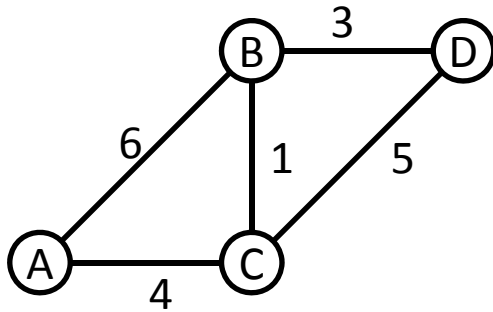


Dijkstra algorithm in time $O(N^2)$, small example



Weight matrix

	A	B	C	D
A	0	6	4	0
B	6	0	1	3
C	4	1	0	5
D	0	3	5	0

Progress Table - initial state

node dist pred closed

A	0	null	false
B	inf	null	false
C	inf	null	false
D	inf	null	false

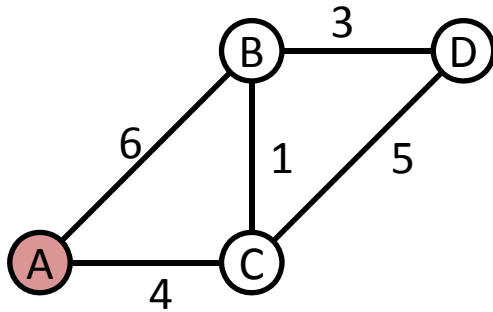
Dijkstra repeatedly updates the table

Init: for each node X : $\text{dist}[X] = \text{infinity}$; $\text{pred}[X] = \text{null}$; $\text{dist}[\text{start}] = 0$; $\text{closed}[X] = \text{false}$;

for i in $[1..N-1]$

1. *Find the open node X with smallest distance from start.* **// $O(N)$**
 $\text{minDist} = \text{infinity}$
 for Y in nodes: *// Apply linear (slowest!) search,*
 if $\text{not}(\text{closed}[Y])$ and $\text{dist}[Y] < \text{minDist}$:
 $X = Y$; $\text{minDist} = \text{dist}[Y]$
2. *// Recalculate the distances of all non-closed neighbours of X .* **// $O(\text{deg}(X)) \subset O(N)$**
 for YY in neighbours X :
 if $\text{dist}(YY) > \text{dist}(X) + \text{edgeWeight}(YY, X)$
 $\text{dist}(YY) = \text{dist}(X) + \text{edgeWeight}(YY, X)$; $\text{pred}(YY) = X$
3. $\text{closed}[X] = \text{true}$ **// $O(1)$**

Dijkstra algorithm in time $O(N^2)$, small example



Weight matrix

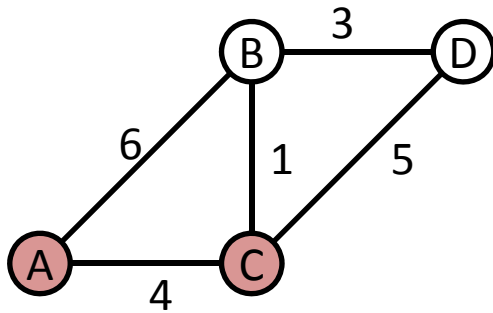
	A	B	C	D
A	0	6	4	0
B	6	0	1	3
C	4	1	0	5
D	0	3	5	0

Progress Table

node dist pred closed

A	0	null	true
B	6	A	false
C	4	A	false
D	inf	null	false

A is the closest, neighbours updated

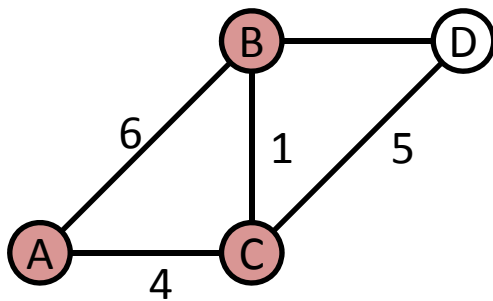


	A	B	C	D
A	0	6	4	0
B	6	0	1	3
C	4	1	0	5
D	0	3	5	0

node dist pred closed

A	0	null	true
B	5	C	false
C	4	A	true
D	9	C	false

C is the closest, neighbours updated



	A	B	C	D
A	0	6	4	0
B	6	0	1	3
C	4	1	0	5
D	0	3	5	0

node dist pred closed

A	0	null	true
B	5	C	true
C	4	A	true
D	8	B	false

B is the closest, neighbours updated

 closed

 updated values in each step