

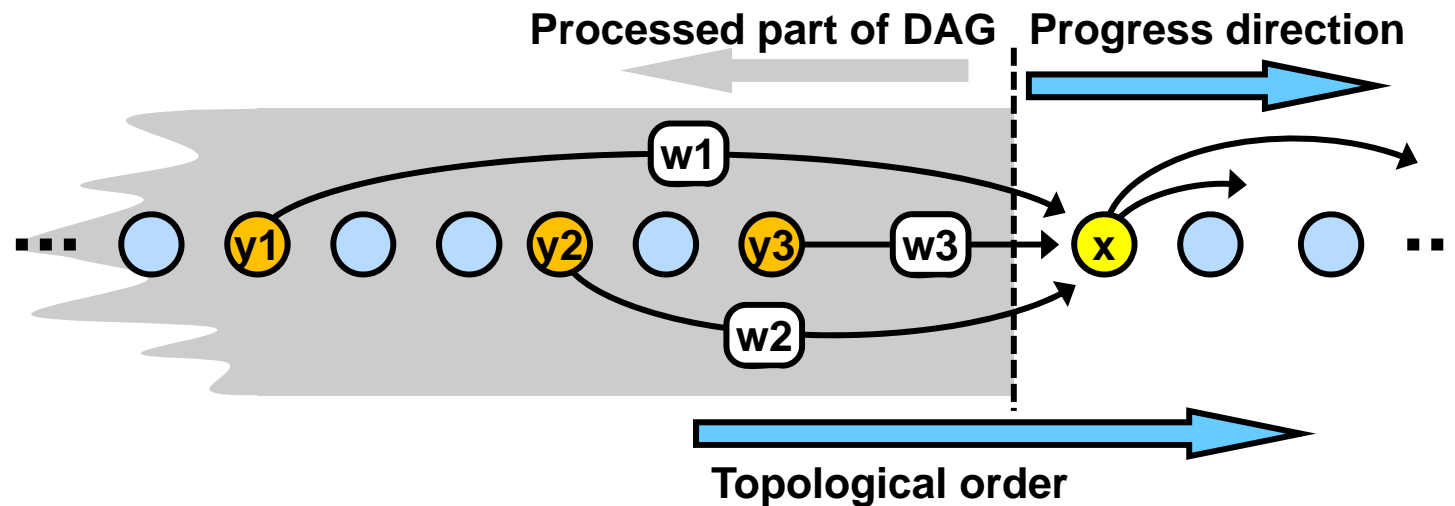
Longest path in DAG

We process nodes of DAG in their topological order.

Denote by $d[x]$ length of the path which ends in x and its length is maximal.

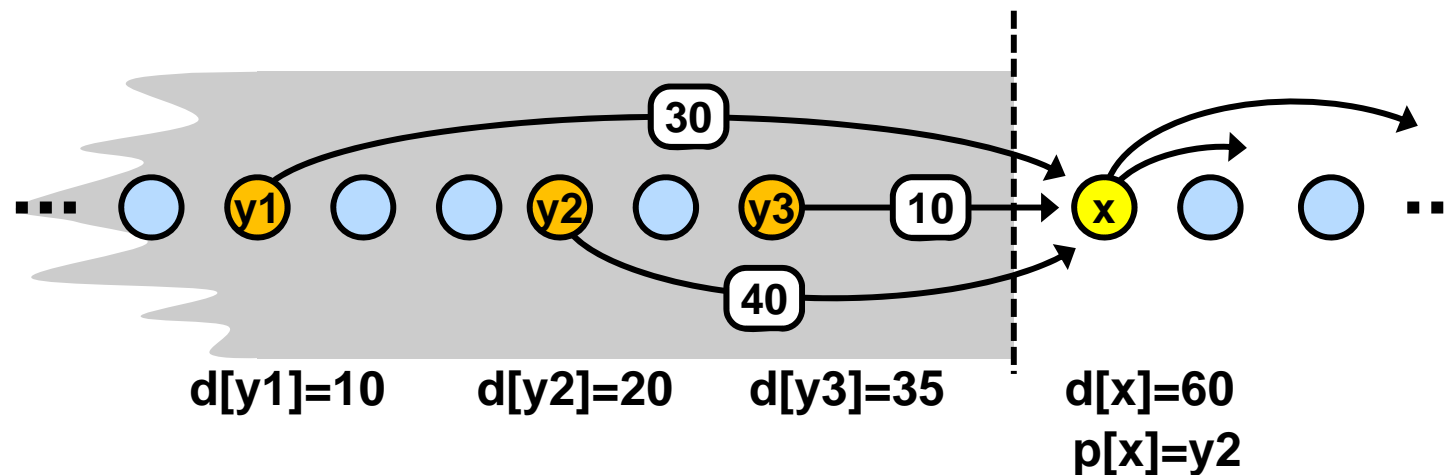
Charakteristic DP view "from the end to the beginning":

- $d[x]$ is set when values of d are known for all previous (= already processed) nodes in the topological order.
- $d[x]$ is the maximum of values $\{d[y_1] + w_1, d[y_2] + w_2, \dots, d[y_k] + w_k\}$, where $(y_1, x), (y_2, x), \dots$ are all edges ending in x and w_1, w_2, \dots are their respective weights.



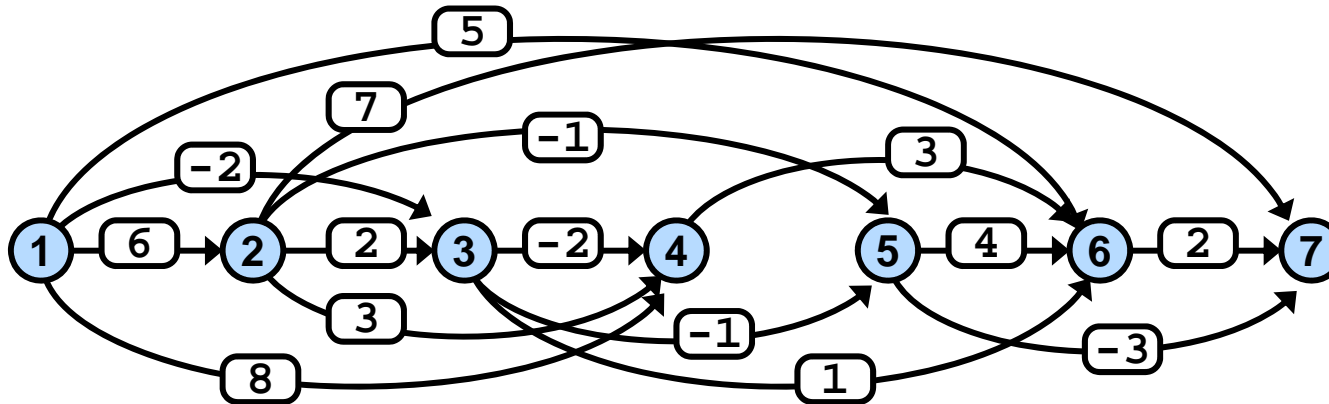
Longest path in DAG

- $d[x]$ is the maximum of values $\{d[y_1] + w_1, d[y_2] + w_2, \dots, d[y_k] + w_k\}$, where $(y_1, x), (y_2, x), \dots$ are all edges ending in x and w_1, w_2, \dots are their respective weights.
- If all values $\{d[y_1] + w_1, d[y_2] + w_2, \dots, d[y_k] + w_k\}$ are negative then none of them contributes to the longest path and the value of $d[x]$ is reset: $d[x] = 0$.
- The node y_j , for which the value $d[y_j] + w_j$ is maximal and non-negative, is set as a predecessor of x on the longest path.



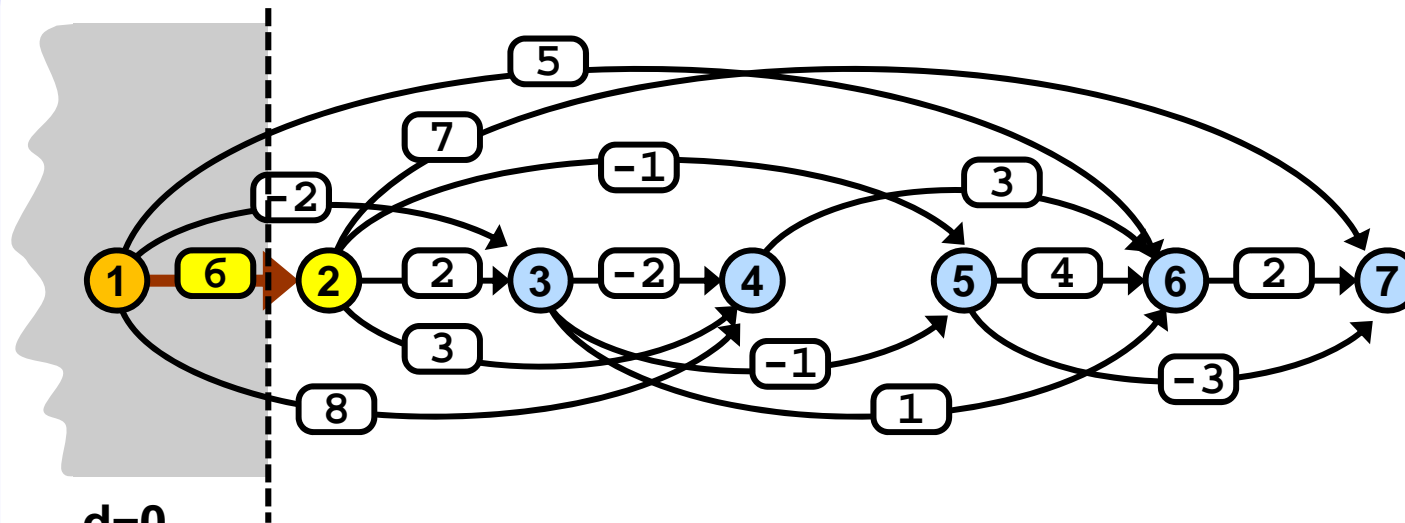
Longest path in DAG

Example



Find the longest path and its length.

Longest path in DAG



$d=0$

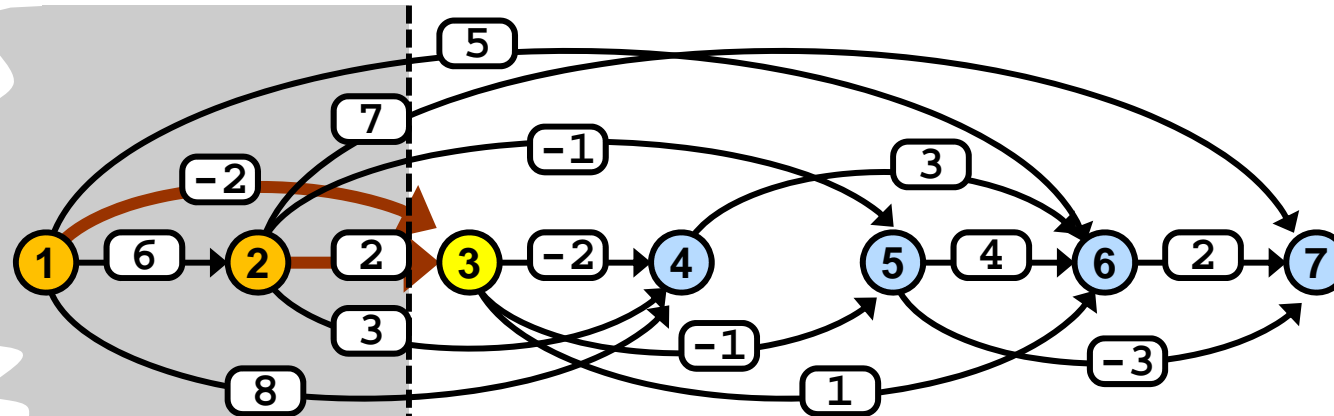
$p=nil$

$$d = \max \{0+6\}$$

$$= 6$$

$p=1$

Longest path in DAG



$d=0$

$p=\text{nil}$

$d=6$

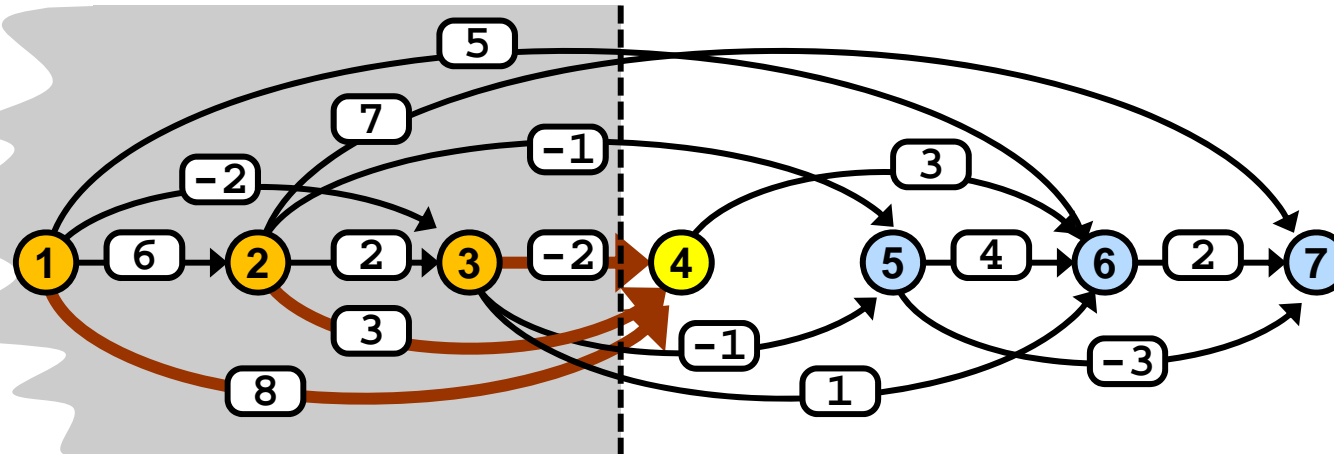
$p=1$

$$d = \max \{0 + -2, 6 + 2\}$$

$$= 8$$

$$p = 2$$

Longest path in DAG



$d=0$
 $p=\text{nil}$

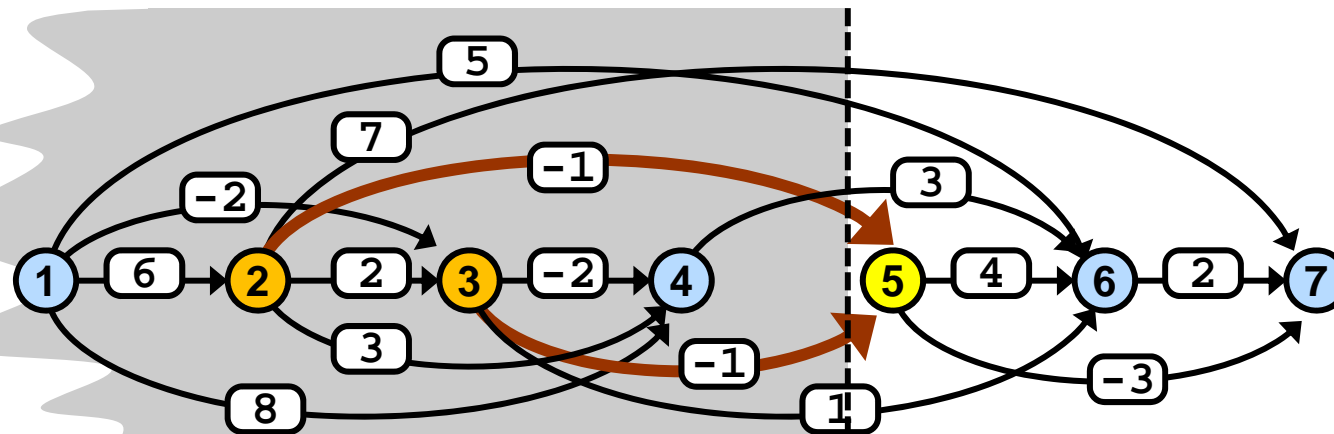
$d=6$
 $p=1$

$d=8$
 $p=2$

$d = \max \{0+8,$
 $6+3,$
 $8+-2\}$
 $= 9$

$p = 2$

Longest path in DAG



d=0
p=nil

d=6
p=1

d=8
p=2

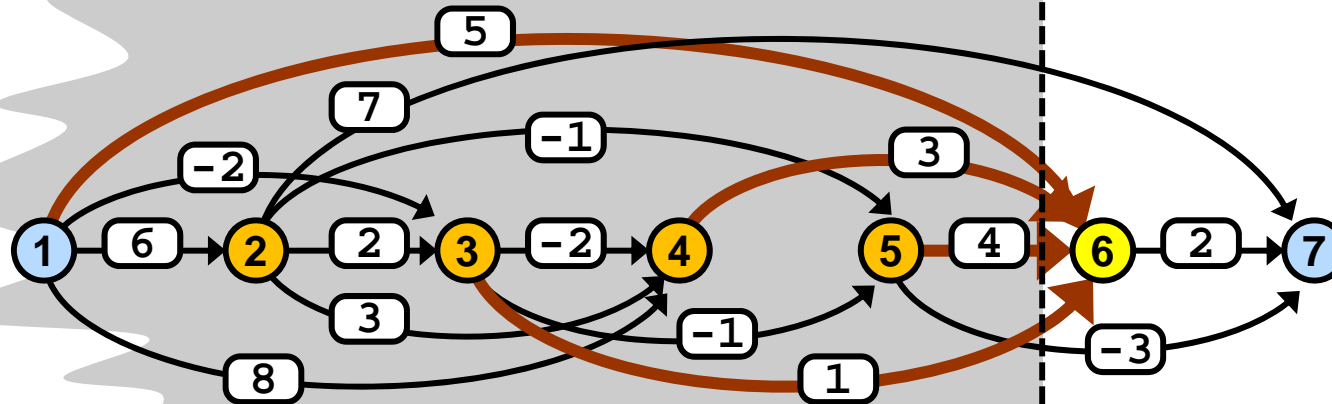
d=9
p=2

$d = \max \{6 + -1, 8 + -1\}$

= 7

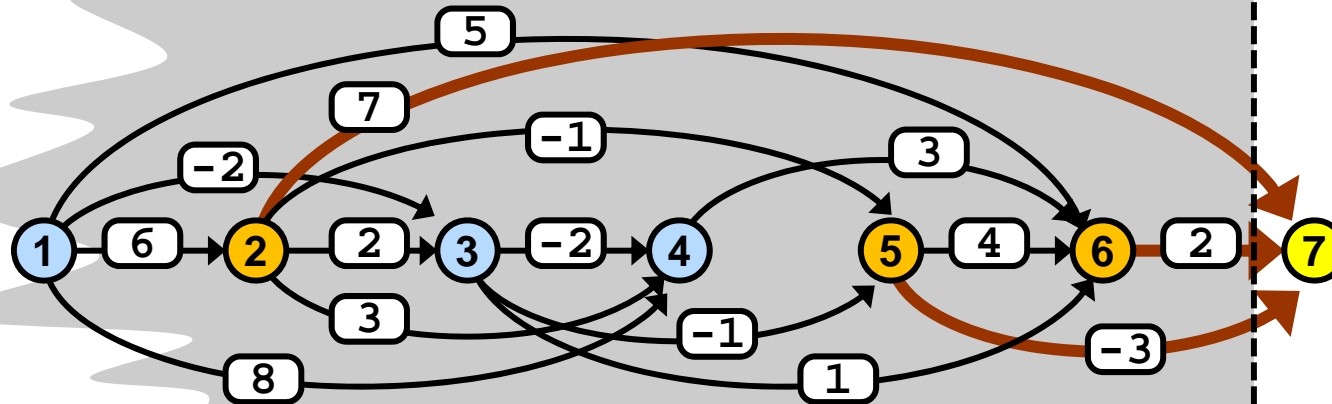
p = 3

Longest path in DAG



d=0	d=6	d=8	d=9	d=7	
p=nil	p=1	p=2	p=2	p=3	d = max {0+5,
					8+1,
					9+3,
					7+4}
					= 12
					p = 4

Longest path in DAG



d=0
p=nil

d=6
p=1

d=8
p=2

d=9
p=2

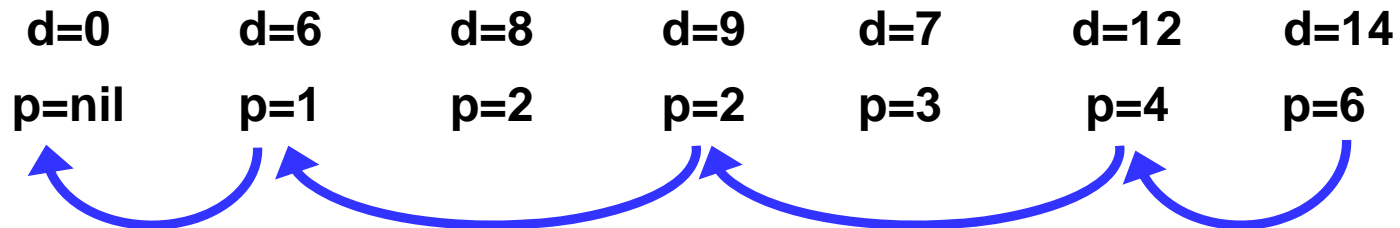
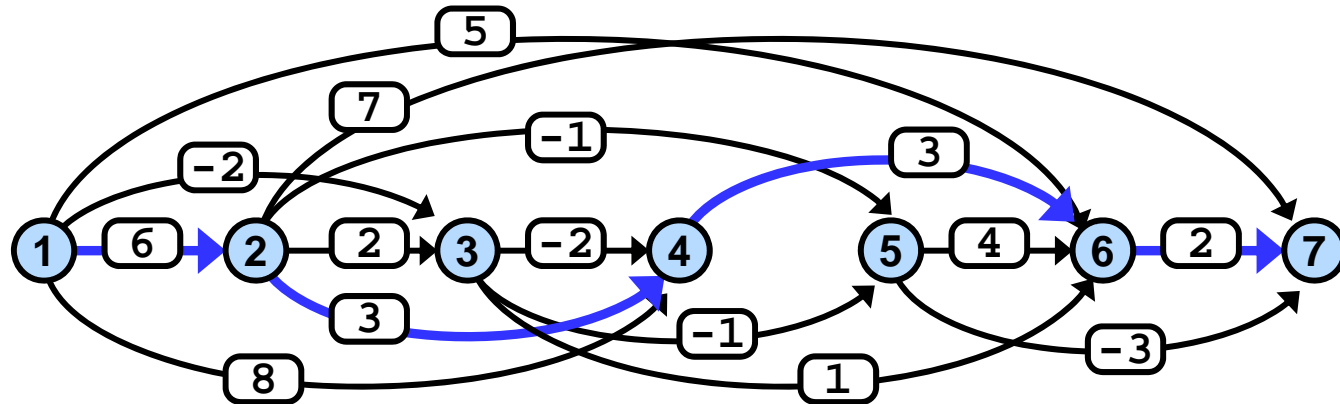
d=7
p=3

d=12
p=4

$d = \max \{6+7, 7+-3, 12+2\}$
 $= 14$

p = 6

Longest path in DAG



Length of the longest path: 14

The longest path itself: 1 -- 2 -- 4 -- 6 -- 7

Longest path in DAG

0. allocate memory for distance and predecessor of each node

```
1. for each x in V(G) {
    x.dist = negInfinity
    x.pred = null
}
```

// supposing nodes are processed
// in ascending topological order

```
2. for each node x in V(G)
    for each edge e = (y, x) in E(G)
        if (x. dist < y.dist + e.weight) {
            x. dist = y.dist + e.weight
            x.pred = y;
        }
    if (x. dist < 0) x.dist = 0; // avoid negative path lengths
}
```

0. Complexity $\Theta(N)$

1. Complexity $\Theta(N)$

2. Complexity $\Theta(M)$,
each edge is visited exactly
once and it is processed in
constant time.

Complexity: $\Theta(N+M)$

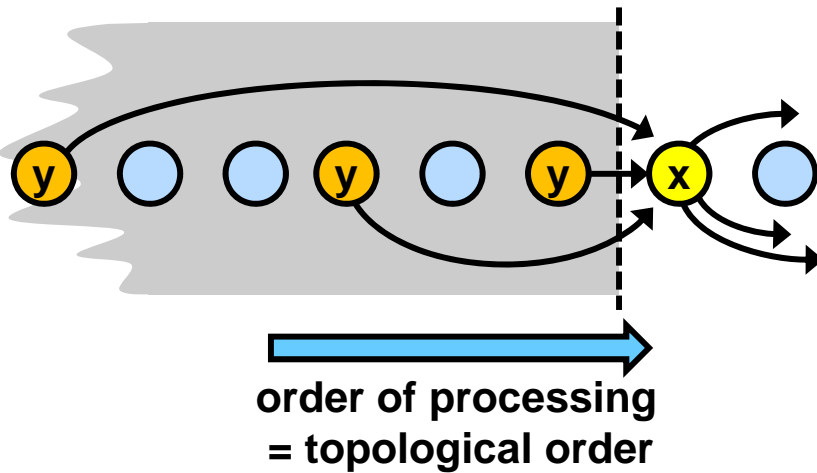
Longest path in DAG

Variant I

```

2. for each node x in  $V(G)$  {
  for each edge  $e = (y, x)$  in  $E(G)$ 
  if ( $x.\text{dist} < y.\text{dist} + e.\text{weight}$ ) {
     $x.\text{dist} = y.\text{dist} + e.\text{weight}$ 
     $x.\text{pred} = y$ ;
  }
  if ( $x.\text{dist} < 0$ )  $x.\text{dist} = 0$ ;
}

```

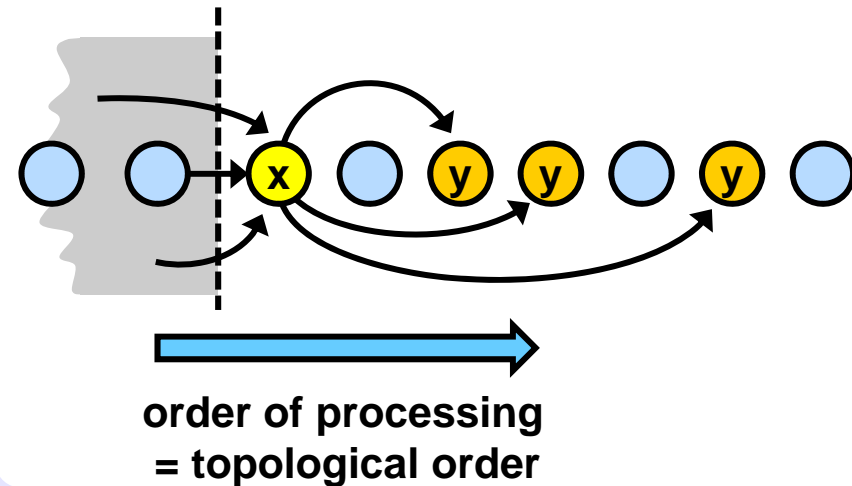


Variant II

```

2. for each node x in  $V(G)$  {
  if ( $x.\text{dist} < 0$ )  $x.\text{dist} = 0$ ;
  for each edge  $e = (x, y)$  in  $E(G)$ 
  if ( $y.\text{dist} < x.\text{dist} + e.\text{weight}$ ) {
     $y.\text{dist} = x.\text{dist} + e.\text{weight}$ 
     $y.\text{pred} = x$ ;
  }
}

```

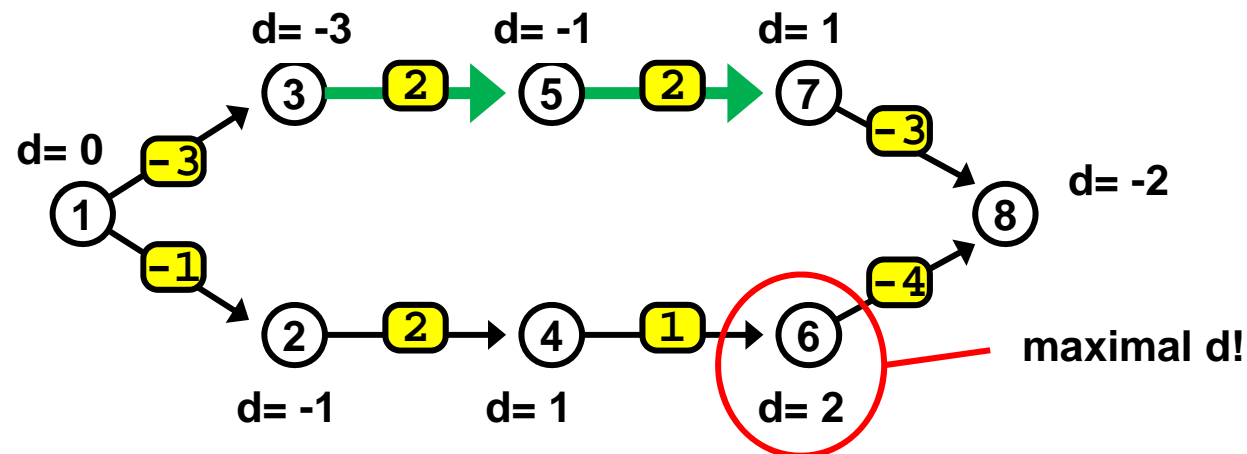


Longest path in DAG

Warning

Algorithms presented in the literature and on the web often solve the maximum path in DAG problem only for non-negative edge weights and do not mention explicitly this limitation. Those algorithms cannot handle DAG containing negative weight edges.

Incorrect result produced by algorithm expecting only non-negative edge weights

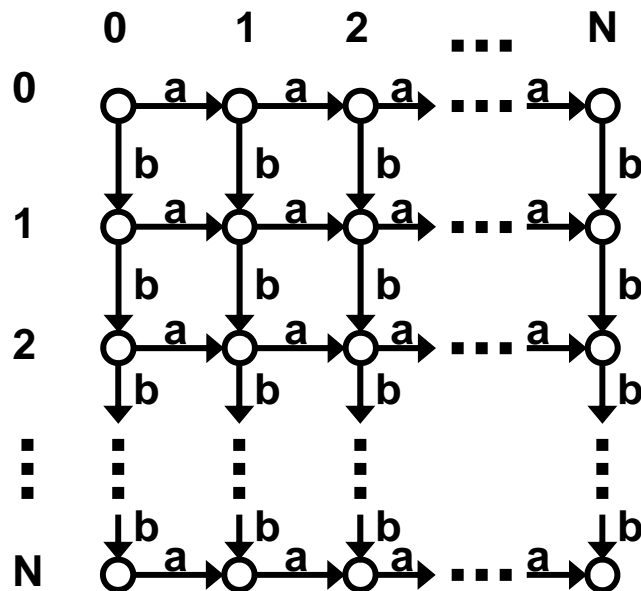


Actual maximum path is 3 -- 5 -- 7 which weight is 4.
Algorithm limited to non-negative weights finds only suboptimal path 1 -- 2 -- 4 -- 6 which weight is 2.

Longest path in DAG

Problem of reconstructing optimal paths
 -- the number of paths can be too high.

Example



Each path from the root to the leaf is optimal, its weight is $N \cdot (a+b)$.

Number of all paths is $\text{Comb}(2N, N)$, it holds that $2^N < \text{Comb}(2N, N) < 4^N$.

The number of optimal paths thus grows exponentially with the value of N .

N	# of optimal paths
1	2
10	184756
20	137846528820
30	118264581564861424
40	107507208733336176461620