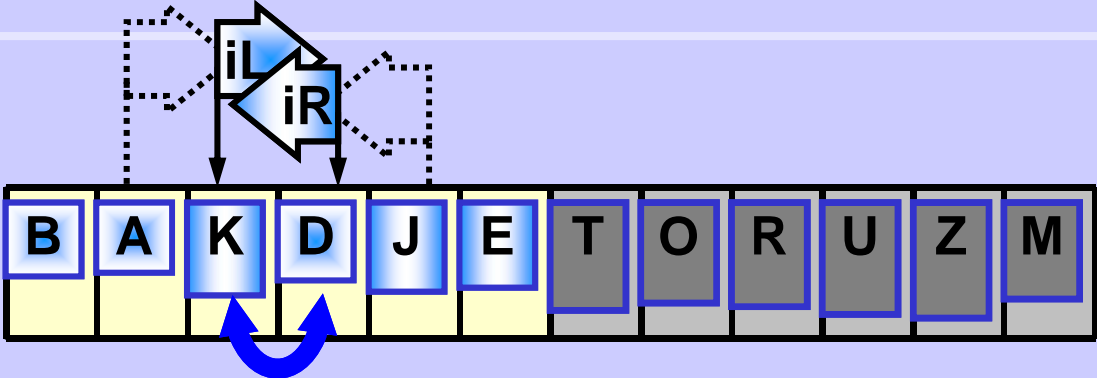
Step 1



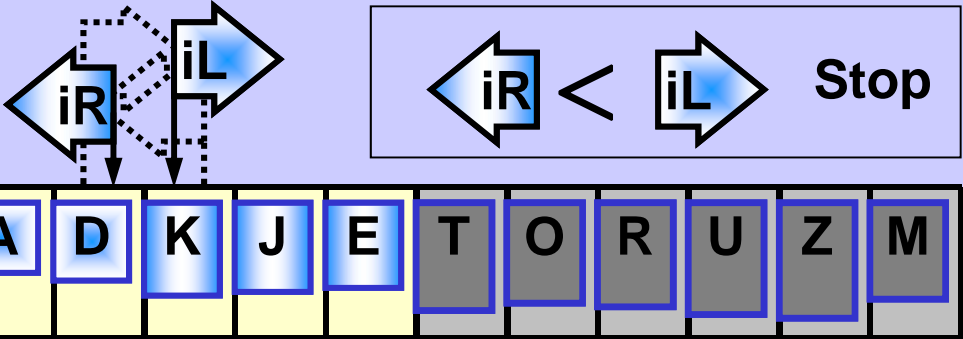
Quicksort

Dividing

Pivot

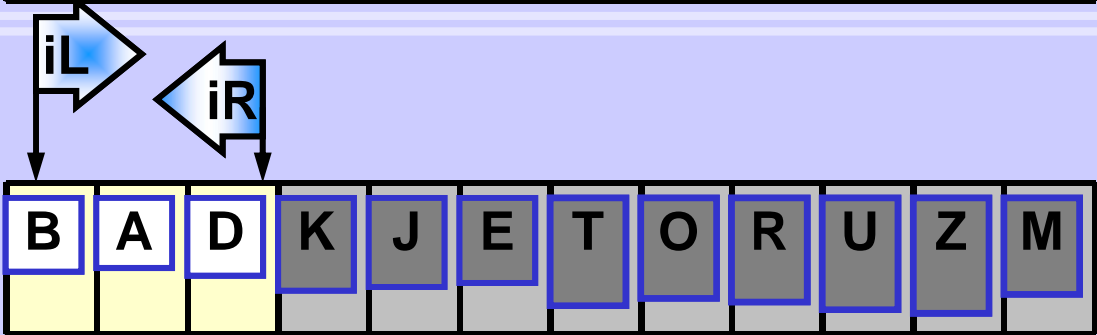


Step 2



Divide!

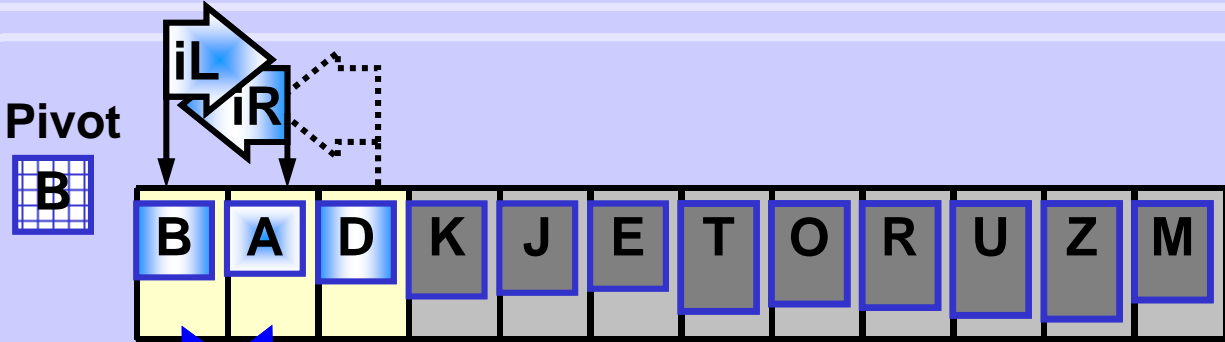
Pivot



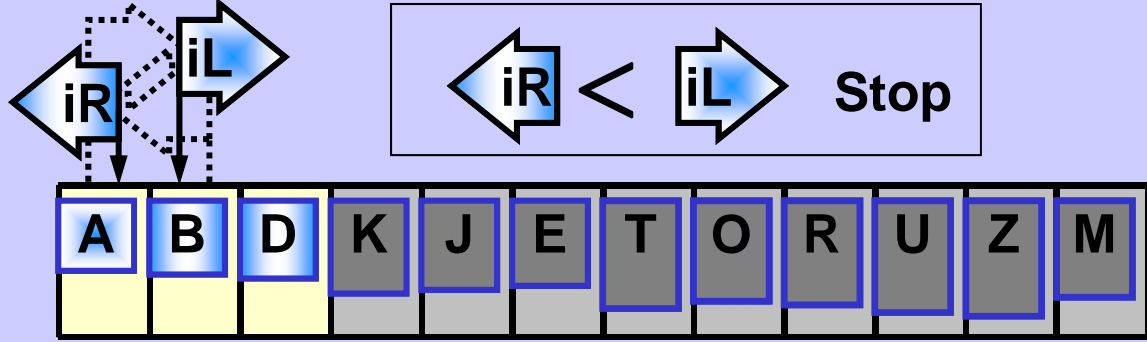
Init

Quicksort

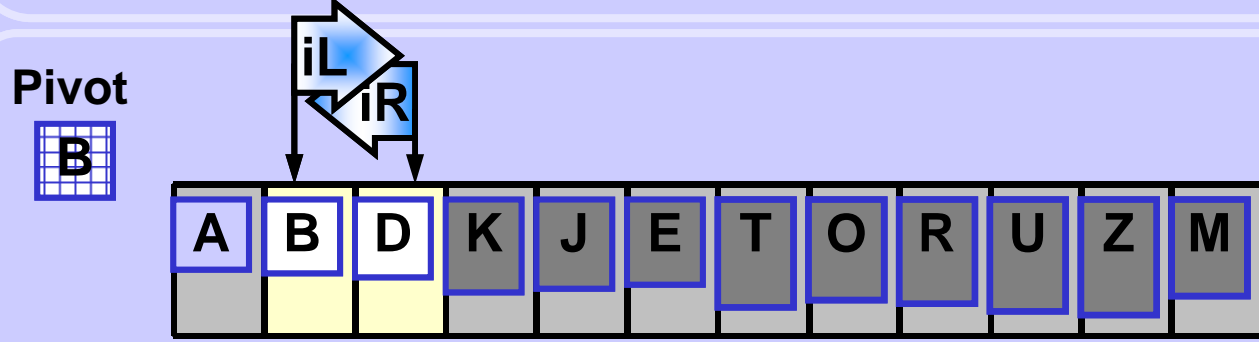
Dividing



Step 1



Divide!



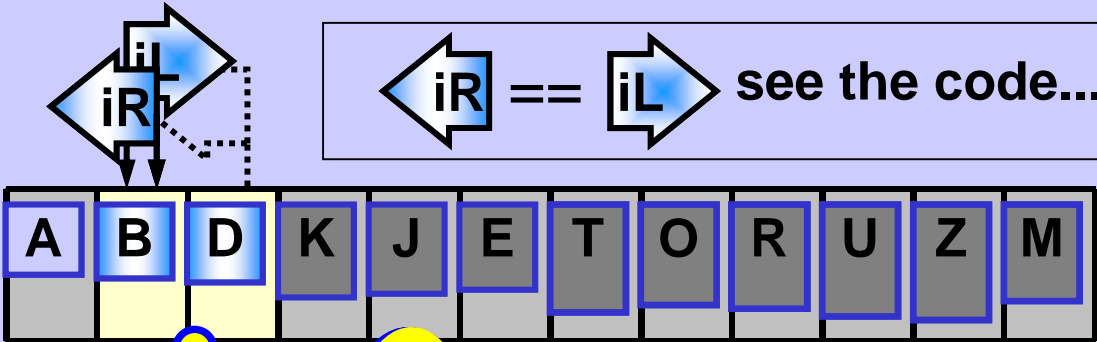
Init

Quicksort

Dividing

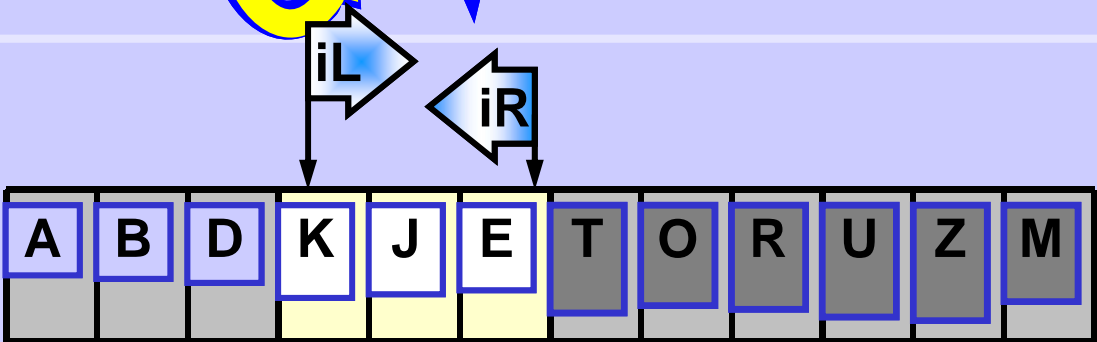
Step 1

Pivot
B



Next part

Pivot
K



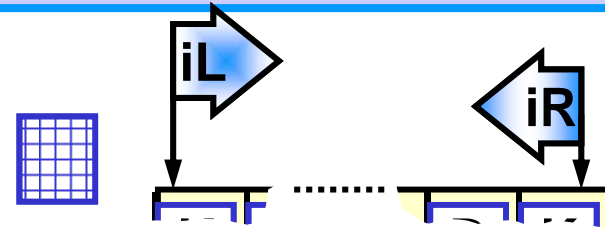
Init

etc...

etc...

Quicksort

```
def qSort( a, low, high ):
    iL = low; iR = high;
    pivot = a[low]
```



```
while True:
```

```
    if iL > iR: break
```

```
    while a[iL] < pivot: iL += 1
```

```
    while a[iR] > pivot: iR -= 1
```

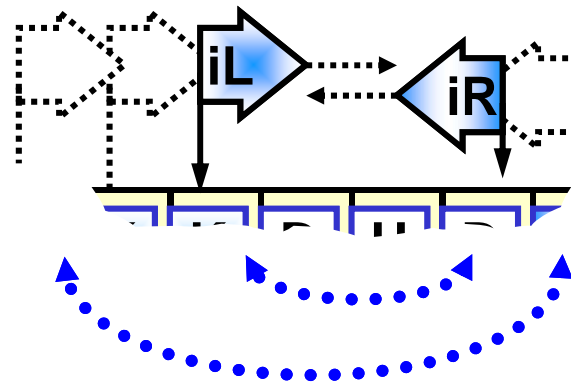
```
    if iL < iR:
```

```
        swap(a, iL, iR)
```

```
        iL += 1; iR -= 1
```

```
    else:
```

```
        if iL == iR: iL += 1; iR -= 1
```



```
if low < iR: qSort( a, low, iR )
```

```
if iL < high: qSort( a, iL, high )
```

Divide!

Quicksort

Init: Left index is set to the first element of the current segment, right index is set to its last element, a pivot value is selected.

Loop (dividing into "small" and "big") :

Left index moves to the right

and stops at element which value is greater or equal to the pivot.

Right index moves to the left

and stops at element which value is smaller or equal to the pivot.

If the left index is still to the left of the right index then

the corresponding elements are swapped

and both indices are moved by one position in their respective directions.

Else if the indices are equal then they are just moved by one in their respective directions.

The loop stops when left index is to the right of the right one.

The recursive calls follow (processing "small" and „big" separately):

Processing segment <beginning, right index>

and the segment <left index, end>

if the segment length is greater than 1.

Quicksort

Asymptotic complexity

Total
tests and moves

$$\Theta(n \cdot \log_2(n))$$

best case

$$\Theta(n \cdot \log_2(n))$$

expected case

$$\Theta(n^2)$$

worst case

Asymptotic complexity of Quicksort is $O(n^2)$...

... but! :

Expected complexity is $\Theta(n \cdot \log_2(n))$ (!!)

Quicksort



Comparing effectivity



N	N²	N × log₂(N)	$\frac{N^2}{N \times \log_2(N)}$
1	1	0	
10	100	33.2	3.0
100	10 000	6 64.4	15.1
1 000	1 000 000	9 965.8	100.3
10 000	100 000 000	132 877.1	752.6
100 000	10 000 000 000	1 660 964.0	6 020.6
1 000 000	1 000 000 000 000	19 931 568.5	50 171.7
10 000 000	100 000 000 000 000	232 534 966.6	430 042.9

Tab. 1

Alternative Partition Method

Literature (e.g. [CLRS] and web as well, e.g. Wikipedia) sometimes mentions the method implemented in the code below. The concept is straightforward, however, in practice it is slower, by factor 2 to 3.

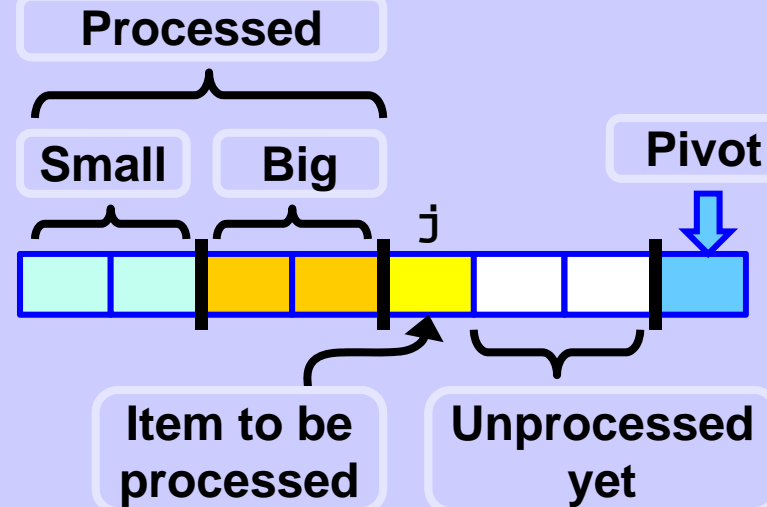
```
# returns index of the first element in the "big" part  
# including the pivot  
def partitionCLRS( a, iL, iR):  
    pivot = a[iR]; iMid = iL-1  
    for j in range(iL, iR):  
        if a[j] <= pivot:  
            iMid += 1  
            swap(a, iMid, j)  
    iMid += 1  
    swap(a, iMid, R)  
    return iMid
```

See the example
on the next slide.

Each element smaller or equal to the pivot is swapped.
About half of all entries are expected to be swapped.
The classic variant performs much less swaps.

Alternative Partition Method

Process from left to right



```

pivot = a[iR]; iMid = iL-1
for j in range(iL, iR):
    if a[j] <= pivot:
        iMid += 1
        swap(a, iMid, j)
iMid += 1
swap(a, iMid, R)
return iMid

```

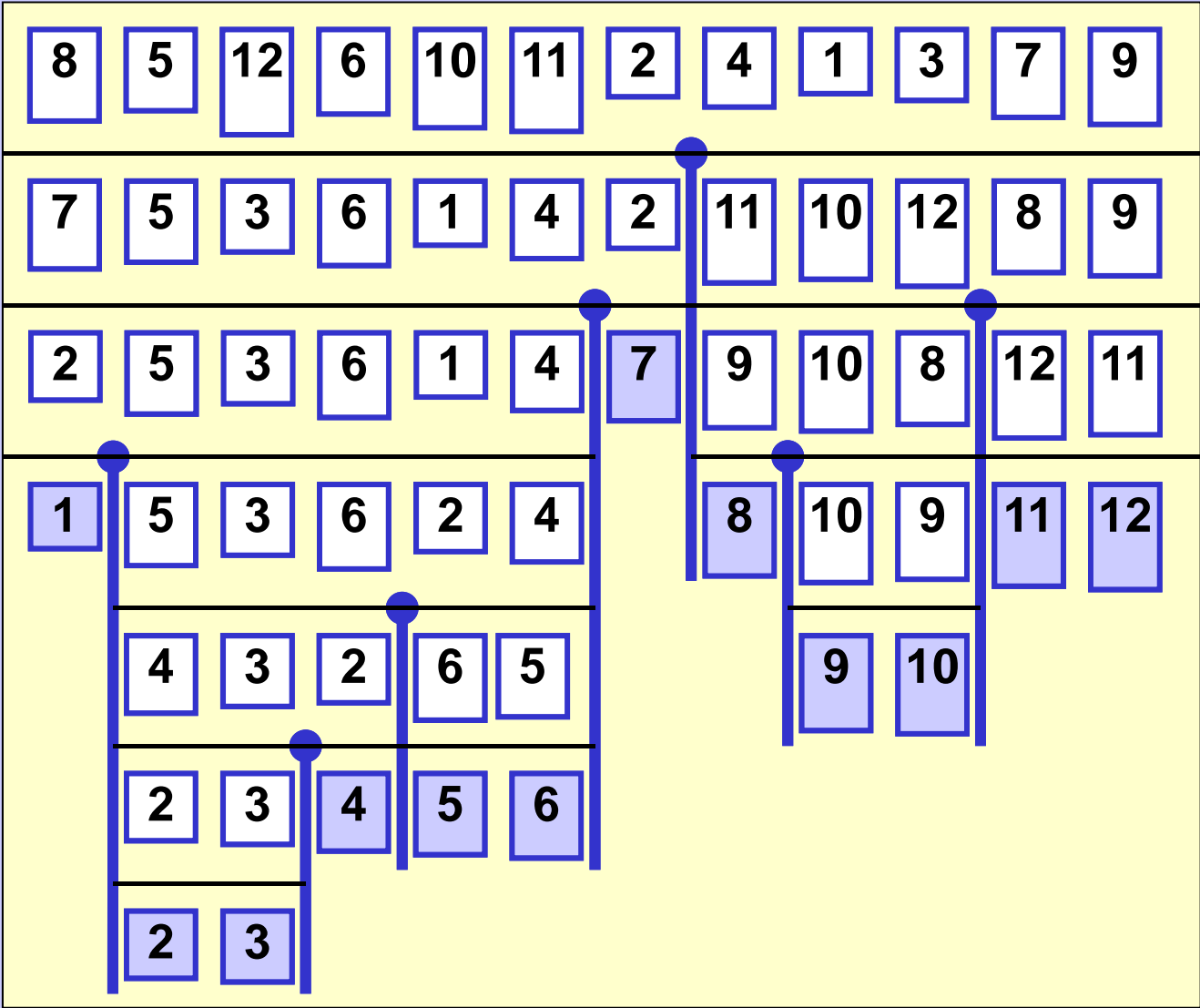
Example



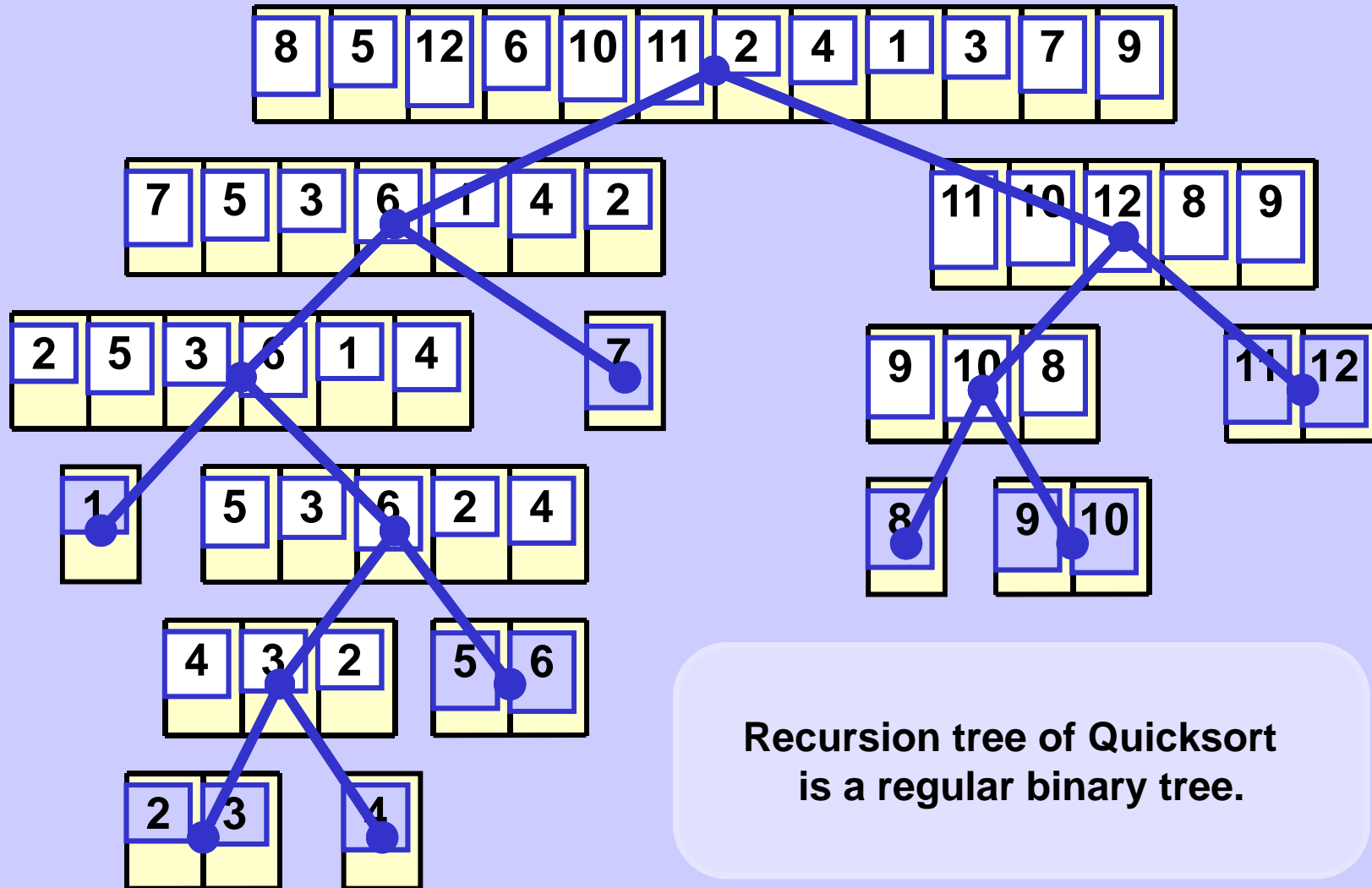
Quicksort

Example

**pivot =
= first
in the
segment**



Quicksort



Stable sort

Stable sort does not change the order of elements with the same value.

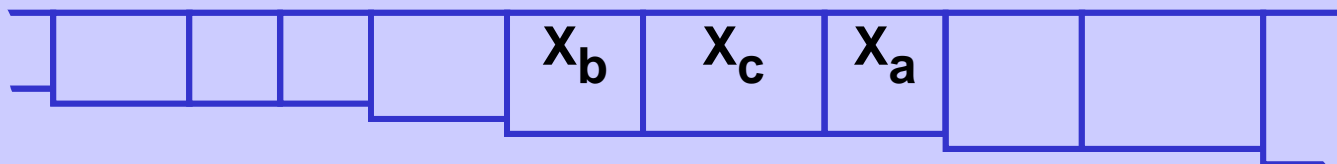
Unsorted data



Values X_i are equal



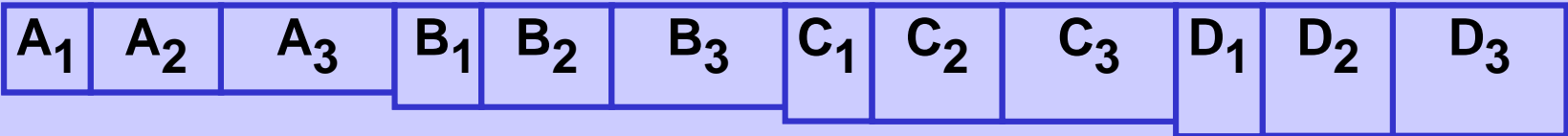
Sorted data



Sort stability



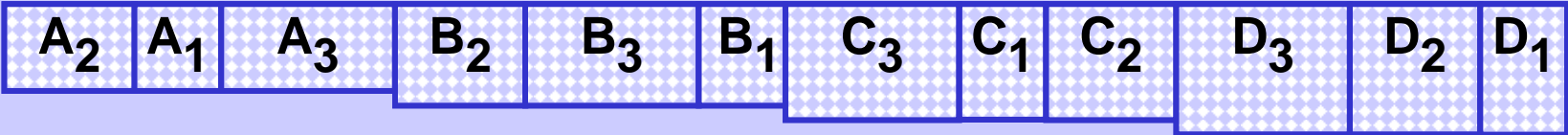
Insert Bubble -- Stable implementation



Insert Bubble -- Unstable implementation



Quicksort Always unstable!!
Select sort



Stable sort

Record:

Name	Surname
------	---------

Input: List sorted
only by names.

Andrew	Cook
Andrew	Amundsen
Andrew	Brown
Barbara	Cook
Barbara	Brown
Barbara	Amundsen
Charles	Amundsen
Charles	Cook
Charles	Brown

Stable sort
Sort records
only by

surnames

Output: List sorted
by surnames
and by names

Andrew	Amundsen
Barbara	Amundsen
Charles	Amundsen
Andrew	Brown
Barbara	Brown
Charles	Brown
Andrew	Cook
Barbara	Cook
Charles	Cook

The order of the records with the same name remains unchanged.

Unstable sort

Record:

Name	Surname
------	---------

Input: List sorted only by names.

Andrew	Cook
Andrew	Amundsen
Andrew	Brown
Barbara	Cook
Barbara	Brown
Barbara	Amundsen
Charles	Amundsen
Charles	Cook
Charles	Brown

QuickSort



Sort records only by surnames

Output: Original order of names is lost.

Sorted

Barbara	Amundsen
Andrew	Amundsen
Charles	Amundsen
Barbara	Brown
Charles	Brown
Andrew	Brown
Charles	Cook
Andrew	Cook
Barbara	Cook

The order of the records with the same name is changed.