



# Functional Programming

## Lecture 6: Imperative scheme and parallelism

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# Last lecture

- We do not need to modify the state
- It breaks nice properties of FP
- It can sometimes be useful
  - random access in  $O(1)$
  - data encapsulation

# "Classes and objects"

Encapsulate data and related functions

Closures combine data with functions

Assignment allows modifying the data

```
(define (make-account balance)
  (define (withdraw amount)
    (if (>= balance amount)
        (begin (set! balance (- balance amount))
                balance)
        "Insufficient funds"))
  (define (deposit amount)
    (set! balance (+ balance amount))
    balance)
  (define (dispatch m)
    (cond ((eq? m 'withdraw) withdraw)
          ((eq? m 'deposit) deposit)
          (else (error "Unknown request -- MAKE-ACCOUNT"
                        m))))
  dispatch)
```

# Lists modifications

In R5RS, we can modify lists using

```
set-car!, set-cdr!
```

List are immutable by default with #lang scheme

Need to use `mcons`, `mcar`, `set-mcar!`, ...

# Queue

```
(define (make-q)
  (let ((front '()) (back '()))
    (define (in x)
      (let ((new (mcons x '())))
        (if (null? front)
            (begin (set! front new) (set! back new))
            (begin (set-mcdr! back new) (set! back new))))))
    (define (out)
      (let ((x (mcar front)))
        (set! front (mcdr front))
        x))
    (define (dispatch m)
      (cond ((eq? m 'in) in)
            ((eq? m 'out) out)))
    dispatch))
```

# Circular lists

```
(define (make-cyclic-list! ls)
  (define (cyc! xs)
    (if (null? (cdr xs))
        (begin (set-cdr! xs ls) ls)
        (cyc! (cdr xs))))
  (cyc! ls))
```

# List?

```
(define (my-list? li)
  (define (help? fast slow)
    (cond ((null? fast) #t)
          ((eq? fast slow) #f)
          ((null? (cdr fast)) #t)
          ((eq? (cdr fast) slow) #f)
          ((not (pair? (cdr fast))) #f)
          (#t (help? (caddr fast) (cdr slow)))))
  (cond ((null? li) #t)
        ((pair? li) (help? (cdr li) li))
        (#t #f)))
```

# Hash table

There are many variants of hash tables

Create a has table comparing with `equal?`

```
(make-hash)
```

**Associate** `v` with `key` in `hash`

```
(hash-set! hash key v)
```

```
(hash-ref hash key [failure-result])
```

```
(hash-ref! hash key to-set)
```

```
hash-remove!, hash-update!
```



# Memoization

```
(define (memoize f)
  (let ((table (make-hash)))
    (lambda args
      (hash-ref! table
                 args
                 ;; If the entry isn't there
                 (lambda ()
                   (begin
                     (display "X")
                     (apply f args)) ))))))
```

# Concurrency and Parallelism

- Thread (concurrency)
  - preempt each other without cooperation
  - share state: variables, function definitions, etc.
  - in Racket, they run on one OS thread
- Futures (parallelism)
  - evaluate an expression in parallel to the main program
  - block on operations that may not run safely in parallel
- Places (parallelism)
  - separate instances of scheme
  - communicate using message passing

# Threads

Run on single OS thread

No speed-up

Waiting for slow/external event: I/O, sockets, etc.

Operations on threads

`(thread thunk)` returns thread descriptor

`thread-suspend`, `thread-resume`, `kill-thread`

Channels

`(make-channel)`, `(channel-put ch v)`

`(channel-get ch)`, `(channel-try-get ch)`

# Thread map

Executes a given function on each element of a list in a separate thread and returns the results

Useful for demonstration purposes

```
(define (thread-map f xs)
  (let ((channels
        (map (lambda (x)
              (let ((c (make-channel)))
                (thread
                  (thunk (channel-put c (f x))))
                  c
                )
              )
            ) xs)))
    (map channel-get channels)))
```

# Futures

`(require racket/future)`

`(future thunk)`

Starts evaluating an expression (given as `thunk`)

Blocks when an operation may not be safely executed

Returns a "future"

`(touch future)`

Finish evaluating the expression in the main thread

If the expression is already evaluated, return the result

As in *promise*, additional touches just return the result

# Future map

Executes a given function on each element of a list in parallel and returns the results

```
(define (future-map f list)
  (let ((res
        (map (lambda (x)
              (future (lambda () (f x))))
             list)))
    (map touch res)))
```

Futures can be visualized and analyzed using

```
(require future-visualizer)
(visualize-futures expr)
```

# Home assignment 3

## Genetic programming

Evolution inspired local search in structured data

Survival of the fittest!!!

Individual: program for the robot in the maze

Population: collection of the programs

New generation: selection, mutation, cross-over

Fitness function: see Home assignment 2

# Summary

- We do not need to modify the state
- It breaks nice properties of FP
- It can sometimes be useful
  - random access in  $O(1)$
  - objects
  - circular data structures
  - memoization
- Concurrency and parallelism