# Search Problems (Where reasoning consists of exploring alternatives) 



- Declarative knowledge creates alternatives:
- Which pieces of knowledge to use?
- How to use them?
- Search is a about exploring alternatives.
It is a major approach to exploit knowledge


## Example: 8-Puzzle



Initial state


Goal state

State: Any arrangement of 8 numbered tiles and an empty tile on a $3 \times 3$ board

## 8-Puzzle: Successor Function



## 8-Puzzle: Successor Function



## Search is about the exploration of alternatives

Across history, puzzles and games requiring the exploration of alternatives have been considered a challenge for human intelligence:

- Chess originated in Persia and India about 4000 years ago
- Checkers appear in 3600-year-old Egyptian paintings
- Go originated in China over 3000 years ago

So, it's not surprising that Al uses games to design and test algorithms


# $\left(n^{2}-1\right)$-puzzle algorithm scalability 



| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |

## | 5-Puzzle

Introduced (?) in 1878 by Sam Loyd, who dubbed himself "America's greatest puzzle-expert"


Joumalist and Adverticing Expert, URILINAK.
Gatmes, Kovatios, Suyplements, Soupantrs, Lia., fut Nerspajcre.
 HOR ADVTRTISING IUR:COSFS.

Allher of the tamona




$$
\therefore \text { ㄱ, } 30 \% 9.6
$$

New York, upore/5 1903

## | 5 -Puzzle

Sam Loyd offered $\$ 1,000$ of his own money to the first person who would solve the following problem:

| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |$\stackrel{?}{\mid}$| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
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## But no one ever won the prize !!

# Stating a Problem as a Search Problem 



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- State space S


# Stating a Problem as a Search Problem 



- State space S
- Successor function:
$x \in S \rightarrow \operatorname{SUCCESSORS}(x) \in 2^{S}$


# Stating a Problem as a Search Problem 



# Stating a Problem as a Search Problem 



## Stating a Problem as a Search Problem



## State Graph

- Each state is represented by a distinct node
- An arc (or edge) connects a node s to a node s' if $s^{\prime} \in \operatorname{SUCCESSORS}(s)$
- The state graph may contain more than one connected component



## Solution to the Search

- A solution is a path connecting the initial node to a goal node (any one)



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## Solution to the Search

- A solution is a path connecting the initial node to a goal node (any one)
- The cost of a path is the sum of the arc costs along this path
- An optimal solution is a solution path of minimum cost
- There might be no solution!



## How big is the state space of the ( $n^{2}-1$ )-puzzle?

- 8-puzzle $\rightarrow$ ?? states


## How big is the state space of the ( $n^{2}-I$ )-puzzle?

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- 8 -puzzle $\rightarrow 9$ ! $=362,880$ states
- 15 -puzzle $\rightarrow 16!\sim 2.09 \times 10^{13}$ states


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- 24-puzzle $\rightarrow 25$ ! ~ $10^{25}$ states


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(n²-I)-puzzle?

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But only half of these states are reachable from any given state
(but you may not know that in advance)

## Permutation Inversions

| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |

## Permutation Inversions

- Let the goal be:

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| 1 | 2 | 3 | 4 | $\mathrm{n}_{2}=0$ | $\mathrm{n}_{3}=0$ | $\mathrm{n}_{4}=0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5 | 10 | 7 | 8 | $n_{5}=0$ | $\mathrm{n}_{6}=0$ | $\mathrm{n}_{7}=1$ |
| 9 | 6 | 11 | 12 | $\mathrm{n}_{11}=0$ | $\mathrm{n}_{12}=0$ | $\mathrm{n}_{13}=0$ |
| 13 | 14 | 15 |  | $\mathrm{n}_{14}=0$ | $\mathrm{n}_{15}=0$ |  |

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- A tile $j$ appears after a tile $i$ if either $j$ appears on the same row as $i$ to the right of $i$, or on another row below the row of $i$.

| 1 | 2 | 3 | 4 | $\mathrm{n}_{2}=0$ | $\mathrm{n}_{3}=0$ | $\mathrm{n}_{4}=0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5 | 10 | 7 | 8 | $n_{5}=0$ | $\mathrm{n}_{6}=0$ | $\mathrm{n}_{7}=$ |
| 9 | 6 | 11 | 12 | $\mathrm{n}_{11}=0$ | $\mathrm{n}_{12}=0$ |  |
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- A tile $j$ appears after a tile $i$ if either $j$ appears on the same row as $i$ to the right of $i$, or on another row below the row of $i$.
- For every $i=I, 2, \ldots, I 5$, let $n_{i}$ be the number of tiles $j$ < $i$ that appear after tile i (permutation inversions)

| 1 | 2 | 3 | 4 | $\mathrm{n}_{2}=0$ | $\mathrm{n}_{3}=0$ | $\mathrm{n}_{4}=0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
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- A tile $j$ appears after a tile $i$ if either $j$ appears on the same row as $i$ to the right of $i$, or on another row below the row of $i$.
- For every $\mathrm{i}=\mathrm{I}, 2, \ldots, \mathrm{I} 5$, let $\mathrm{n}_{\mathrm{i}}$ be the number of tiles j < i that appear after tile i (permutation inversions)
- $N=n_{2}+n_{3}+\ldots+n_{15}+$ row number of empty tile

- Proposition: ( N mod 2 ) is invariant under any legal move of the empty tile
- Proof:
- Any horizontal move of the empty tile leaves N unchanged
- A vertical move of the empty tile changes N by an even increment $( \pm I \pm I \pm I \pm I)$

$s=$| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 |  | 7 |
| 9 | 10 | 11 | 8 |
| 13 | 14 | 15 | 12 |


$s^{\prime}=$| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 11 | 7 |
| 9 | 10 |  | 8 |
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$N\left(s^{\prime}\right)=N(s)+3+1$

- Proposition: ( N mod 2 ) is invariant under any legal move of the empty tile
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- It can be shown that this is also a sufficient condition
- $\rightarrow$ The state graph consists of two connected components of equal size


## 15-Puzzle

Sam Loyd offered \$1,000 of his own money to the first person who would solve the following problem:

| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |$\stackrel{?}{\square}$| 1 | 2 | 3 | 4 |
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| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |$\stackrel{?}{ } \quad$| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
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$N=4$
$N=5$

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| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 |  |$\stackrel{?}{\square}$| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 15 | 14 |  |

$N=4$
$N=5$
So, the second state is not reachable from the first, and Sam Loyd took no risk with his money ...

## What is the Actual State Space?

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a) The set of all states?
[e.g., a set of I6! states for the I5-puzzle]

## What is the Actual State Space?

a) The set of all states? [e.g., a set of 16 ! states for the 15 -puzzle]
b) The set of all states reachable from a given initial state? [e.g., a set of $16!/ 2$ states for the 15 -puzzle]

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In general, the answer is a)
[because one does not know in advance which states are reachable]

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In general, the answer is a)
[because one does not know in advance which states are reachable]
But a fast test determining whether a state is reachable from another is very useful, as search techniques are often
inefficient when a problem has no solution

## Searching the State Space



- It is often not feasible (or too expensive) to build a complete representation of the state graph


## 8-, I 5-, 24-Puzzles

## 100 millions states/sec

## 8-, | 5-, 24-Puzzles

## 8 -puzzle $\rightarrow 362,880$ states



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100 millions states/sec

## Searching the State Space

- Often it is not feasible (or too expensive) to build a complete representation of the state graph
- A problem solver must construct a solution by exploring a small portion of the graph


## Searching the State Space



## Searching the State Space



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## Searching the State Space



## Simple Problem-Solving-Agent Algorithm

1. $I \leftarrow$ sense/read initial state
2. GOAL? $\leftarrow$ select/read goal test
3. Succ $\leftarrow$ select/read successor function
4. solution $\leftarrow$ search (I, GOAL?, Succ)
5. perform(solution)

## State Space

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- Each state is an abstract representation of a collection of possible worlds sharing some crucial properties and differing on non-important details only



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## Successor Function

- It implicitly represents all the actions that are feasible in each state



## Successor Function

- It implicitly represents all the actions that are feasible in each state
- Only the results of the actions (the successor states) and their costs are returned by the function
- The successor function is a "black box": its content is unknown
E.g., in assembly planning, the successor function may be quite complex (collision, stability, grasping, ...)


## Path Cost

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- An arc cost is a positive number measuring the "cost" of performing the action corresponding to the arc, e.g.:
- I in the 8-puzzle example
- expected time to merge two sub-assemblies


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[This condition guarantees that, if path becomes arbitrarily long, its cost also becomes arbitrakily large]


## Goal State

| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 4 | 5 | 6 |
| 7 | 8 |  |

## Goal State

- It may be explicitly described:


| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 4 | 5 | 6 |
| 7 | 8 |  |

("a" stands for "any" other than 1,5 , and 8)

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- or partially described:


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## Goal State

- It may be explicitly described:
- or partially described:


| 1 | 2 | 3 |
| :--- | :--- | :--- |
| 4 | 5 | 6 |
| 7 | 8 |  |

- or defined by a condition, e.g., the sum of every row, of every column, and of every diagonal equals 30

| 15 | 1 | 2 | 12 |
| :---: | :---: | :---: | :---: |
| 4 | 10 | 9 | 7 |
| 8 | 6 | 5 | 11 |
| 3 | 13 | 14 |  |

## Other examples

## 8-Queens Problem

Place 8 queens in a chessboard so that no two queens are in the same row, column, or diagonal.


A solution


Not a solution

## Formulation \# I



- States: all arrangements of 0, I, $2, \ldots, 8$ queens on the board
- Initial state: 0 queens on the board
- Successor function: each of the successors is obtained by adding one queen in an empty square
- Arc cost: irrelevant
- Goal test: 8 queens are on the board, with no queens attacking each other


## Formulation \# I

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- Arc cost: irrelevant
- Goal test: 8 queens are on the board, with no queens attacking each other
$\rightarrow \sim 64 \times 63 x . . . \times 57 \sim 3 \times 10^{14}$ states


## Formulation \#2

- States: all arrangements of $\mathrm{k}=0$,

 $\mathrm{I}, 2, \ldots, 8$ queens in the $k$ leftmost columns with no two queens attacking each other
- Initial state: 0 queens on the board
- Successor function: each successor is obtained by adding one queen in any square that is not attacked by any queen already in the board, in the leftmost empty column
- Arc cost: irrelevant


## Formulation \#2

- States: all arrangements of $\mathrm{k}=0$,

 $\mathrm{I}, 2, \ldots, 8$ queens in the $k$ leftmost columns with no two queens attacking each other
- Initial state: 0 queens on the board
- Successor function: each successor is obtained by adding one queen in any square that is not attacked by any queen already in the board, in the leftmost empty column
- Arc cost: irrelevant


## $\rightarrow 2,057$ states

## n-Queens Problem

- A solution is a goal node, not a path to this node (typical of design problem)
- Number of states in state space:
- 8-queens $\rightarrow 2,057$
- 100-queens $\rightarrow 10^{52}$
- But techniques exist to solve n -queens problems efficiently for large values of $n$
They exploit the fact that there are many solutions well distributed in the state space


## Path Planning



## Path Planning



What is the state space?

## Formulation \# I



## Formulation \# I



Cost of one horizontal/vertical step $=1$ Cost of one diagonal step $=\sqrt{ } 2$

## Optimal Solution



This path is the shortest in the discretized state space, but not in the original continuous space

## Formulation \#2



## Formulation \#2

sweep-line



## Formulation \#2



Formulation \#2


## Formulation \#2



## Formulation \#2



## Formulation \#2



## Formulation \#2



## States



## Successor Function



## Solution Path



A path-smoothing post-processing step is usually needed to shorten the path further

## Formulation \#3



## Formulation \#3



Cost of one step: length of segment

## Formulation \#3



Cost of one step: length of segment

## Formulation \#3



Cost of one step: length of segment

## Solution Path



The shortest path in this state space is also the shortest in the original continuous space

## Assumptions in Basic Search

- The world is static
- The world is discretizable
- The world is observable
- The actions are deterministic

But many of these assumptions can be removed, and search still remains an important problem-solving tool

## Search and AI

- Search methods are ubiquitous in Al systems. They often are the backbones of both core and peripheral modules
- An autonomous robot uses search methods:
- to decide which actions to take and which sensing operations to perform,
- to quickly anticipate collision,
- to plan trajectories,
- to interpret large numerical datasets provided by sensors into compact symbolic representations,
- to diagnose why something did not happen as expected,
- etc...
- Many searches may occur concurrently and sequentially


## Applications

Search plays a key role in many applications, e.g.:

- Route finding: airline travel, networks
- Package/mail distribution
- Pipe routing,VLSI routing
- Comparison and classification of protein folds
- Pharmaceutical drug design
- Design of protein-like molecules
- Video games

