

Relaxation heuristics

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PUI Tutorial
Week 3

- Any questions regarding the lecture?



- THANK YOU for filling the feedback form
- mutexes - will be in more lectures
- speeeeeed - always let me know (tell me, use chat)
- more STRIPS/FDR/PDDL examples - will do!

Obtaining heuristics

- Many different possible ways to obtain a heuristic in general
- General idea:

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Obtaining heuristics

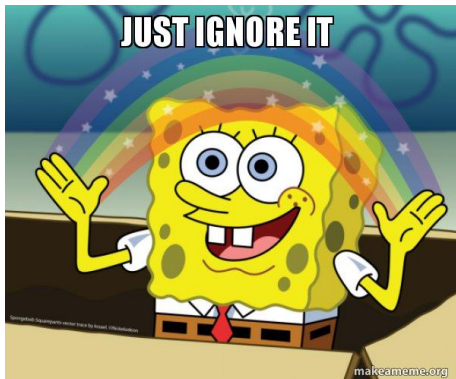
- Many different possible ways to obtain a heuristic in general
- General idea: solve a **simplified** version of the problem
 - relaxation
 - abstraction
- This week: **relaxation**



- Relaxation is
 - general design technique
 - usually ignoring something
 - simplifying the problem

Relaxation heuristic

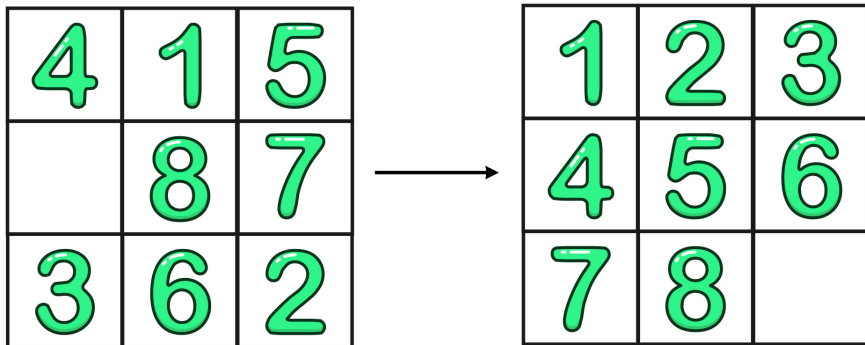
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- **Example:** 8-puzzle in STRIPS

Relaxation heuristic

- **Example:** 8-puzzle in STRIPS



- How would you formulate it?

STRIPS = $\langle F, O, S_i, g \rangle$

F = Facts

n1p1 - number 1 on position 1

n1p2 - number 1 on position 2

...

n8p9 - number 8 on position 9

free(p1) - position 1 is free

...

free(p9) - position 8 is free

next(p1,p2) - position 1 and 2 are adjacent

...

next(p8,p9) - position 8 and 9 are adjacent

O = Operators \langle pre, add, del \rangle

move-n1-p1-p2 (move number 1 from position 1 to position 2)

- *preconditions: {n1p1, free(p2)}*
- *add effects: {n1p2, free(p1)}*
- *delete effects: {n1p1}*

...

S_i (initial state) = {n1p2, n2p9, n3p7, n4p1, n5p3, n6p8, n7p6, n8p5}

g (goal state specification) = {n1p1, n2p2, n3p3, n4p4, n5p5, n6p6, n7p7, n8p8}

- What do we relax?

- What do we relax? Delete effects

- What do we relax? Delete effects
- **Delete relaxation**

Relaxed STRIPS planning task

Relaxation of a STRIPS planning task $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$ is the planning task $\Pi^+ = \langle F, O^+, s_{init}, s_{goal}, c \rangle$ which contains set of relaxed operators.

Relaxation heuristic

- What do we relax? Delete effects
- **Delete relaxation**

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Relaxation of operators

Relaxation of operator $o = \langle pre(o), add(o), del(o) \rangle$ is operator $o^+ = \langle pre(o), add(o), \emptyset \rangle$.

- Let's try it out! <http://editor.planning.domains>

- Operator $move(n, p1, p2)$
 - h^* - works with STRIPS definition Π
 - h^+ - works with relaxed STRIPS definition Π^+

h^+ heuristic

The h^+ heuristic computes length of the optimal relaxed plan π^+ which is an optimal solution to the relaxed problem Π^+ .

- Is h^+ admissible?

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- Is h^+ admissible? Yes, if it's optimal.
- Computing h^+ is still very hard though.
- We can compute an **estimate** of h^+ .
 - h^{add}
 - h^{max}

h^{add} heuristic

- STRIPS planning task $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$
- $h^{add}(s)$ gives estimate of the distance from s to a state that satisfies s_{goal}
- $h^{add}(s) = \sum_{f \in s} \Delta_0(s, f)$, where
 - $\Delta_0(s, o) = \sum_{f \in pre(o)} \Delta_0(s, f), \forall o \in O$
 - $\Delta_0(s, f) = \begin{cases} 0 & \text{if } f \in s, \\ \text{inf} & \text{if } \forall o \in O : f \notin add(o), \\ \min\{c(o) + \Delta_0(s, o) \mid o \in O, f \in add(o)\} & \text{otherwise.} \end{cases}$

h^{max} heuristic

- STRIPS planning task $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$
- $h^{max}(s)$ gives estimate of the distance from s to a state that satisfies s_{goal}
- $h^{max}(s) = \max_{f \in s_{goal}} \Delta_1(s, f)$, where
 - $\Delta_1(s, o) = \max_{f \in pre(o)} \Delta_1(s, f), \forall o \in O$
 - $\Delta_1(s, f) = \begin{cases} 0 & \text{if } f \in s, \\ \inf & \text{if } \forall o \in O : f \notin add(o), \\ \min\{c(o) + \Delta_1(s, o) \mid o \in O, f \in add(o)\} & \text{otherwise.} \end{cases}$

Exercise h^{add} , h^{max}

Compute $h^{max}(s_{init})$ for the following problem $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$:

$$F = \{a, b, c, d, e, f, g\}$$

$$O =$$

	pre	add	del	c
o_1	{a}	{c,d}	{a}	1
o_2	{a,b}	{e}	\emptyset	1
o_3	{b,e}	{d,f}	{a,e}	1
o_4	{b}	{a}	\emptyset	1
o_5	{d,e}	{g}	{e}	1

$$s_{init} = \{a, b\} \quad s_{goal} = \{f, g\}$$

Algorithm 1: Algorithm for computing $h^{\max}(s)$.

Input: $\Pi = \langle \mathcal{F}, \mathcal{O}, s_{init}, s_{goal}, c \rangle$, state s

Output: $h^{\max}(s)$

```
1 for each  $f \in s$  do  $\Delta_1(s, f) \leftarrow 0$ ;  
2 for each  $f \in \mathcal{F} \setminus s$  do  $\Delta_1(s, f) \leftarrow \infty$ ;  
3 for each  $o \in \mathcal{O}$ ,  $pre(o) = \emptyset$  do  
4   | for each  $f \in add(o)$  do  $\Delta_1(s, f) \leftarrow \min\{\Delta_1(s, f), c(o)\}$  ;  
5 end  
6 for each  $o \in \mathcal{O}$  do  $U(o) \leftarrow |pre(o)|$ ;  
7  $C \leftarrow \emptyset$ ;  
8 while  $s_{goal} \not\subseteq C$  do  
9   |  $k \leftarrow \arg \min_{f \in \mathcal{F} \setminus C} \Delta_1(s, f)$ ;  
10  |  $C \leftarrow C \cup \{k\}$ ;  
11  | for each  $o \in \mathcal{O}, k \in pre(o)$  do  
12  |   |  $U(o) \leftarrow U(o) - 1$ ;  
13  |   | if  $U(o) = 0$  then  
14  |   |   | for each  $f \in add(o)$  do  
15  |   |   |   |  $\Delta_1(s, f) \leftarrow \min\{\Delta_1(s, f), c(o) + \Delta_1(s, k)\}$ ;  
16  |   |   |   | end  
17  |   |   | end  
18  |   | end  
19 end  
20  $h^{\max}(s) = \max_{f \in s_{goal}} \Delta_1(s, f)$ ;
```

Exercise h^{add} , h^{max}

Compute $h^{add}(s_{init})$ for the following problem $\Pi = \langle F, O, s_{init}, s_{goal}, c \rangle$:

$$F = \{a, b, c, d, e, f, g\}$$

$$O =$$

	pre	add	del	c
o_1	{a}	{c,d}	{a}	1
o_2	{a,b}	{e}	\emptyset	1
o_3	{b,e}	{d,f}	{a,e}	1
o_4	{b}	{a}	\emptyset	1
o_5	{d,e}	{g}	{e}	1

$$s_{init} = \{a, b\} \quad s_{goal} = \{f, g\}$$

What's the difference in the algorithm?

Heuristic dominance

Admissible heuristic h_1 dominates an admissible heuristic h_2 if for every state s $h_1(s) \geq h_2(s)$

- h^+ is **admissible, consistent**
- h^{max} is **admissible, consistent**
- h^{add} is **not admissible, nor consistent** but can be very informative
- $h^{max} \leq h^+ \leq h^*$

- Relaxation heuristics
- Delete relaxation
- h^+ , h^{max} and h^{add} heuristics
- Implementation of h^{max}



[Feedback form link](#)

