



Functional Programming

Lecture 9: Haskell Types

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What is a Type?

A type is a name for a collection of related values. For example, in Haskell the basic type

Bool

contains the two logical values:

False

True

Type Errors

Applying a function to one or more arguments of the wrong type is called a type error.

```
> 1 + False  
error ...
```

1 is a number and False is a logical value, but + requires two numbers.

Types in Haskell

If evaluating an expression e would produce a value of type t , then e has type t , written

$$e :: t$$

Every well formed expression has a type, which can be automatically calculated at compile time using a process called type inference.

Static typing

- All type errors are found at compile time
 - safer: if it compiles, there is not type mismatch
 - faster: no need for type checks at run time
- In GHCi, the `:type` command calculates the type of an expression, without evaluating it:

```
> :type not False
not False :: Bool
```

Basic Types

Haskell has a number of basic types, including:

- | | |
|---------|--------------------------------|
| Bool | - logical values |
| Char | - single characters |
| String | - strings of characters |
| Int | - fixed-precision integers |
| Integer | - arbitrary-precision integers |
| Float | - floating-point numbers |

List Types

A list is a sequence of values of the same type:

```
[False, True, False] :: [Bool]
['a', 'b', 'c', 'd'] :: [Char]
```

In general, for any type a

$[a]$ is the type of lists with elements of type a

The type of a list says nothing about its length

The type of the elements can be arbitrary

Tuple Types

A tuple is a sequence of values of different types:

```
(False, True) :: (Bool, Bool)
```

```
(False, 'a', True) :: (Bool, Char, Bool)
```

The type of n-tuples whose i-th element has type t_i is

(t_1, t_2, \dots, t_n)

The type of a tuple encodes its size

The type of the components is unrestricted

Function Types

A function is a mapping from values of one type to values of another type:

```
not  :: Bool → Bool
even :: Int  → Bool
```

In general:

$t1 \rightarrow t2$ is the type of functions that map values of type $t1$ to values to type $t2$.

Function Types

The arrow \rightarrow is typed at the keyboard as `->`

The argument and result types are unrestricted

It is encouraged to write types above each function

```
add :: (Int,Int) -> Int
add (x,y) = x+y
```

```
zeroto :: Int -> [Int]
zeroto n = [0..n]
```

Curried Functions

Functions with multiple arguments are also possible by returning functions as results:

$$\begin{aligned} \text{add}' &:: \text{Int} \rightarrow (\text{Int} \rightarrow \text{Int}) \\ \text{add}' \ x \ y &= x+y \end{aligned}$$

add and add' produce the same final result, but add take arguments in a different form

$$\begin{aligned} \text{add} &:: (\text{Int}, \text{Int}) \rightarrow \text{Int} \\ \text{add}' &:: \text{Int} \rightarrow (\text{Int} \rightarrow \text{Int}) \end{aligned}$$

Curried Functions

Transparently works for multiple arguments

```
mult :: Int → (Int → (Int → Int))  
mult x y z = x*y*z
```

mult takes an integer x and returns a function mult x , which in turn takes an integer y and returns a function mult $x y$, which finally takes an integer z and returns the result $x*y*z$.

Partial function application

Curried functions are more flexible than functions on tuples, because useful functions can often be made by partially applying a curried function.

```
add' 1 :: Int → Int
```

```
take 5 :: [Int] → [Int]
```

```
drop 5 :: [Int] → [Int]
```

Currying Conventions

To avoid excess parentheses when using curried functions, two conventions are adopted:

- The arrow \rightarrow associates to the right.

$\boxed{\text{Int} \rightarrow \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}}$

means $\text{Int} \rightarrow (\text{Int} \rightarrow (\text{Int} \rightarrow \text{Int}))$

- As a consequence, it is then natural for function application to associate to the left.

$\boxed{\text{mult } x \ y \ z}$

means $((\text{mult } x) \ y) \ z$

Polymorphic Functions

A function is called polymorphic (“of many forms”) if its type contains type variables.

```
length :: [a] → Int
```

Type variables can be instantiated to different types in different circumstances

```
> length [False, True]
2
> length [1, 2, 3, 4]
4
```

a = Bool

a = Int

Many of the functions defined in the standard prelude are polymorphic.

```
fst :: (a,b) → a
```

```
head :: [a] → a
```

```
take :: Int → [a] → [a]
```

```
zip :: [a] → [b] → [(a,b)]
```

```
id :: a → a
```


Overloaded Functions

A polymorphic function is called overloaded if its type contains one or more class constraints.

$$(+)\ ::\ \text{Num}\ a\ \Rightarrow\ a\ \rightarrow\ a\ \rightarrow\ a$$

For any numeric type a , $(+)$ takes two values of type a and returns a value of type a .

Constrained type variables can be instantiated to any types that satisfy the constraints:

```
> 1 + 2
3

> 1.0 + 2.0
3.0

> 'a' + 'b'
ERROR
```

a = Int

a = Float

Char is not a numeric type

Haskell has a number of type classes, including:

`Num` - Numeric types

`Eq` - Equality types

`Ord` - Ordered types

For example:

```
(+) :: Num a => a -> a -> a
```

```
(==) :: Eq a => a -> a -> Bool
```

```
(<) :: Ord a => a -> a -> Bool
```

Hints and Tips

- When defining a new function in Haskell, it is useful to begin by writing down its type;
- Within a script, it is good practice to state the type of every new function defined;
- When stating the types of polymorphic functions that use numbers, equality or orderings, take care to include the necessary class constraints.

Type Declarations

In Haskell, a new name for an existing type can be defined using a type declaration.

```
type String = [Char]
```

String is a synonym for the type [Char].

Type declarations make other types easier to read.

```
type Pos = (Int,Int)
```

```
left :: Pos → Pos
```

```
left (x,y) = (x-1,y)
```

Parametrized Types

Like function definitions, type declarations can also have parameters. With

```
type Pair a = (a, a)
```

we can define:

```
mult :: Pair Int → Int  
mult (m, n) = m*n  
  
copy :: a → Pair a  
copy x = (x, x)
```

Type declarations can be nested:

```
type Pos = (Int,Int)
type Trans = Pos → Pos
```



However, they cannot be recursive:

```
type Tree = (Int, [Tree])
```



Data Declarations

Define a completely new type by specifying its values

```
data Bool = False | True
```

Values False and True are the constructors for the type
Type and constructor names begin with a capital letter

Values of new types can be used in the same ways as those of built in types. Given

```
data Answer = Yes | No | Unknown
```

we can define:

```
answers :: [Answer]
answers = [Yes, No, Unknown]

flip :: Answer → Answer
flip Yes      = No
flip No       = Yes
flip Unknown  = Unknown
```

Parametric Constructors

The constructors in a data declaration can also have parameters. Given

```
data Shape = Circle Float  
           | Rect Float Float
```

we can define:

```
square :: Float → Shape  
square n = Rect n n
```

Circle and Rect can be viewed as functions that construct values of type Shape

New composed data types can still be decomposed by pattern matching

```
area :: Shape → Float
area (Circle r) = pi * r^2
area (Rect x y) = x * y
```

Parametric Data Declarations

One of the most common Haskell types

```
data Maybe a = Nothing | Just a
```

allows defining safe operations.

```
safediv :: Int → Int → Maybe Int  
safediv _ 0 = Nothing  
safediv m n = Just (m `div` n)
```

```
safehead :: [a] → Maybe a  
safehead [] = Nothing  
safehead xs = Just (head xs)
```

Recursive Types

New types can be declared in terms of themselves.
That is, types can be recursive. (just not with type keyword)

```
data Nat = Zero | Succ Nat
```

A value of type Nat is either Zero, or Succ n where n :: Nat. Nat contains infinite sequence of values:

```
Zero
```

```
Succ Zero
```

```
Succ (Succ Zero)
```

We can use pattern matching and recursion to translate from Int to Nat and back.

```
nat2int :: Nat → Int
```

```
nat2int Zero      = 0
```

```
nat2int (Succ n) = 1 + nat2int n
```

```
int2nat :: Int → Nat
```

```
int2nat 0 = Zero
```

```
int2nat n = Succ (int2nat (n-1))
```

Two naturals can be added by converting them to integers, adding, and then converting back:

```
add :: Nat → Nat → Nat
add m n = int2nat (nat2int m + nat2int n)
```

However, using recursion the function add can be defined without the need for conversions:

```
add Zero      n = n
add (Succ m) n = Succ (add m n)
```

Example: Arithmetic Expressions

Recursive typed can represent tree structures, such as expressions from numbers, plus, multiplication.

```
data Expr = Val Int
          | Add Expr Expr
          | Mul Expr Expr
```

1 + 2 * 3

```
Add (Val 1) (Mul (Val 2) (Val 3))
```


Using recursion, it is now easy to define functions that process expressions. For example:

```
size :: Expr → Int
```

```
size (Val n)    = 1
```

```
size (Add x y)  = size x + size y
```

```
size (Mul x y)  = size x + size y
```

```
eval :: Expr → Int
```

```
eval (Val n)    = n
```

```
eval (Add x y)  = eval x + eval y
```

```
eval (Mul x y)  = eval x * eval y
```

Homework assignment 4

Evaluating a log of card game Sedma

- we provide the basic types to use
- just implementing the function (no I/O)
- will need implementing instances
 - next lecture
 - use deriving for now
- deadline is two weeks from your lab

Summary

- Everything has a type known in compile time
 - basic values
 - functions
 - data structures
- Types are key for data structures in Haskell
- Types can be instances of classes
 - polymorphic functions