

Combinatorial Optimization

Lab 02: Examples

March 1, 2019

Abstract: During the second lab of the course, we went through several simple examples of the ILP programs. This document provides short descriptions of the problems. The source codes in C++, Java and Python programming languages are located in the respective folders. Even though the problems are simple, they are \mathcal{NP} -hard (except for the LP example). Students should be able to formulate these problems formally and write the corresponding ILP models.

For a short description of the programming interfaces see, e.g., [C++](#), [Java](#), [Python](#). We also encourage you to go through the [Gurobi documentation](http://www.gurobi.com/documentation/) (<http://www.gurobi.com/documentation/>) – you can find there descriptions of the individual methods, as well as many useful examples.

1 Two partition problem

Given multiset S of positive integers, we ask if it can be partitioned to two subsets S_1, S_2 , such that sum of the numbers in S_1 equals to the sum of the numbers in S_2 , and each number is assigned (either to S_1 or to S_2).

Example 1 Multiset $S = \{1, 1, 1, 2, 2, 3\}$ can be partitioned, e.g., to $S_1 = \{2, 3\}$, $S_2 = \{1, 1, 1, 2\}$. Note that it is not a unique solution.

Example 2 Multiset $S = \{2, 4\}$ cannot be partitioned.

2 Maximum independent set

Given graph $G = (V, E)$, we are trying to find subset $V^* \subseteq V$, such that V^* is as large as possible and for each pair of vertices $u, v \in V^*$, $u \neq v$, it holds that there is no edge $e = \{u, v\}$ in the original graph.

3 Minimum vertex cover

Given graph $G = (V, E)$, we are looking for subset $V^* \subseteq V$, such that V^* is as small as possible and each edge $e = \{u, v\}$ is ‘covered’ by at least one vertex $x \in V^*$, i.e. $\forall e = \{u, v\} \in E \exists x \in V^*$ such that $x = u$ or $x = v$.

4 SAT (boolean satisfiability problem)

Given formula φ (in CNF), we are looking for truth assignment (True / False) to each variable, such that φ is satisfied (evaluates to True).

Example 1 Formula $\varphi = (x_1 \vee x_2) \wedge (\neg x_2 \vee x_3)$ can be satisfied by $x_1 := \text{True}$, $x_3 := \text{True}$, $x_2 := \text{False}$.

Example 2 Formula $\varphi = (x_1) \wedge (\neg x_1)$ cannot be satisfied.