### Lecture 9: Haskell Types

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A **type** is a name for a collection of related values (e.g., basic, composed, functions, etc.). For example, in Haskell the basic type

### Bool

contains the two logical values:

True

False

Applying a function to one or more arguments of the wrong type is called a **type error**.

Prelude> 1 + False

1 is a number and False is a logical value, but + requires two numbers.

- Applying a function to the wrong number of arguments
- Permuting the arguments of a function
- Forgetting the application of a (conversion) function

Java types: A function defined to return a string may return

- a String
- null
- an exception
- in addition to that, it can also return a "modified state of the world" (e.g. it can print a line of text to the console)
- not return at all (e.g. run into an infinite loop)

Functional languages can be (and often are) completely rigorous about types

- If it says it returns a String, it does that and only that
- No side effects!

Types in functional languages often more expressive

• They allow representing more complex properties of programs

- Types are checked without executing a function
  - E.g., type error in an unreachable branch is still detected
  - infinite stream can have a clear type
- All type errors are found at compile time
  - + safer: if it compiles, there is not type mismatch
  - + faster: no need for type checks at run time
  - + clearer?: can serve as documentation / specification
  - sometimes more verbose code
  - slower compilation
  - more complex compiler implementation

If evaluating an expression e would produce a value of type t, then e  $\ensuremath{\text{has type}}$  t, written

Every well-formed expression has a type automatically calculated at **compile time** using a process called **type inference**.

In GHCi, the :type command calculates the type of an expression, without evaluating it:

> :type not False not False :: Bool Haskell has a number of **basic types**, including:

Bool
Char
String
Int
Integer
Float

logical values single characters strings of characters fixed-precision integers arbitrary-precision integers floating-point numbers A tuple is a sequence of values of possibly different types:

```
(False,'a',True) :: (Bool,Char,Bool)
(False,True) :: (Bool,Bool)
```

The type of n-tuples whose i-th element has type ti is (t1,t2,...,tn)

- The type of a tuple encodes its size
- The type of the components is unrestricted

A list is a sequence of values of the same type:

```
[False,True,False] :: [Bool]
['a','b','c','d'] :: [Char]
```

In general, for any type a [a] is the type of lists with elements of type a

- The type of a list says nothing about its length
- The type of the elements can be arbitrary (not only basic)

Types of functions are denoted using ->

```
add :: (Int,Int) -> Int
add (x,y) = x+y
zeroto :: Int -> [Int]
zeroto n = [0..n]
```

- The argument and result types are unrestricted
- It is encouraged to write types above each function

Functions with multiple arguments are also possible by returning **functions as results**:

```
add :: (Int,Int) -> Int
add (x,y) = x+y
add' :: Int -> (Int -> Int)
add' x y = x+y
```

add and add' produce the same final result, but add take arguments in a different form

#### It transparently works for multiple arguments

mult :: Int -> (Int -> (Int -> Int))
mult x y z = x\*y\*z

mult takes an integer x and returns a function mult x, which in turn takes an integer y and returns a function mult x y, which finally takes an integer z and returns the result x\*y\*z.

Curried functions are more flexible than functions on tuples, because useful functions can often be made by **partially applying** a curried function.

```
add' 1 :: Int -> Int
take 5 :: [Int] -> [Int]
drop 5 :: [Int] -> [Int]
```

# Currying Conventions

To avoid excess parentheses when using curried functions, two conventions are adopted:

1) The arrow operator -> associates to the **right**. Int -> Int -> Int -> Int means Int -> (Int -> (Int -> Int))

2) As a consequence, it is then natural for function application to associate to the left.

mult x y z means ((mult x) y) z

Unless tupling is explicitly required, all functions in Haskell are normally defined in curried form.

A function is called **polymorphic** (of many forms) if its type contains type variables.

Type variables can be instantiated to different types in different circumstances

> length [False,True] -- a = Bool
2
> length [1,2,3,4] -- a = Int
4

Many of the functions defined in the standard prelude are polymorphic.

```
fst :: (a,b) -> a
head :: [a] -> a
take :: Int -> [a] -> [a]
zip :: [a] -> [b] -> [(a,b)]
id :: a -> a
```

## Overloaded functions

A polymorphic function is called **overloaded** if its type contains one or more class constraints.

(+) :: Num a => a -> a -> a

For any numeric type a, (+) takes two values of type a and returns a value of type a.

Constrained type variables can be instantiated to any types that satisfy the constraints:

> 1 + 2 -- a = Int
3
> 1.0 + 2.0 -- a = Float
3.0
> 'a' + 'b' -- Char is not numeric
ERROR

#### Haskell has a number of type classes, including:

Numeric types

Equality types

Ord Ordered types

For example, you can verify by calling :type:

(+) :: Num a => a -> a -> a (==) :: Eq a => a -> a -> Bool (<) :: Ord a => a -> a -> Bool

Eq

- When defining a new function in Haskell, it is useful to begin by writing down its type;
- Within a script, it is good practice to state the type of every new function defined;
- When stating the types of polymorphic functions that use numbers, equality or orderings, take care to include the necessary class constraints.

In Haskell, a new **name** for an **existing type** can be defined using a **type** declaration.

Type declarations make other types easier to read.

```
type Pos = (Int,Int)
left :: Pos -> Pos
left (x,y) = (x-1,y)
```

Like function definitions, type declarations can also have **parameters**. With

we can define:

```
mult :: Pair Int -> Int
mult (m,n) = m*n
copy :: a -> Pair a
copy x = (x,x)
```

Type declarations can be nested:

```
type Pos = (Int,Int)
```

```
type Trans = Pos -> Pos
```



However, they cannot be recursive:

```
type Tree = (Int,[Tree])
```



Define a completely new type by specifying its values

data Bool = False | True

Values False and True are the **constructors** for the type

Type and constructor names begin with a capital letter

Values of new types can be used in the same ways as those of built in types. Given

data Answer = Yes | No | Unknown

we can define:

```
answers :: [Answer]
answers = [Yes,No,Unknown]
flip :: Answer -> Answer
flip Yes = No
flip No = Yes
flip Unknown = Unknown
```

### Parametric constructors

The constructors in a data declaration can have parameters. Given

```
data Shape = Circle Float | Rect Float Float
```

we can define:

square :: Float -> Shape
square n = Rect n n

Circle and Rect can be viewed as  $\ensuremath{\textbf{functions}}$  that construct values of type Shape

New composed data types can be decomposed by **pattern matching** 

```
area :: Shape -> Float
area (Circle r) = pi * r<sup>2</sup>
area (Rect x y) = x * y
```

One of the most common Haskell types

data Maybe a = Nothing | Just a

allows defining safe operations.

```
safediv :: Int -> Int -> Maybe Int
safediv _ 0 = Nothing
safediv m n = Just (m `div` n)
safehead :: [a] -> Maybe a
safehead [] = Nothing
safehead xs = Just (head xs)
```

New types can be declared in terms of themselves. That is, types can be **recursive**. (just not with type keyword)

data Nat = Zero | Succ Nat

A value of type Nat is either Zero, or Succ n where n :: Nat. Nat contains infinite sequence of values:

Zero Succ Zero Succ (Succ Zero)

. . .

We can use pattern matching and recursion to translate from Int to Nat and back.

```
nat2int :: Nat -> Int
nat2int Zero = 0
nat2int (Succ n) = 1 + nat2int n
int2nat :: Int -> Nat
int2nat 0 = Zero
int2nat n = Succ (int2nat (n-1))
```

Two naturals can be added by converting them to integers, adding, and then converting back:

add :: Nat -> Nat -> Nat add m n = int2nat (nat2int m + nat2int n)

However, using recursion the function add can be defined without the need for conversions:

add Zero n = nadd (Succ m) n = Succ (add m n)

## Records

Purely positional data declarations are impractical with a large number of fields. Therefore, the fields can be named:

This allows to define records in arbitrary order

And access fields using automatically generated functions, e.g.,

```
firstName :: Person -> String
```

Recursive typed can represent tree structures, such as **expressions** from numbers, plus, multiplication.

```
data Expr = Val Int
| Add Expr Expr
| Mul Expr Expr
```

1+2\*3

Add (Val 1) (Mul (Val 2) (Val 3))

Using recursion, it is now easy to define functions that process expressions. For example:

```
size :: Expr -> Int
size (Val n) = 1
size (Add x y) = size x + size y
size (Mul x y) = size x + size y
eval :: Expr -> Int
eval (Val n) = n
eval (Add x y) = eval x + eval y
eval (Mul x y) = eval x * eval y
```

## Homework assignment 4

- Everything has a type known in compile time
  - basic values
  - functions
  - data structures
- Types are key for data structures in Haskell

### • Algebraic types

compose complex types from simpler as products and unions

- Types can be instances of classes
  - polymorphic functions