

## 14. Computing marginal probabilities for GRFs

Consider a GRF for pairs  $(x, s)$  on a graph  $(V, E)$ , where  $x: V \rightarrow F$  is a field of features and  $s: V \rightarrow K$  is a field of hidden variables

$$p_u(x, s) = \frac{1}{Z(u)} \exp \left[ \sum_{i \in V} u_i(x_i, s_i) + \sum_{j \in E} u_{ij}(s_i, s_j) \right]$$

Computing its marginal probabilities on nodes and edges like

$$p_u(s_i), p_u(s_i | x), p_u(s_i, s_j), p_u(s_i, s_j | x)$$

for  $i, j \in V$ ,  $\{i, j\} \in E$  is needed for

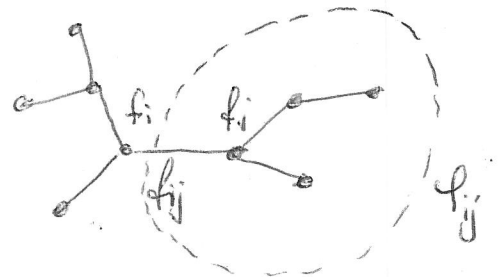
- (1) Inference with locally additive loss functions (e.g. Hamming distance)  $\Rightarrow$  optimal decision is based on  $p(s_i | x)$
- (2) Learning model parameters  $u_i, u_{ij}$ . This requires to compute node and edge marginals from  $u$ - $s$  and vice versa. (see next section)

Computing the partition function  $Z(u)$  and marginal probs for general GRFs is NP-hard. We have to rely on approximation algorithms.

### A. Belief propagation

Recall Sec. 10, computing marginal probs of a Markov model on a tree:

$$\begin{aligned} p(s) &= \frac{1}{Z} \prod_{i \in V} f_i(s_i) \prod_{j \in E} f_{ij}(s_i, s_j) \stackrel{!}{=} \\ &= \prod_{i \in V} p(s_i) \prod_{j \in E} \frac{p(s_i, s_j)}{p(s_i) p(s_j)} \end{aligned}$$



Hence, 
$$p(s_i) \propto f_i(s_i) \prod_{j \in \mathcal{N}_i} \psi_{ij}(s_i)$$

$$p(s_i, s_j) \propto f_i(s_i) f_{ij}(s_i, s_j) f_j(s_j) \prod_{l \in \mathcal{N}_i \setminus j} \psi_{il}(s_i) \prod_{m \in \mathcal{N}_j \setminus i} \psi_{jm}(s_j)$$

It follows that

$$\frac{p(s_i, s_j)}{p(s_i) p(s_j)} = \frac{f_{ij}(s_i, s_j)}{\psi_{ij}(s_i) \psi_{ji}(s_j)}$$

- recall that  $\psi$ -s are defined on oriented edges
- the above formula is an equivalent transformation in the  $(+, x)$ -domain.

The recursive definition of the  $\psi$ -s is

$$\psi_{ij}(s_i) = \sum_{s_j \in \mathcal{K}} f_{ij}(s_i, s_j) f_j(s_j) \prod_{l \in \mathcal{N}_j \setminus i} \psi_{jl}(s_j)$$

Belief propagation for random fields on general graphs (aka message passing):

(1) repeatedly recompute  $\psi$ -s according to the formula above, until (hopefully) a fixpoint  $\psi^*$  is reached

(2) estimate marginals from

$$p(s_i) \propto f_i(s_i) \prod_{j \in \mathcal{N}_i} \psi_{ij}^*(s_i)$$

$$p(s_i, s_j) \propto f_{ij}(s_i, s_j) \frac{p(s_i) p(s_j)}{\psi_{ij}^*(s_i) \psi_{ji}^*(s_j)}$$

Remark 1

(1) Quite often BP gives reasonable estimates for unary marginals. However, it inherently fails to estimate pairwise marginals.

(2) In the log-domain (i.e. replacing  $(+, \times)$  by  $(\min, +)$ )

this gives an approximation algorithm for solving  $(\min, +)$ -problems

### B. Sampling

Let  $S = \{S_i \mid i \in V\}$  be a  $K$ -valued random field with joint p.d.  $p(S)$  and let  $F: K^V \rightarrow \mathbb{R}$  be a random variable. How can we estimate its expectation

$$\mathbb{E}_p(F) = \sum_{S \in K^V} p(S) F(S)$$

by sampling?

- generate an i.i.d. sample  $\{S^j \in K^V \mid j=1, \dots, \ell\}$  of realisations from  $p(S)$
- estimate the expectation by  $\mathbb{E}_p(F) \approx \frac{1}{\ell} \sum_{j=1}^{\ell} F(S^j)$

How to sample from  $p(S)$ ? Theorem 1, Sec. 1  $\Rightarrow$  design a homogeneous Markov chain with transition probability  $T(S|S')$ ,  $S, S' \in K^V$  s.t.

- the chain is irreducible and a-periodic
- its stationary p.d. is  $p(S)$

In practice:

- Design a set of simple (sparse) transition prob. matrices  $B_m$ ,  $m \in M$  s.t.  $p(S)$  is stationary for all of them.

- Compose  $T$  by

$$T = \prod_{m \in M} B_m \quad \text{or} \quad T = \sum_{m \in M} \alpha_m B_m \quad (\alpha_m \geq 0, \sum_{m \in M} \alpha_m = 1)$$

- prove that  $T$  is irreducible and a-periodic.

Gibbs Sampler (for GRFs)

Design  $B_i, i \in V$  by

$$B_i(s|s') = \begin{cases} 0 & \text{if } s_{V \setminus i} \neq s'_{V \setminus i} \\ p(s_i | s'_{V \setminus i}) & \text{otherwise} \end{cases}$$

Stationarity of  $p(s)$

$$\begin{aligned} \sum_{s' \in K^V} B_i(s|s') p(s') &= \sum_{k \in K} p(s_i | s_{V \setminus i} = k) p(s'_{V \setminus i} = k, s_{V \setminus i}) \\ &= p(s_i | s_{V \setminus i}) p(s_{V \setminus i}) = p(s) \end{aligned}$$

It is easy to see that  $T = \prod_{i \in V} B_i$  and  $T = \sum_{i \in V} \alpha_i B_i$  are irreducible and a-periodic if  $p(s)$  is strictly positive.

Remark 2

- A Gibbs sampler is easy to implement
- Gibbs samplers are very slow: long "burn-in" time and slow mixing.

C. Mean field approximation

If only unary marginals needed  $\Rightarrow$  approximate  $p(s)$  by an factorising distribution  $q(s) = \prod_{i \in V} q_i(s_i)$  with smallest KL-divergence from  $p$

$$D_{KL}(q || p) = \sum_{s \in K^V} q(s) \log \frac{q(s)}{p(s)} \rightarrow \min_q$$

For a GRF on a graph this reads

$$\sum_{i \in V} \sum_{s_i \in K} q_i(s_i) \log q_i(s_i) - \sum_{i \in V} \sum_{s_i \in K} q_i(s_i) u_i(s_i) -$$

$$- \sum_{j \in E} \sum_{s_i, s_j \in K} q_i(s_i) q_j(s_j) u_{ij}(s_i, s_j) \rightarrow \min_{q \geq 0}$$

$$\text{s.t. } \sum_{s_i \in K} q_i(s_i) = 1 \quad \forall i \in V$$

This can be solved approximately e.g. by block-coordinate ascent.