Description Logics – Querying

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Outline

1. What if OWL is not enough?
2. Complex Queries
    - Evaluation of Conjunctive Queries in $\mathcal{ALC}$
3. Modeling Error Explanation
    - Black-box methods
    - Algorithms based on CS-trees
    - Algorithm based on Reiter's Algorithm
    - Algorithm based on Reiter's Algorithm

Problems

- OWL cannot express everything, can we do more?
- Is there any more powerful “query” language (beyond consistency checking)?
- What to do if an ontology is inconsistent?

What if OWL is not enough?

SROIQ (OWL) Revision

$$Man \sqsubseteq Person$$

$$Man \sqsubseteq \neg Woman$$

$$Man \sqcap \exists hasChild \cdot Man \sqsubseteq FatherOfSons$$

$$hasSon \sqsubseteq hasChild$$

$$hasParent \circ hasBrother \sqsubseteq hasUncle$$

$$trans(hasDescendant)$$

$$sym(hasSpouse)$$

$$fun(hasMother)$$

$$hasWife \sqsubseteq hasHusband^{-}$$

How to express hasStepSibling?

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$$\begin{aligned} hasSpouse(?m1, ?f), hasSpouse(?m2, ?f),\\ hasChild(?m1, ?c1), hasChild(?m2, ?c2),\\ hasChild(?f, ?c1), hasChild(?f, ?c2), ?c1! =?c2\\ \rightarrow hasStepSibling(?c1, ?c2) \end{aligned}$$

OWL2-DL + rules undecidable

… unless variables in rules are restricted to match named individuals only.

DL-safe Rules

A rule is DL-safe, if its variables are *distinguished*, i.e. thet can only match **named individuals** in the ontology. Consistency checking of OWL2-DL + DL-safe rules is decidable.

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Complex Queries

What if we need to answer a complex query?

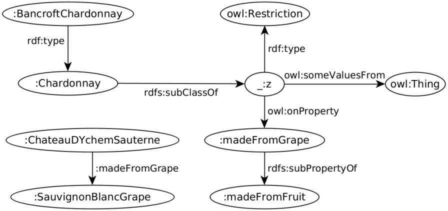
How many czech writers died in the Czech Republic according to DBPedia ?

```
PREFIX dbo: <http://dbpedia.org/ontology/>
PREFIX dbr: <http://dbpedia.org/resource/>
PREFIX dcterms: <http://purl.org/dc/terms/>
SELECT COUNT(?x)
{
  ?x   dbo:deathPlace  dbr:Czech_Republic ;
       dcterms:subject dbr:Category:Czech_writers .
}
```

at the following endpoint:

`http://dbpedia-live.openlinksw.com/sparql/`

To remind – SPARQL Evaluation Semantics



```
PREFIX : <http://ex.org/e1>
SELECT ?x
WHERE { ?x :madeFromFruit _:y }
```

Simple-entailment No result.

$$Chardonay(BancroftChardonnay).$$
$$Chardonay \sqsubseteq \exists madeFromGrape \cdot \top$$
$$madeFromGrape \sqsubseteq madeFromFruit$$

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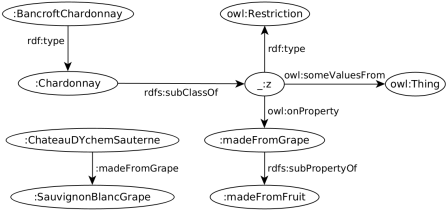

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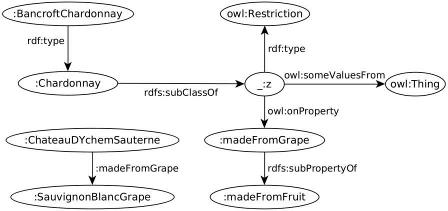

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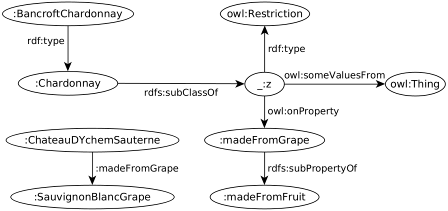

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RDF-entailment No result.

RDFS-entailment One result: `?x=:ChateauDYchemSauterne`.

OWL-entailment Two results: `?x=:ChateauDYchemSauterne` and `?x=:BancroftChardonnay`.

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$$Chardonay \sqsubseteq \exists madeFromGrape \cdot \top$$
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Conjunctive Queries

Metaqueries

Query Types

Conjunctive (ABox) queries – queries asking for individual tuples complying with a graph-like pattern.

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Example

"Find all mothers and their daughters having at least one brother." :

$$\begin{aligned} Q(?x, ?z) \quad \leftarrow \quad & Woman(?x), hasChild(?x, ?y), hasChild(?x, ?z), \\ & Man(?y), Woman(?z) \end{aligned}$$

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Metaqueries – queries asking for individual/concept/role tuples. There are several languages for metaqueries, e.g. SPARQL-DL, OWL-SAIQL, etc.

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Example

"Find all people together with their type." in SPARQL-DL:

$$Q(?x,?c) \leftarrow TYPE(?x,?c), SUBCLASSOF(?c, Person)$$

Conjunctive (ABox) queries

Conjunctive (ABox) queries are analogous to database SELECT-PROJECT-JOIN queries.

Conjunctive Query

$$Q(?x_1, \ldots, ?x_D) \leftarrow t_1, \ldots t_T,$$

where each t_i is either

- $C(y_k)$ (where C is a concept)
- $R(y_k, y_l)$ (where R is a role)

and y_i is either (i) an individual, or (ii) variable from a new set V (variables will be differentiated from individuals by the prefix "?"). We need all $?x_i$ to be present also in one of t_i.

Conjunctive ABox Queries – Semantics

- Conjunctive queries of the form $Q()$ are called *boolean* – such queries only test existence of a relational structure in each model $\mathcal{I}$ of the ontology $\mathcal{K}$.
- Consider any interpretation $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$. *Evaluation* η is a function from the set of individuals and variables into $\Delta^{\mathcal{I}}$ that coincides with $\mathcal{I}$ on individuals.
- Then $\mathcal{I} \models_{\eta} Q()$, iff
  - $\eta(y_k) \in C^{\mathcal{I}}$ for each atom $C(y_k)$ from $Q()$ and
  - $\langle \eta(y_k), \eta(y_l) \rangle \in R^{\mathcal{I}}$ for each atom $R(y_k, y_l)$ from $Q()$
- Interpretation $\mathcal{I}$ is a model of $Q()$, iff $\mathcal{I} \models_{\eta} Q()$ for some η.
- Next, $\mathcal{K} \models Q()$ ($Q()$ is satisfiable in $\mathcal{K}$) iff $\mathcal{I} \models Q()$ whenever $\mathcal{I} \models \mathcal{K}$

Conjunctive ABox Queries – Variables

- Queries without variables are not practically interesting. For queries with variables we define semantics as follows. An N-tuple $\langle i_1, \ldots, i_n \rangle$ is a *solution* to $Q(?x_1, \ldots, ?x_n)$ in theory $\mathcal{K}$, whenever $\mathcal{K} \models Q'()$, for a boolean query Q' obtained from Q by replacing all occurences of $?x_1$ in all t_k by an individual i_1, etc.

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distinguished occur in the query head as well as body, e.g. $?x, ?z$ in the previous example. These variables are evaluated as domain elements that are necessarily interpretations of some individual from $\mathcal{K}$. That individual is the binding to the distinguished variable in the query result.

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undistinguished occur only in the query body, e.g. $?y$ in the previous example. Their can be interpretated as any domain elements.

Conjunctive Queries – Examples

Example

Let's have a theory $\mathcal{K}_4 = (\emptyset, \{(\exists R_1 \cdot C_1)(i_1), R_2(i_1, i_2), C_2(i_2)\})$.

- Does $\mathcal{K} \models Q_1()$ hold for $Q_1() \leftarrow R_1(?x_1, ?x_2)$?
- What are the solutions of the query $Q_2(?x_1) \leftarrow R_1(?x_1, ?x_2)$ for $\mathcal{K}$?
- What are the solutions of the query $Q_3(?x_1, ?x_2) \leftarrow R_1(?x_1, ?x_2)$ for $\mathcal{K}$?

Evaluation of Conjunctive Queries in $\mathcal{ALC}$

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Satisfiability of $\mathcal{ALC}$ Boolean Queries

- Satisfiability of the boolean query $Q()$ having a tree shape can be checked by means of the **rolling-up technique**. $\Longrightarrow$

Rolling-up Technique

- Each two atoms $C_1(y_k)$ and $C_2(y_k)$ can be replaced by a single query atom of the form $(C_1 \sqcap C_2)(y_k)$.

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  - (i) C, whenever y_l is a variable,
  - (ii)$C \sqcap Y_l$, whenever y_l is an individual. Y_l is a *representative concept* of individual y_l occuring neither in $\mathcal{K}$ nor in Q. For each y_l it is necessary to extend ABox of $\mathcal{K}$ with concept assertion $Y_l(y_l)$.

Satisfiability of $\mathcal{ALC}$ Boolean Queries (2)

... after rolling-up the query we obtain the query $Q()' \leftarrow C(y)$, that is satisfied in $\mathcal{K}$, iff $Q()$ is satisfied in $\mathcal{K}$:

- **If y is an individual, then $Q'()$ is satisfied, whenever $\mathcal{K} \models C(y)$ (i.e. $\mathcal{K} \cup \{(\neg C)(y)\}$ is inconsistent)**

Example

Consider a query $Q_4() \leftarrow R_1(?x_1, ?x_2), R_2(?x_1, ?x_3), C_2(?x_3)$. This query can be rolled-up into the query $Q_4' \leftarrow (\exists R_1 \cdot \top \sqcap \exists R_2 \cdot C_2)(?x_1)$. This query is satisfiable in $\mathcal{K}_4$, as $\mathcal{K}_4 \cup \{(\exists R_1 \cdot \top \sqcap \exists R_2 \cdot C_2) \sqsubseteq \bot\}$ is inconsistent.

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- **If y is a variable, then $Q'()$ is satisfied, whenever $\mathcal{K} \cup \{C \sqsubseteq \bot\}$ is inconsistent. Why ?**

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Satisfiability of Boolean Queries in $\mathcal{ALC}$ (3)

… and what to do with queries with distinguished variables ?

- Let's consider just queries that form "connected component" and contain for some variable y_k at least two query atoms of the form $R_1(y_1, y_k)$ and $R_2(y_2, y_k)$.

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- Question: *Why is it enough to take just one connected component?*
- **Let's make use of the tree model property of $\mathcal{ALC}$. Each pair of atoms $R_1(y_1, y_k)$ and $R_2(y_2, y_k)$ can be satisfied only if y_k is interpreted as a domain element, that is an interpretation of an individual – y_k can be treated as distinguished. Why (see next slide) ?**

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- For $\mathcal{SHOIN}$ and $\mathcal{SROIQ}$ there is no sound and complete decision procedure for general boolean queries.

$\mathcal{ALC}$ Model Example

Queries with Distinguished Variables – naive pruning

Consider arbitrary query $Q(?x_1, \ldots, ?x_D)$. How to evaluate it ?

- **naive way**: Replace each distinguished variable x_i with each individual occuring in $\mathcal{K}$. *Solutions* are those D-tuples $\langle i_1, \ldots, i_D \rangle$, for which a boolean query created from Q by replacing each x_k with i_k is satisfiable.

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Remind that $\mathcal{K}_4 = (\emptyset, \{(\exists R_1 \cdot C_1)(i_1), R_2(i_1, i_2), C_2(i_2)\})$. The query

$$Q_5(?x_1) \leftarrow R_1(?x_1, ?x_2), R_2(?x_1, ?x_3), C_2(?x_3)$$

has *solution* $\langle i_1 \rangle$ as

$$Q_5'() \leftarrow R_1(i_1, ?x_2), R_2(i_1, ?x_3), C_2(?x_3)$$

can be rolled into $Q_5''()$ for which $\mathcal{K}_4 \models Q_5''$:

$$Q_5''() \leftarrow (\exists R_1 \cdot \top \sqcap \exists R_2 \cdot C_2)(i_1)$$

Queries with Distinguished Variables – naive pruning

… another example

The query

$$Q_6(?x_1, ?x_3) \leftarrow R_1(?x_1, ?x_2), R_2(?x_1, ?x_3), C_2(?x_3)$$

has *solution* $\langle i_1, i_2 \rangle$ as

$$Q_6'() \leftarrow R_1(i_1, ?x_2), R_2(i_1, i_2), C_2(i_2)$$

can be rolled into Q_6'' for which $\mathcal{K}_4 \cup \{\mathbf{I_2(i_2)}\} \models Q_6''$.

$$Q_6''() \leftarrow (\exists R_1 \cdot \top \sqcap \exists R_2 \cdot (C_2 \sqcap I_2))(i_1).$$

Similarly $Q_7(?x_1, ?x_2) \leftarrow R_1(?x_1, ?x_2), R_2(?x_1, ?x_3), C_2(?x_3)$ has no solution.

Queries with Distinguished Variables – iterative pruning

- **… a bit more clever strategy than replacing all variables**: First, let's replace just the first variable $?x_1$ with each individual from $\mathcal{K}$, resulting in Q_2. If the subquery of Q_2 containing all query atoms from Q_2 without distinguished variables is not a logical consequence of $\mathcal{K}$, then we do not need to test potential bindings for other variables.

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- Many other optimizations are available.

Queries with Distinguished Variables – iterative pruning

For the query $Q_6(?x_1, ?x_3)$, the naive strategy needs to check four different bindings (resulting in four tableau algorithm runs)

$$\begin{gathered}\langle i_1, i_1\rangle,\\ \langle \mathbf{i_1}, \mathbf{i_2}\rangle,\\ \langle i_2, i_1\rangle,\\ \langle i_2, i_2\rangle.\end{gathered}$$

Out of them only $\langle i_1, i_2\rangle$ is a solution for Q_6. Consider only partial binding $\langle i_2\rangle$ for $?x_1$. Applying this binding to Q_6 we get $Q_7(?x_3) = R_1(i_2, ?x_2), R_2(i_2, ?x_3), C_2(?x_3)$. Its distinguished-variable-free subquery is $Q_7'() = R_1(i_2, ?x_2)$ and $\mathcal{K}_4 \not\models Q_7'$. Because of **monotonicity** of $\mathcal{ALC}$, we do not need to check the two bindings for $?x_3$ in this case which saves us one tableau algorithm run.

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Modeling Error Explanation

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- We can start iterating through all axioms in the theory and look, “what went wrong”.
- ... but hardly in case we have **hundred thousand axioms**
- A solution might be to ask the computer to *localize the axioms causing the problem for us*.

DNA

images/l11/dna2e_color.png

MUPS – example

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Example

Consider theory $\mathcal{K}_5 = (\{\alpha_1, \alpha_2, \alpha_3\}, \emptyset)$

$$\begin{aligned} \alpha_1 &: \mathit{Person} \sqsubseteq \exists \mathit{hasParent} \cdot (\mathit{Man} \sqcap \mathit{Woman}) \sqcap \forall \mathit{hasParent} \cdot \neg \mathit{Person}, \\ \alpha_2 &: \mathit{Man} \sqsubseteq \neg \mathit{Woman}, \\ \alpha_3 &: \mathit{Man} \sqcup \mathit{Woman} \sqsubseteq \mathit{Person}. \end{aligned}$$

Unsatisfiability of *Person* comes independently from two axiom sets (MUPSes), namely $\{\alpha_1, \alpha_2\}$ and $\{\alpha_1, \alpha_3\}$. Check it yourself !

MUPS

Currently two approaches exist for searching all MUPSes for given concept:

black-box methods perform many satisfiability tests using existing inference engine.

- ☺ flexible and easily reusable for another (description) logic
- ☹ time consuming

glass-box methods all integrated into an existing reasoning (typically tableau) algorithm.

- ☺ efficient
- ☹ hardly reusable for another (description) logic.

Glass-box methods

- For $\mathcal{ALC}$ there exists a complete algorithm with the following idea:
  - tableau algorithm for $\mathcal{ALC}$ is extended in such way that it “remembers which axioms were used during completion graph construction”.
  - for each completion graph containing a clash, the axioms that were used during its construction can be transformed into a MUPS.
- Unfortunately, complete glass-box methods do not exist for OWL-DL and OWL2-DL. The same idea (tracking axioms used during completion graph construction) can be used also for these logics, but only as a preprocessing reducing the set of axioms used by a black-box algorithm.

Black-box methods

Task formulation

- Let's have *a set of axioms* X of given DL and *reasoner* R for given DL. We want to find MUPSes for :
  1. concept unsatisfiability, '
  2. theory (ontology) inconsistency,
  3. arbitrary entailment.
- It can be shown (see [**k2006droo**]) that w.l.o.g. we can deal only with *concept unsatisfiability*.
- **MUPS:** Let's denote $MUPS(C, Y)$ a minimal subset $MUPS(C, Y) \subseteq Y \subseteq X$ causing unsatisfiability of C.
- **Diagnose:** Let's denote $DIAG(C, Y)$ a minimal subset $DIAG(C, Y) \subseteq Y \subseteq X$, such that if $DIAG(C, Y)$ is removed from Y, the concept C becomes satisfiable.

Task formulation (2)

- Let's focus on concept C unsatisfiability. Denote

$$R(C,Y) = \left\{ \begin{array}{ll} \mathit{true} & \text{iff} Y \not\models (C \sqsubseteq \bot) \\ \mathit{false} & \text{iff} Y \models (C \sqsubseteq \bot)) \end{array} \right\}.$$

- There are many methods (see [**bsw2003famus**]). We introduce just two of them:
  - Algorithms based on CS-trees.
  - Algorithm for computing a single MUPS[**k2006droo**] + Reiter algorithm [**r1987tdfp**].

Algorithms based on CS-trees

CS-trees

- A naive solution: test for each set of axioms from $\mathcal{T} \cup \mathcal{A}$ for $\mathcal{K} = (\mathcal{T}, \mathcal{A})$, whether the set causes unsatisfiability – minimal sets of this form are MUPSes.
- *Conflict-set trees (CS-trees)* systematize exploration of all these subsets of $\mathcal{T} \cup \mathcal{A}$. The main gist :

  > *If we found a set of axioms X that do not cause unsatisfiability of C (i.e. $X \not\models C \sqsubseteq \bot$), then we know (and thus can avoid asking reasoner) that $Y \not\models C \sqsubseteq \bot$ for each $Y \subseteq X$.*

- CS-tree is a representation of the state space, where each state s has the form (D, P), where
  - D is a set of axioms that *necessarily has to be part of all MUPSes* found while exploring the subtree of s.
  - P is a set of axioms that *might be part of some MUPSes* found while exploring the subtree of s.

CS-tree Exploration – Example

Example

A CS-tree for unsatisfiability of *Person* (abbr. *Pe*, not to be mixed with the set P) in $\mathcal{K}_5 = \{\alpha_1, \alpha_2, \alpha_3\}$:

$$\underbrace{Pe \sqsubseteq \exists hP \cdot (M \sqcap W) \sqcap \forall hP \cdot \neg Pe}_{\alpha_1}, \quad \underbrace{M \sqsubseteq \neg W}_{\alpha_2}, \quad \underbrace{M \sqcup W \sqsubseteq Pe}_{\alpha_3}.$$

`images/l11/cstree.png`

In gray states, the concept *Person* is satisfiable ($R(Pe, D \cup P) = true$). States with a dotted border are pruned by the algorithm.

CS-tree Exploration

The following algorithm is exponential in the number of tableau algorithm runs.

1. (Init) The root of the tree is an initial state $s_0 = (\emptyset, \mathcal{K})$ – apriori, we don't know any axiom being necessarily in a MUPS ($D_{s_0} = \emptyset$), but potentially all axioms can be there ($P_{s_0} = \mathcal{T} \cup \mathcal{A}$). Next, we define $Z = (s_0)$ and $R = \emptyset$

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3. (Test) If $R(C, D_s \cup P_s) = \mathit{true}$ then no subset of $D_s \cup P_s$ can cause unsatisfiability – we continue with step 2.

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4 (Finding an unsatisfiable set) We add $D_s \cup P_s$ into R and remove from R all $s' \in R$ such that $D_s \cup P_s \subseteq s'$. For $P_s = \alpha_1, \ldots, \alpha_N$ we push to Z a new state $(D_s \cup \{\alpha_1, \ldots, \alpha_{i-1}\}, P_s \setminus \{\alpha_1, \ldots, \alpha_i\})$ – we continue with step 2.

CS-tree Exploration (2)

- Soundness : Step 4 is important – here, we cover all possibilities. It always holds that $D_s \cup P_s$ differs to $D_s' \cup P_s'$ by just one element, where s' is a successor of s.

CS-tree Exploration (2)

- Soundness : Step 4 is important – here, we cover all possibilities. It always holds that $D_s \cup P_s$ differs to $D'_s \cup P'_s$ by just one element, where s' is a successor of s.
- Finiteness : Set $D_s \cup P_s$ is finite at the beginning and gets smaller with the tree depth. Furthermore, in step 4 we generate only finite number of states.

Algorithm based on Reiter's Algorithm

Another Approach – Reiter's Algorithm

There is an alternative to CS-trees:

1. Find a single (arbitrary) MUPS (*singleMUPS* in the next slides).
2. "remove the source of unsatisfiability provided by MUPS" (Reiter's algorithm in the next slides) from the set of axioms go explore the remaining axioms in the same manner.

Algorithm based on Reiter's Algorithm

Finding a single $MUPS(C, Y)$ – example

Example

The run of $singleMUPS(Person, \mathcal{K}_5)$ introduced next.

Finding a single $MUPS(C, Y)$ – example

Example

The run of *singleMUPS*$(Person, \mathcal{K}_5)$ introduced next.

1.PHASE :

$$\mathcal{K}_5 = \{\alpha_1, \alpha_2, \alpha_3\} \quad R(Person, \{\alpha_1\}) = true$$

$$S = \{\alpha_1\}$$

Finding a single $MUPS(C, Y)$ – example

Example

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1.PHASE :

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$$\mathcal{K}_5 = \{\alpha_1, \alpha_2, \alpha_3\} \qquad R(Person, \{\alpha_1, \alpha_2\}) = \mathit{false}$$
$$S = \{\alpha_1, \alpha_2\}$$

2.PHASE :

$$S = \{\alpha_1, \alpha_2\} \qquad R(Person, \{\alpha_1, \alpha_2\} - \{\alpha_1\}) = \mathit{true}$$
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singleMUPS(C, Y) – finding a single MUPS

The following algorithm is polynomial in the number of tableau algorithm applications – the computational complexity stems from the complexity of tableau algorithm itself.

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1. (Initialization) Denote $S = \emptyset$, $K = \emptyset$
2. (Finding superset of MUPS) While $R(C, S) = \mathit{false}$, then $S = S \cup \{\alpha\}$ for some $\alpha \in Y \setminus S$.

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The following algorithm is polynomial in the number of tableau algorithm applications – the computational complexity stems from the complexity of tableau algorithm itself.

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2. (Finding superset of MUPS) While $R(C, S) = \mathit{false}$, then $S = S \cup \{\alpha\}$ for some $\alpha \in Y \setminus S$.
3. (Pruning found set) For each $\alpha \in S \setminus K$ evaluate $R(C, S \setminus \{\alpha\})$. If the result is *false*, then $K = K \cup \{\alpha\}$. The resulting K is itself a MUPS.

Finding all MUPSes – Reiter Algorithm, example

Example (continued)

images/l11/reiter.png

Finding all MUPSes – Reiter Algorithm, example

Example (continued)

images/l11/reiter.png

Finding all MUPSes – Reiter Algorithm

- Reiter algorithm runs $\mathit{singleMUPS}(C, Y)$ multiple times to construct so called "Hitting Set Tree", nodes of which are pairs $(\mathcal{K}_i, M_i)$, where $\mathcal{K}_i$ lacks some axioms comparing to $\mathcal{K}$ and $M_i = \mathit{singleMUPS}(C, \mathcal{K}_i)$, or $M_i =$ "SAT'', if C is satisfiable w.r.t. $\mathcal{K}_i$.

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- Paths from the root to leaves build up *diagnoses* (i.e. minimal sets of axioms, each of which removed from $\mathcal{K}$ causes satisfiability of C).

Finding all MUPSes – Reiter Algorithm

- Reiter algorithm runs *singleMUPS*(C, Y) multiple times to construct so called "Hitting Set Tree", nodes of which are pairs $(\mathcal{K}_i, M_i)$, where $\mathcal{K}_i$ lacks some axioms comparing to $\mathcal{K}$ and $M_i = singleMUPS(C, \mathcal{K}_i)$, or $M_i =$ "*SAT*", if C is satisfiable w.r.t. $\mathcal{K}_i$.
- Paths from the root to leaves build up *diagnoses* (i.e. minimal sets of axioms, each of which removed from $\mathcal{K}$ causes satisfiability of C).
- Number of *singleMUPS*(C, Y) calls is at most exponential w.r.t. the initial axioms count. Why ?

Finding all MUPSes – Reiter Algorithm (2)

1 (Initialization) Find a single MUPS for C in $\mathcal{K}$, and construct the root $s_0 = (\mathcal{K}, singleMUPS(C, \mathcal{K}))$ of the hitting set tree. Next, set $Z = (s_0)$.

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2. (Depth First Search) If Z is empty, STOP.
3. (Test) Otherwise pop an element from Z and denote it as $s_i = (\mathcal{K}_i, M_i)$. If $M_i = \text{``}SAT''$, then go to step 2.
4. (Decomposition) For each $\alpha \in M_i$ insert into Z a new node $(\mathcal{K}_i \setminus \{\alpha\}, singleMUPS(\mathcal{K}_i \setminus \{\alpha\}, C))$. Go to step 2.

Modeling Error Explanation – Summary

- finding MUPSes is the most common way for explaining modeling errors.
- black-box vs. glass box methods. Other methods involve e.g. incremental methods [**bsw2003famus**].
- the goal is to find MUPSes (and diagnoses) – what to do in order to solve a modeling problem (unsatisfiability,inconsistency).
- above mentioned methods are quite universal – they can be used for many other problems that are not related with description logics.