



K-means Clustering and its Generalization

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 $\begin{array}{ll} \text{Given:} & \mathcal{T} = \{\mathbf{x}_l\}_{l=1}^L, & \text{the set of observations} \\ & K & \text{the number of desired cluster prototypes} \\ \text{Output:} & (\mathbf{c}_k)_{k=1}^K, & \text{the set of cluster prototypes (etalons)} \\ & \{\mathcal{T}_k\}_{k=1}^K & \text{the clustering (partitioning) of the data} \\ & \cup_{k=1}^K \mathcal{T}_k = \{x_l\}_{l=1}^L, \ \mathcal{T}_i \cap \mathcal{T}_j = \emptyset \text{ for } i \neq j \end{array}$

The result is obtained by solving the following optimization problem:

$$(\mathbf{c}_1, \mathbf{c}_2, ..., \mathbf{c}_K; \mathcal{T}_1, \mathcal{T}_2, ..., \mathcal{T}_K) = \operatorname*{argmin}_{\operatorname{all} \mathbf{c}'_k, \mathcal{T}'_k} J(\mathbf{c}'_1, \mathbf{c}'_2, ..., \mathbf{c}'_K; \mathcal{T}'_1, \mathcal{T}'_2, ..., \mathcal{T}'_K),$$

where





- Given: $\mathcal{T} = \{\mathbf{x}_l\}_{l=1}^L$, the set of observationsKthe number of desired cluster prototypesOutput: $(\mathbf{c}_k)_{k=1}^K$,
 $\{\mathcal{T}_k\}_{k=1}^K$ the set of cluster prototypes (etalons)
the clustering (partitioning) of the data
 $\cup_{k=1}^K \mathcal{T}_k = \{x_l\}_{l=1}^L$, $\mathcal{T}_i \cap \mathcal{T}_j = \emptyset$ for $i \neq j$
 - 1. Initialize \mathbf{c}_k (e.g. by assigning random \mathbf{x}_l to \mathbf{c}_k)
 - 2. Assignment optimization: $\mathcal{T}_k = \{ \mathbf{x} \in \mathcal{T} : \forall j, ||\mathbf{x} - \mathbf{c}_k||_2^2 \le ||\mathbf{x} - \mathbf{c}_j||_2^2 \}$
 - 3. Prototype optimization: $\mathbf{c}_{k} = \frac{1}{|\mathcal{T}_{k}|} \sum_{\mathbf{x} \in \mathcal{T}_{k}} \mathbf{x}$
 - 4. Terminate if $\mathcal{T}_k^{t+1} = \mathcal{T}_k^t$, $\forall k$; else go to 2







Number of clusters K=3

Initialization: $\mathbf{c}_k = \operatorname{random} \mathbf{x}_l,$ without replacement

K-means: an Example





Optimizing partitions: Euclidean Distances A B C D E F c_1 (5 4,5 0 1 1,4 4 c_2 (5,7 5 1 0 1 3 c_3 (6,4 5,8 1,4 1 0 3,2) Sum of squares = J¹(.) = 9.0 Optimizing prototypes:

$$c_{1} = \left(\frac{1+2+4}{3}, \frac{1+1+5}{3}\right) = (2.3, 2.3)$$

$$c_{2} = \left(\frac{5+8}{2}, \frac{5+5}{2}\right) = (6.5, 5)$$

$$c_{3} = (5, 6)$$

K-means: an Example





Optimizing partitions: Euclidean Distances A B C D E F $c_1 \begin{pmatrix} 1,9 & 1,4 & 3,1 & 3,8 & 4,5 & 6,3 \\ 6,8 & 6 & 2,5 & 1,5 & 1,8 & 1,5 \\ 6,4 & 5,8 & 1,4 & 1 & 0 & 3,2 \end{pmatrix}$ Sum of squares = J²(.) = 1.78

Optimizing prototypes:

$$c_{1} = \left(\frac{1+2}{2}, \frac{1+1}{2}\right) = (1.5,1)$$

$$c_{2} = (8,5)$$

$$c_{3} = \left(\frac{4+5+5}{3}, \frac{5+5+6}{3}\right) = (4.7,5.3)$$

K-means: an Example



(4/4)



Optimizing partitions: Euclidean Distances A B C D E F $c_1 = \begin{pmatrix} 0,5 & 0,5 & 4,7 & 5,3 & 6,1 & 7,6 \\ 8,1 & 7,2 & 4 & 3 & 3,2 & 0 \\ 5,7 & 5,1 & 0,7 & 0,5 & 0,7 & 3,3 \end{pmatrix}$

Sum of squares = $J^3(.) = 0.31$

Assignment unchanged \Rightarrow terminate





- If neither Step 3 nor Step 2 changes $J(\cdot)$, the algorithm terminates.
- Step 3 (cluster centre optimization) reduces J(·), because for a fixed assignment T_k, the mean over the data points in T_k is the optimal solution for the squared error.
- Step 2 (assignment optimization) reduces J(·) because for every x₁, the contribution to the cost function either stays the same, or gets lower.
- The fact that $J(\cdot)$ is reduced implies that no assignment is repeated during the run of the algorithm.
- Since there is a finite number of assignmens (how many?) the k-means algorithm converges, in a finite number of steps, to a local minimum.





- Alternatively, \mathcal{T}_{k} is initialised, and steps 2. and 3. are swapped
- The k-means algorithm is not a guaranteed global minimum optimizer. This is easily proved by a counter-example.
- Efficiency. The complexity of Step 2. (assignment optimization) dominates, as for every observation the nearest prototype is sought. Trivially implemented, this requires $L \times K$ operations. Any idea for a speed-up?





ln:	$\mathcal{T} = \{\mathbf{x}_{\mathbf{I}}\}_{I=1}^{L}$,	the set of observations
	d(.,.)	"distance function" (may not be a metric)
Out:	$(\mathbf{c_k})_{k=1}^{K}$,	the set of cluster prototypes (etalons)
	$\{\mathcal{T}_k\}_{k=1}^{K}$	the clustering (partitioning) of the data

- 1. Initialize c_k (e.g. by assigning random x_l to c_k)
- 2. Assignment optimization: $\mathcal{T}_k = \{ \mathbf{x} \in \mathcal{T} : \forall j, d(\mathbf{x}, \mathbf{c_k}) \le d(\mathbf{x}, \mathbf{c_j}) \}$
- 3. Prototype optimization: $\mathbf{c_k} = \arg \min_{\mathbf{c}} \sum_{\mathbf{x} \in \mathcal{T}_k} d(\mathbf{x}, \mathbf{c})$
- 4. Terminate If $\mathcal{T}_k^{t+1} = \mathcal{T}_k^t$, $\forall k$; else go to 2





ln:	$\mathcal{T} = \{\mathbf{x}_l\}_{l=1}^L$,	the set of observations
	d(.,.)	$ \mathbf{c}-\mathbf{x} _1$, ie. $d(.,.)$ is the L1-metric
Out:	$(\mathbf{c}_k)_{k=1}^{K}$,	the set of cluster prototypes (etalons)
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- 1. Initialize \mathbf{c}_k (e.g. by assigning random \mathbf{x}_l to \mathbf{c}_k)
- 2. Assignment optimization: $\mathcal{T}_k = \{ \mathbf{x} \in \mathcal{T} : \forall j, d(\mathbf{x}, \mathbf{c}_k) \le d(\mathbf{x}, \mathbf{c}_j) \}$
- 3. Prototype optimization: $\mathbf{c}_k = \text{median}\{\mathcal{T}_k\}$
- 4. Terminate If $\mathcal{T}_k^{t+1} = \mathcal{T}_k^t$, $\forall k$; else go to 2

Median is the minimizer of the L1-norm in a cluster, ie. median $\{\mathcal{T}_k\} = \mathbf{c}_{\mathbf{k}}^{\star} = \arg\min_{\mathbf{c}} \sum_{\mathbf{x} \in \mathcal{T}_k} ||\mathbf{x} - \mathbf{c}_{\mathbf{k}}||_1$





- In: $\mathcal{T} = \{\mathbf{x}_l\}_{l=1}^L$, observations \mathbf{x}_l are strings $d(s_1, s_2)$ is the Levenshtein distance, ie. the number of edit operations to transform s_1 into s_2 Out: $(\mathbf{c}_k)_{k=1}^K$, the set of cluster prototypes, \mathbf{c}_k are strings $\{\mathcal{T}_k\}_{k=1}^K$ the clustering (partitioning) of the data
 - 1. Initialize \mathbf{c}_k
 - 2. Assignment optimization: $\mathcal{T}_k = \{ \mathbf{x} \in \mathcal{T} : \forall j, d(\mathbf{x}, \mathbf{c}_k) \leq d(\mathbf{x}, \mathbf{c}_j) \}$
 - 3. Prototype optimization: $\mathbf{c}_k = \arg \min_{\mathbf{c}} \sum_{\mathbf{x} \in \mathcal{T}_k} d(\mathbf{x}, \mathbf{c})$
 - 4. Terminate If $\mathcal{T}_k^{t+1} = \mathcal{T}_k^t$, $\forall k$; else go to 2

K-means Generalization: Clustering Strings: Notes



- the calculation of d(.,.) might be non-trivial
- It might be very hard to minimize ∑_{x∈T_k} d(x, c) over the space of all strings.
 The minimization can be restricted to c ∈ T.
- Is the algorithm guaranteed to terminate if Step 2. (Step 3.) is only improving J(·), not finding the minimimum (given fixed T or c_k respectively)?





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 - 1. Initialize \mathbf{c}_k (e.g. by assigning random \mathbf{x}_l to \mathbf{c}_k)
 - 2. Assignment optimization: $\mathcal{T}_k = \{ \mathbf{x} \in \mathcal{T} : \forall j, ||\mathbf{x} - \mathbf{c}_k||_2 \le ||\mathbf{x} - \mathbf{c}_j||_2 \}$
 - 3. Prototype optimization: no closed-form solution for *geometric median*. Use e.g. iterative Weiszfeld's algorithm. $\mathbf{c}_k = \operatorname{argmin}_{\mathbf{c}} \sum_{x \in \mathcal{T}_k} \|\mathbf{x} - \mathbf{c}\|_2$
 - 4. Terminate if $\mathcal{T}_k^{t+1} = \mathcal{T}_k^t$, $\forall k$; else go to 2





macros_rpz.tex
sfmath.sty
cmpitemize.tex

Thank you for your attention.