# ParamILS: Iterated Local Search in Parameter Configuration Space

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Substantial part of this material is based on the paper Frank Hutter, Holger H. Hoos, Kevin Leyton-Brown, and Thomas Stützle: ParamILS: An Automatic Algorithm Configuration Framework, Journal of Artificial Intelligence Research (JAIR), volume 36, pp. 267-306, October 2009. See http://www.cs.ubc.ca/ hutter/papers/Hutter09PhD.pdf



http://cw.felk.cvut.cz/doku.php/courses/a0m33eoa/start

# Algorithm Configuration (Parameter Tuning) Problem

Let  $\mathcal{A}$  denote an algorithm, whose parameters are to be optimized for a **probability distribution** of problem instances,  $\mathcal{D}$ .

 ${\mathcal D}$  may be given

- implicitly, as through a random instance generator or a distribution over such generators, or
- as the uniform distribution over a finite sample of problem instances.

With each of the algorithm parameters,  $p_1 \dots p_k$ , a domain of possible values is associated and the parameter configuration space,  $\Theta$ , is the cross-product of these domains (or a subset thereof).

• We assume that all parameter domains are finite sets.

This assumption can be met by discretizing all numerical parameters to a finite number of values.

 While parameters may be ordered, we do not exploit such ordering relations => all parameters are finite and categorical.

The elements of  $\Theta$  are called parameter configurations,  $\theta_i$ , and  $\mathcal{A}(\theta)$  denotes algorithm  $\mathcal{A}$  with parameter configuration  $\theta \in \Theta$ .

The objective of the parameter configuration (parameter tuning) problem is to find the parameter configuration  $\theta \in \Theta$  resulting in the best performance of  $\mathcal{A}$  on distribution  $\mathcal{D}$ .

There are many ways of measuring the **cost**,  $o(A, \theta, I, s)$ , of running algorithm A with parameter configuration  $\theta$  on an instance I, using seed s in case of randomized algorithm. For example, we might be interested in

- the computational resources consumed by the given algorithm (such as runtime, memory or communication bandwidth), or
- the approximation error, or
- the improvement achieved over an instance-specific reference cost,
- the quality of the solution found.

The behaviour of the algorithms can vary significantly between multiple runs on different instances or when randomized algorithms are run repeatedly with fixed parameters on a single problem instance.

Therefore, the cost of a candidate solution  $\theta$  is defined as

$$c(\theta) = m(O_{\theta})$$

the statistical population parameter m of the cost distribution  $O_{\theta}(\mathcal{A}, \theta, \mathcal{D})$ , over instances drawn from distribution of instances,  $\mathcal{D}$ , and multiple independent runs.

An optimal solution,  $\theta^*$ , minimizes  $c(\theta)$ :

$$\theta^* \in \arg\min_{\theta\in\Theta} c(\theta).$$

For example, we might aim to minimize mean runtime or median solution cost.

The  $O_{\theta}(\mathcal{A}, \theta, \mathcal{D})$  is typically unknown, so we can only acquire approximations of their statistics,  $c(\theta)$ , based on a limited number of samples (i.e. the cost of single executions of  $\mathcal{A}(\theta)$ ) – let's denote an approximation of  $c(\theta)$  based on N samples by  $\widehat{c}_N(\theta)$ .

- For deterministic algorithms, the algorithm A is run on N ≤ M instances (M is the size of the finite training set of instances).
- For randomized, algorithms, we can run multiple runs with different seeds if M < N.

### Manually-executed local search in parameter configuration space

- 1. Begin with some initial parameter configuration.
- 2. Experiment with modifications to single parameter values at a time, accepting new configurations whenever they result in improved performance (**iterative first improvement procedure**).
- 3. Repeat step 2 until no single-parameter change yields an improvement.

This procedure stops as soon as it **reaches a local optimum** (a parameter configuration that cannot be improved by modifying a single parameter value).

#### Issues when trying to automate the process

- Which parameter configurations should be evaluated?
- Which problem instances should be used and how many runs should be performed on each instance?

Stochastic nature of the algorithm configuration problem.

• Which cutoff time  $\kappa_i$  (the maximum amount of time the configured algorithm is allowed to use) should be used for each run?

Employs **Iterated Local Search** that builds a chain of local optima by iterating through a main loop consisting of:

- 1. a solution perturbation to escape from local optima,
- 2. a subsidiary local search procedure and
- 3. an acceptance criterion to decide whether to keep or reject a newly obtained candidate solution.

### $ParamILS(\theta_0, r, p_{restart}, s)$

- 1. uses a combination of default and random settings for initialization,
  - $\theta_0$  is the initial parameter configuration, and
  - r is the number of randomly chosen configurations for initialization,
- 2. uses a one-exchange neighborhood (one parameter is modified in each search step),
- 3. employs iterative first improvement as a subsidiary local search procedure,
- 4. uses a fixed number, s, of random moves for perturbation,
- 5. always accepts better or equally-good parameter configurations,
- 6. re-initializes the search at random with probability  $p_{restart}$ .

# ParamILS(N): Algorithm

```
Input : Initial configuration \theta_0 \in \Theta, algorithm parameters r, p_{restart}, and s.
    Output : Best parameter configuration \theta found.
         r i = 1, ..., r \operatorname{do}
\theta \leftarrow \operatorname{random} \theta \in \Theta;
if \operatorname{better}(\theta, \theta_0) then \theta_0 \leftarrow \theta;
 1 for i = 1, ..., r do
                                                                          Initialization
 2
 3
 4 \theta_{ils} \leftarrow IterativeFirstImprovement (\theta_0); First iteration of the ILS procedure
 5 while not TerminationCriterion() do
          \theta \leftarrow \theta_{ils};
 6
          // ===== Perturbation
          for i = 1, \ldots, s do \theta \leftarrow random \theta' \in Nbh(\theta);
 7
                                                                                           Main body
          // ==== Basic local search
                                                                                          of the ILS procedure
          \theta \leftarrow IterativeFirstImprovement(\theta);
 8
          // ===== AcceptanceCriterion
          if better(\theta, \theta_{ils}) then \theta_{ils} \leftarrow \theta;
 9
          with probability p_{restart} do \theta_{ils} \leftarrow random \theta \in \Theta;
10
11 return overall best \theta_{inc} found;
12 Procedure IterativeFirstImprovement (\theta)
13 repeat
          \theta' \leftarrow \theta:
14
          foreach \theta'' \in Nbh(\theta') in randomized order do
15
                if \underline{better}(\theta'', \theta') then \theta \leftarrow \theta''; break;
16
                                                                                     How to define
                                                                                     better(\theta'', \theta')?
17 until \theta' = \theta:
18 return \theta;
```

Basic variant, BasicILS(N), uses procedure  $better_N(\theta_1, \theta_2)$  that compares two cost approximations  $\widehat{c}_N(\theta_1)$  and  $\widehat{c}_N(\theta_2)$  based on exactly N samples from the respective cost distributions  $O_{\theta}(\mathcal{A}, \theta_1, \mathcal{D})$  and  $O_{\theta}(\mathcal{A}, \theta_2, \mathcal{D})$  – the same N instances are used for all configurations  $\theta_i$ .

Procedure  $better_N(\theta_1, \theta_2)$  simply compares estimates  $\hat{c}_N(\theta_1)$  and  $\hat{c}_N(\theta_2)$  based on the same N instances using the same random seeds.

• It updates the best-so-far solution,  $\theta_{inc}$ .

	Input	: Parameter configuration $\theta_1$ , parameter configuration $\theta_2$
	Output	: True if $\theta_1$ does better than or equal to $\theta_2$ on the first N instances; false otherwise
	Side Effect	: Adds runs to the global caches of performed algorithm runs $\mathbf{R}_{\theta_1}$ and $\mathbf{R}_{\theta_2}$ ; potentially
		updates the incumbent $\theta_{inc}$
1	$\hat{c}_N(\boldsymbol{\theta}_2) \leftarrow \boldsymbol{\theta}_2$	$objective(\boldsymbol{ heta}_2,N)$
2	$\hat{c}_N(\boldsymbol{\theta}_1) \leftarrow \boldsymbol{\theta}_1$	$objective(\boldsymbol{\theta_1},N)$
3	return $\hat{c}_N($	$oldsymbol{ heta}_1) \leq \hat{c}_N(oldsymbol{ heta}_2)$

### The question is how to choose the optimal number of training instances, N?

- Using too small N leads to good training performance, but poor generalization to previously unseen test benchmarks.
- On the other hand, we cannot evaluate every parameter configuration on an enormous training set - if we did, search progress would be unreasonably slow.

*FocusedILS* is a variant of ParamILS that **adaptively varies the number of training samples** considered from one parameter configuration to another in order **to focus samples on promising configurations**.

•  $N(\theta)$  denotes the number of runs available to estimate the cost statistic  $c(\theta)$ .

The question is how to compare two parameter configurations  $\theta_1$  and  $\theta_2$  for which  $N(\theta_1) \leq N(\theta_2)$ ?

• What if we computed the empirical statistics based on the available number of runs for each configuration?

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• What if we computed the empirical statistics based on the available number of runs for each configuration?

Can lead to systematic bias if, for example, the first instances are easier than the average ones.

**Domination**: Configuration  $\theta_1$  dominates  $\theta_2$  when at least as many runs have been conducted on  $\theta_1$  as on  $\theta_2$ , and the performance of  $\mathcal{A}(\theta)$  on the first  $N(\theta_1)$  runs is at least as good as that of  $\mathcal{A}(\theta_2)$  on all of its runs.

 $\theta_1$  dominates  $\theta_2$  if and only if  $N(\theta_1) \ge N(\theta_2)$  and  $\widehat{c}_{N(\theta_2)}(\theta_1) \le \widehat{c}_{N(\theta_2)}(\theta_2)$ .

ParamILS version, called FocusedILS, encodes a comparison strategy based on the domination in procedure  $better_{Foc}(\theta_1, \theta_2)$ .

Procedure  $better_{Foc}(\theta_1, \theta_2)$ 

- 1. first it acquires one additional run for the configuration i having smaller  $N(\theta_i)$ , or one run for both configurations if  $N(\theta_1) = N(\theta_2)$ ;
- 2. then, it continues performing runs in this way until one configuration dominates the other. It returns true if  $\theta_1$  dominates  $\theta_2$ , and false otherwise.

It keeps track of the total number, B, of configurations evaluated since the last improving step.

- Whenever  $better_{Foc}(\theta_1, \theta_2)$  returns true, B extra (bonus) runs are performed for  $\theta_1$  and B is reset to 0.
- This way it is ensured that many runs are performed with good configurations  $\implies$  the error made in every comparison of two configurations  $\theta_1$  and  $\theta_2$  decreases on expectation.

: Parameter configuration  $\theta_1$ , parameter configuration  $\theta_2$ Input : True if  $\theta_1$  dominates  $\theta_2$ , false otherwise Output Side Effect: Adds runs to the global caches of performed algorithm runs  $\mathbf{R}_{\theta_1}$  and  $\mathbf{R}_{\theta_2}$ ; updates the global counter B of bonus runs, and potentially the incumbent  $\theta_{inc}$ 1  $B \leftarrow B + 1$ 2 if  $N(\theta_1) < N(\theta_2)$  then  $\begin{array}{c|c} \mathbf{3} & \boldsymbol{\theta}_{min} \leftarrow \boldsymbol{\theta}_1; \boldsymbol{\theta}_{max} \leftarrow \boldsymbol{\theta}_2 \\ \mathbf{4} & \quad \text{if } N(\boldsymbol{\theta}_1) = N(\boldsymbol{\theta}_2) \text{ then } B \leftarrow B+1 \end{array}$ 5 else  $\theta_{min} \leftarrow \theta_2; \theta_{max} \leftarrow \theta_1$ 6 repeat 7  $i \leftarrow N(\boldsymbol{\theta}_{min}) + 1$ 8  $\hat{c}_i(\boldsymbol{\theta}_{max}) \leftarrow objective(\boldsymbol{\theta}_{max}, i) // If N(\boldsymbol{\theta}_{min}) = N(\boldsymbol{\theta}_{max}), adds a new run to \mathbf{R}_{\boldsymbol{\theta}_{max}}.$ 9  $\hat{c}_i(\boldsymbol{\theta}_{min}) \leftarrow objective(\boldsymbol{\theta}_{min}, i) \parallel Adds a new run to \mathbf{R}_{\boldsymbol{\theta}_{min}}.$ 10 until dominates( $\theta_1, \theta_2$ ) or dominates( $\theta_2, \theta_1$ ) 11 if dominates( $\theta_1, \theta_2$ ) then // ==== Perform B bonus runs. $\hat{c}_{N(\theta_1)+B}(\theta_1) \leftarrow objective(\theta_1, N(\theta_1)+B) \parallel Adds B \text{ new runs to } \mathbf{R}_{\theta_1}.$ 12  $B \leftarrow 0$ 13 return true 14 15 else return false 16 **Procedure** dominates( $\theta_1, \theta_2$ ) 17 if  $N(\theta_1) < N(\theta_2)$  then return false 18 return  $objective(\theta_1, N(\theta_2)) < objective(\theta_2, N(\theta_2))$ 

**Often, the computational resources are wasted** with evaluating a parameter configuration that is much worse than other, previously-seen configurations.

Ex.: Let's assume a case where parameter configuration  $\theta_1$  takes a total of 10 seconds to solve N = 100 instances (i.e. it has a mean runtime of 0.1 seconds per instance), and another parameter configuration  $\theta_2$  takes 100 seconds to solve the first of these instances.

Clearly, when comparing the mean runtimes of  $\theta_1$  and  $\theta_2$  based on this set of instances, it is not necessary to run  $\theta_2$  on remaining 99 instances. Instead, we can terminate the first run of  $\theta_2$  after  $10 + \epsilon$  seconds, which is a lower bound on  $\theta_2$ 's mean runtime of  $0.1 + \epsilon/100$ . This lower bound exceeds the mean runtime of  $\theta_1$ , so we can already be sure that  $\theta_2$  cannot do better than  $\theta_1$ .

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Adaptive capping is based on the idea of avoiding unnecessary runs of the algorithm  $\mathcal{A}$  by developing bounds on the performance measure to be optimized.

- **Trajectory-preserving capping** provably does not change *BasicILS*'s search trajectory, but can lead to large computational savings.
- **Aggressive capping** potentially yielding even better performance.

**Trajectory-Preserving Capping** (for the case of **optimizing the mean of non-negative cost function**) – implements bounded evaluation of a parameter configuration  $\theta$ , procedure  $objective(\theta, N, bound)$ , based on N runs so that

- it sequentially performs runs for  $\theta$  and after each run computes a lower bound on  $\widehat{c}_N(\theta)$  based on  $i \leq N$  runs performed so far (i.e. the lower bound  $\widehat{c}_N(\theta)$  is calculated as the sum of runtimes of each of the *i* runs, divided by N);
- once the lower bound exceeds the *bound* passed as an argument, the remaining runs for  $\theta_2$  are skipped and a large constant, worstPossibleObjective, is returned.

Procedure  $better_N(\theta_1, \theta_2)$  is modified as follows

 $\begin{array}{ll} 1 & bound \leftarrow \infty \\ 2 & \widehat{c}_N(\theta_2) \leftarrow objective(\theta_2, N, bound) \\ 3 & bound \leftarrow \widehat{c}_N(\theta_2) \\ 4 & \widehat{c}_N(\theta_1) \leftarrow objective(\theta_1, N, bound) \\ 5 & \text{return } \widehat{c}_N(\theta_1) \leq \widehat{c}_N(\theta_2) \end{array}$ 

Trajectory-preserving capping typically requires much less runtime than the standard BasicILS.

The *trajectory-preserving capping* computes the upper bound on the cumulative runtime **from the best configuration encountered in the current ILS iteration**.

New ILS iteration starts with a **perturbation of the best-so-far configuration**  $\theta$ .

Frequently, the new parameter configuration  $\theta$  is of poor quality, thus **the capping criterion does not apply as quickly as it could** if the comparison was performed against the overall *incumbent*. **Aggressive Capping** – bounds the evaluation of any configuration by the performance of the overall best-so-far (*incumbent*) configuration,  $\theta_{inc}$ , multiplied by a factor **bound multiplier**, bm.

- Different values of m may imply different search trajectories:
  - $-bm = \infty$  reduces to trajectory-preserving but **no savings strategy**.
  - -bm = 1 is a **very aggressive strategy**, since once we know the evaluated configuration  $\theta$  is worse than the  $\theta_{inc}$ , its evaluation is terminated.

Recommended setting is bm = 2.

When two configurations \(\theta\_1\) and \(\theta\_2\) are compared and the evaluations both are terminated preemptively, the configuration having solved more instances within the allowed time is considered the better one.

Ties are broken to favor moving to a new parameter configuration.

*FocusedILS* varies the number of runs used to evaluate each parameter configuration.

How can the adaptive capping be used here?

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How can the adaptive capping be used here?

• By using separate bounds for every number of runs, N.

# **ParamILS: Final Remarks**

Suitable for configuration problems with many parameters and huge configuration spaces.

Successful applications:

- **SPEAR**, a complete SAT solver
  - 26 parameters,
  - $8.34\cdot10^{17}$  possible configurations,
  - FocusedILS produced configurations that solved test problems about 100 times faster than previous state-of-the-art solvers.
- CPLEX, a prominent solver for mixed integer programming problems with carefully chosen default parameter settings
  - 76 parameters,
  - $-1.9 \cdot 10^{47}$  possible configurations,
  - FocusedILS obtained substantial improvements in terms of both the time required to find optimal solutions (speedup factors from 1.98 to 52.3) and minimizing the optimality gap with factors from 1.26 to 8.65.

Frank Hutter, Holger H. Hoos, Kevin Leyton-Brown, and Thomas Stützle: ParamILS: An Automatic Algorithm Configuration Framework. In *Journal of Artificial Intelligence Research* (JAIR), volume 36, pp. 267-306, October 2009.

Other papers and SW available at http://www.cs.ubc.ca/labs/beta/Projects/ParamILS/