

# Automated Action Planning

## Classical Planning for Non-Classical Planning Formalisms

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# Automated Action Planning

— Classical Planning for Non-Classical Planning Formalisms

Overview

Replanning

Contingent (Stochastic) Planning

Expressiveness and Compilation

Examples

Soft Goals and Net-Benefit Planning

Conformant Planning

Belief space

$K_0$

$K_{T,M}$

# Beyond Classical Planning

## Richer models people are working on

1. Temporal Planning (action have duration)
2. Metric Planning (continuous variables)
3. Planning with Preferences
4. Planning with Resource Constraints
5. Net-benefit Planning (maximize net value of goals achieved)
6. Generalized Planning (complex control structures, such as loops)
7. Multi-agent Planning
8. Planning Under Uncertainty:
  - 8.1 Conformant Planning
  - 8.2 Contingent Planning
  - 8.3 Markov Decision Processes (MDPs)
  - 8.4 Partially Observable MDPs
  - 8.5 Conformant Probabilistic Planning (Fully Unobservable POMDPs)

# How many courses on planning do we need?

## Key Insights:

- 😊 Classical planning offers a wealth of ideas for generating good solutions, fast.
- 😞 Importing these ideas to each of the above non-classical formalisms is difficult, and often simply does not work.

## Yet:

- 😊 Goal oriented sequencing of actions is a fundamental computational problem at the heart of all planning problems.
- 😊 Classical planners have reached a certain performance level that makes them attractive for addressing this problem.

So...

# Two Strategies

## 1. Top-down:

Develop **native solvers** for **more general class of models**

+: generality

–: complexity

## 2. Bottom-up: Extend the scope of 'classical' solvers

+: efficiency

–: generality

We now explore the second approach

# Using Classical Planners within Non-Classical Planners

## Two Key Techniques:

1. **Replanning**: the classical problem is an optimistic view of the original problem
2. **Compilation**: the classical problem is equivalent to the original problem  
(possibly under certain reasonable conditions)

# Replanning

An online method for solving planning problems with some uncertainty

1. Make some assumptions → get a simpler model
2. Solve simpler model
3. Execute until your observation contradict your assumptions
4. Repeat (Replan)

An established technique:

- ▶ Underlies many closed loop controllers
- ▶ Used in motion planning under uncertainty

## Motivation: Why Analyzing the Expressive Power?

- ▶ **Expressive power** is the motivation for designing new planning languages
- ↪ Often there is the question: *Syntactic sugar* or *essential feature*?
- ▶ *Compiling away* or change planning algorithm?
- ▶ If a feature can be compiled away, then it is apparently only *syntactic sugar*.
- ▶ However, a compilation can lead to **much larger planning domain descriptions** or to **much longer plans**.
- ↪ This means the planning algorithm will probably choke, i.e., it cannot be considered as a **compilation**



## Example: DNF Preconditions

- ▶ Assume we have **DNF preconditions** in STRIPS operators
  - ▶ This can be **compiled away** as follows
  - ▶ **Split** each operator with a DNF precondition  $c_1 \vee \dots \vee c_n$  into  $n$  operators with the same effects and  $c_i$  as preconditions
- ↪ If there exists a plan for the original planning task there is one for the new planning task and *vice versa*
- The **planning task** has almost the **same size**
  - The **shortest plans** have the **same size**

## Example: Conditional effects

- ▶ Can we compile away **conditional effects** to STRIPS?
  - ▶ Example operator:  $\langle a, b \triangleright d \wedge \neg c \triangleright e \rangle$
  - ▶ Can be translated into four operators:  
 $\langle a \wedge b \wedge c, d \rangle, \langle a \wedge b \wedge \neg c, d \wedge e \rangle, \dots$
  - ▶ Plan **existence** and plan **size** are identical
  - ▶ **Exponential blowup** of domain description!
- Can this be avoided?

# FDR Planning with Soft Goals

- ▶ Planning with **soft goals** aimed at plans  $\pi$  that maximize **utility**

$$u(\pi) = \sum_{p \in \text{app}_{\pi}(I)} u(p) - \sum_{a \in \pi} \text{cost}(a)$$

- ▶ Best plans achieve best **tradeoff** between **action costs** and **rewards**  
 ~ Note: "do nothing" is always a valid plan.  
 → **Suggests conceptual difference?**
- ▶ Model used in recent planning competitions; **net-benefit track** 2008 IPC
- ▶ Yet soft goals **do not** add expressive power; they can be **compiled away**

## FDR Planning with Soft Goals

- ▶ For each soft goal  $p$ , create **new hard goal**  $p'$  initially false, and **two new actions**:
  - ▶ *collect*( $p$ ) with precondition  $p$ , effect  $p'$  and **cost** 0, and
  - ▶ *forgo*( $p$ ) with an empty precondition, effect  $p'$  and **cost**  $u(p)$
- ▶ Plans  $\pi$  maximize  $u(\pi)$  iff minimize  $cost(\pi) = \sum_{a \in \pi} cost(a)$  in resulting problem
- ▶ **Any helpful in practice?**
- ▶ Compilation yields better results than native soft goal planners in 2008 IPC [KG07]

Domain	IPC-2008 Net-Benefit Track			Compiled Problems			
	Gamer	HSP <sub>P</sub> *	Mips-XXL	Gamer	HSP <sub>F</sub> *	HSP <sub>0</sub> *	Mips-XXL
crewplanning(30)	4	16	8	-	8	21	8
elevators (30)	11	5	4	18	8	8	3
openstacks (30)	7	5	2	6	4	6	1
pegsol (30)	24	0	23	22	26	14	22
transport (30)	12	12	9	-	15	15	9
woodworking (30)	13	11	9	-	23	22	7
total	71	49	55		84	86	50

## Planning without observability: conformant planning

- ▶ Here we consider the second special case of planning with partial observability: planning without observability.
- ▶ Plans are **sequences of actions** because observations are not possible, actions cannot depend on the nondeterministic events or uncertain initial state, and hence the same actions have to be taken no matter what happens.
- ▶ Techniques needed for planning without observability can often be generalized to the general partially observable case.

## Why acting without observation?

- ▶ Conformant planning is like planning to act in an environment while you are blind and deaf.
- ▶ Observations could be **expensive** or it is preferable to have a **simple** plan.
- ▶ Example: Finding **synchronization sequences** in hardware circuits
- ▶ Example: **Initializing a system** consisting of many components that are in unknown states.
- ▶ **Internal** motivation: try to understand the **unobservable case** so that one can better deal with the more complicated **partially observable case**.

## Belief states and the belief space

- ▶ The current state is not in general known during plan execution. Instead, a **set of possible current states** is known.
- ▶ The set of possible current states forms the **belief state**.
- ▶ The set of all belief states is the **belief space**.
- ▶ If there are  $n$  states and none of them can be observationally distinguished from another, then there are  $2^n - 1$  belief states.

# The belief space

1. Let  $B$  be a **belief state** (a set of states).
2. Operator  $o$  is executable in  $B$  if it is executable in **every**  $s \in B$ .
3. When  $o$  is executed, possible next states are  $T = \text{img}_o(B)$ .
4. Belief states can be succinctly represented using Boolean formulae or BDDs.



# The belief space

## Example

### Example (Next slide)

Belief space generated by states over two Boolean state variables.

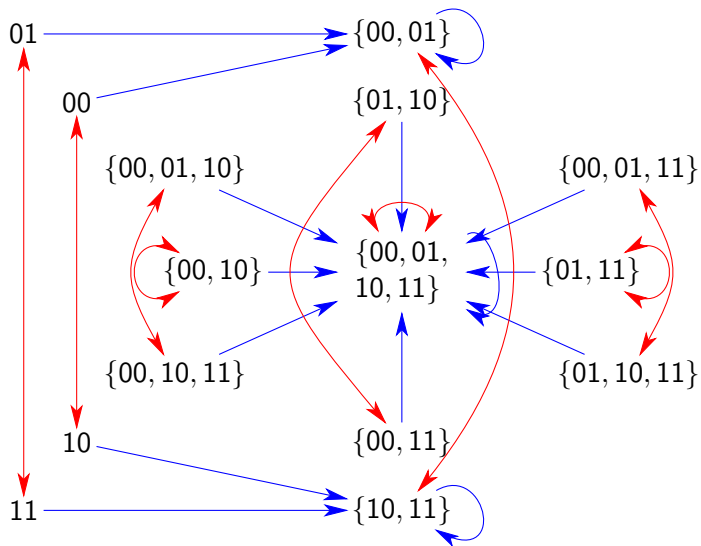
$n = 2$  state variables,  $2^n = 4$  states,  $2^{2^n} - 1 = 15$  belief states

**red action:** complement the value of the first state variable

**blue action:** assign a random value to the second state variable

# The belief space

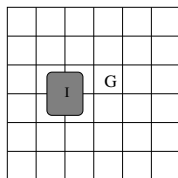
## Example



# Algorithms for unobservable problems

1. Find an operator sequence  $o_1, \dots, o_n$  that reaches a state satisfying  $G$  starting from **any** state satisfying  $I$ .
2.  $o_1$  must be applicable in all states  $B_0 = \{s \in S \mid s \models I\}$  satisfying  $I$ .  
 $o_2$  must be applicable in all states in  $B_1 = \text{img}_{o_1}(B_0)$ .  
 $o_i$  must be applicable in all states in  $B_i = \text{img}_{o_i}(B_{i-1})$  for all  $i \in \{1, \dots, n\}$ .  
 Terminal states must be goal states:  $B_n \subseteq \{s \in S \mid s \models G\}$ .

# Conformant vs. Classical Planning



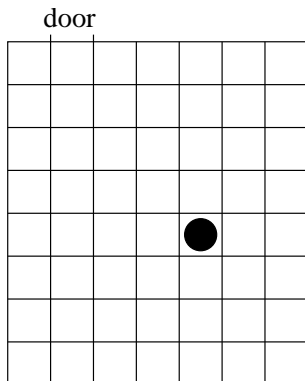
**Problem:** A robot must move from an **uncertain**  $I$  into  $G$  with **certainty**, one cell at a time, in a grid  $n \times n$

- ▶ Conformant and classical planning look similar except for uncertain  $I$  (assuming actions are deterministic).
- ▶ Yet plans can be quite different:  
best **conformant plan** **must** move robot to a corner first! (in order to localize)

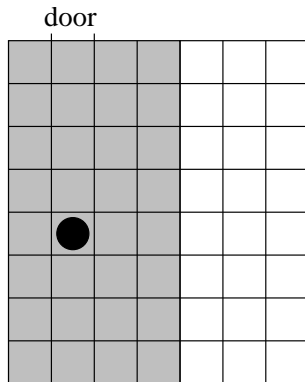
# The belief space

## Example

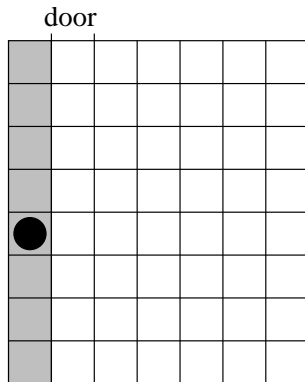
- ▶ A robot without any sensors, anywhere in a room of size  $7 \times 8$ .
- ▶ Actions: go North, South, East, West; if no way, just stay where you are
- ▶ Plan for getting out:  $6 \times$  West,  $7 \times$  North,  $1 \times$  East,  $1 \times$  North
- ▶ On the next slides we depict one possible location of the robot (●) and the change in the belief state at every execution step by gray fields.



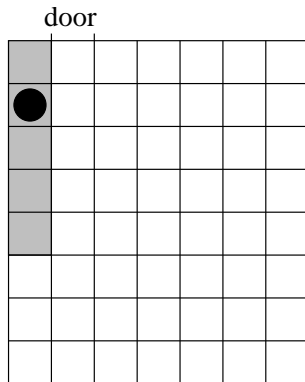
# Example: after WWW



Example: after WWWWWW

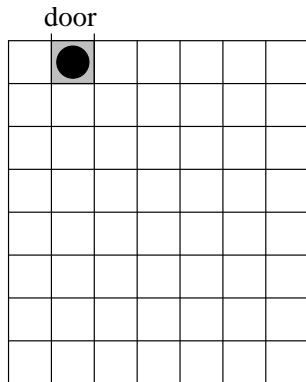


Example: after WWWWNNNN





Example: after WWWWWWNNNNNNNE



# Empirical Troubles with Conformant Planning

## Problems with top-down approach

- ▶ **effective representation** of belief states  $b$
- ▶ **effective heuristic**  $h(b)$  for estimating cost from  $b$  to  $b_G$

Now show: both tackled by **translation** into classical planning!

# Complexity: Classical vs. Conformant Planning

- ▶ **Complexity:** conformant planning harder than classical planning
  - ▶ *because **verification** of a conformant plan **intractable** in worst case*
- ▶ **Idea:** focus on computation of conformant plans that are **easy to verify** (e.g., in linear time in the plan length)
  - ▶ *computation of such plans **no more complex than classical planning***

## Basic Translation: Move to Knowledge Level

Given **conformant** problem  $\Pi = \langle P, I, O, G \rangle$

- ▶  $P$  – set of (all unobservable) *propositional* state variables
- ▶  $O$  – set of operators with conditional effects  $\langle c, e \rangle$
- ▶  $I$  – *prior knowledge* about the initial state (clauses over  $P$ )
- ▶  $G$  – goal description (conjunction over  $P$ )

Define **classical** problem  $K_0(\Pi) = \langle P', I', O', G' \rangle$

- ▶  $P' = \{Kp, K\neg p \mid p \in P\}$
- ▶  $I' = \{Kp \mid \text{clause } p \in I\}$
- ▶  $G' = \{Kp \mid p \in G\}$
- ▶  $O' = O$  but preconds  $p$  replaced by  $Kp$ , and effects  $\langle c, e \rangle$  replaced by  $Kc \rightarrow Ke$  (**supports**) and  $\neg K\neg c \rightarrow \neg K\neg e$  (**cancellation**)

$K_0(\Pi)$  is **sound** but **incomplete**: every classical plan that solves  $K_0(\Pi)$  is a conformant plan for  $\Pi$ , but not vice versa.

## Basic Translation: Move to Knowledge Level

<b>Conformant</b> $\Pi$	$\Rightarrow$	<b>Classical</b> $K_0(\Pi)$
$\langle P, I, O, G \rangle$	$\Rightarrow$	$\langle P', I', O', G' \rangle$
variable $p$	$\Rightarrow$	$Kp, K\neg p$ (two vars)
<b>Init:</b> known var $p$	$\Rightarrow$	$Kp \wedge \neg K\neg p$
<b>Init</b> unknown var $p$	$\Rightarrow$	$\neg Kp \wedge \neg K\neg p$ (both false)
<b>Goal</b> $p$	$\Rightarrow$	$Kp$
<b>Operator</b> $a$ has prec $p$	$\Rightarrow$	$a$ has prec $Kp$
<b>Operator</b> $a: \langle c, p \rangle$	$\Rightarrow$	$\left\{ \begin{array}{l} a : Kc \rightarrow Kp \\ a : K\neg c \rightarrow \emptyset \\ a : \neg K\neg c \rightarrow \neg K\neg p \end{array} \right.$

# Basic Properties and Extensions

- ▶ Translation  $K_0(\Pi)$  is **sound**:
  - ▶ If  $\pi$  is a **classical plan** that solves  $K_0(\Pi)$ , then  $\pi$  is a **conformant plan** for  $\Pi$ .
- ▶ But way **too incomplete**
  - ▶ often  $K_0(\Pi)$  will have no solution while  $\Pi$  does
  - ▶ works when **uncertainty is irrelevant**
- ▶ Extension  $K_{T,M}(\Pi)$  we present now **can** be both **complete and polynomial**

## Idea

- ▶ Given literal  $L$  and tag  $t$ , atom  $KL/t$  means
  - ▶  $K(t_0 \supset L)$ :  $KL$  true if  $t$  is true **initially**

## Example

- ▶ Conformant Problem  $\Pi$ :
  - ▶ Init:  $x_1 \vee x_2, \neg g$
  - ▶ Goal:  $g$
  - ▶ Actions:  $a_1 : x_1 \rightarrow g, a_2 : x_2 \rightarrow g$
- ▶ Classical Problem  $K_{T,M}(\Pi)$ :
  - ▶ Init:  $Kx_1/x_1, Kx_2/x_2, K\neg g, \neg Kg, \neg Kx_1, \neg K\neg x_1, \dots$
  - ▶ After  $a_1$ :  $Kg/x_1, Kx_1/x_1, Kx_2/x_2, \neg K\neg g, \neg Kg, \dots$
  - ▶ After  $a_2$ :  $Kg/x_2, Kg/x_1, Kx_1/x_2, Kx_2/x_2, \neg K\neg g, \neg Kg, \dots$ 
    - ▶ New action  $merge_g$ :  $Kg/x_1 \wedge Kg/x_2 \rightarrow Kg$
  - ▶ After  $merge_g$ :  $Kg, Kg/x_2, Kg/x_1, Kx_1/x_2, Kx_2/x_2, \neg K\neg g, \dots$
  - ▶ Goal satisfied:  $Kg$

Key elements in Translation  $K_{T,M}(\Pi)$ 

- ▶ a set  $T$  of **tags**  $t$ : consistent set of **assumptions** (literals) about the **initial situation**  $I$

$$I \not\models \neg t$$

- ▶ a set  $M$  of **merges**  $m$ : **valid subsets of tags**

$$I \models \bigvee_{L \in m} L$$

- ▶ Semantics of var  $KL/t$ :  $L$  is true given that initially  $t$  (i.e.  $K(t_0 \supset L)$ )



## Example of $T, M$

### Example

Given  $I = \{p \vee q, v \vee \neg w\}$ ,  $T$  and  $M$  can be:

$$\begin{array}{ll}
 T & = \{\{\}, p, q, v, \neg w\} & T' & = \{\{\}, \{p, v\}, \{q, v\}, \dots\} \\
 M & = \{\{p, q\}, \{v, \neg w\}\} & M' & = \dots
 \end{array}$$

## Translation $K_{T,M}(\Pi)$

For conformant  $\langle P, I, O, G \rangle$ ,  $K_{T,M}(\Pi)$  is  $\langle P', I', O', G' \rangle$

- ▶ **P'**:  $KL/t$  for every literal  $L$  and tag  $t \in T$
- ▶ **I'**:  $KL/t$  if  $I \models (t \supset L)$
- ▶ **G'**:  $KL$  for  $L \in G$
- ▶ For every tag  $t$  in  $T$  and  $a : L_1 \wedge \dots \wedge L_n \rightarrow L$  in  $O$ , add to  $O'$ 
  - ▶  $a : KL_1/t \wedge \dots \wedge KL_n/t \rightarrow KL/t$
  - ▶  $a : \neg K \neg L_1/t \wedge \dots \wedge \neg K \neg L_n/t \rightarrow \neg K \neg L/t$
- ▶ **prec**  $L \Rightarrow$  **prec**  $KL$
- ▶ **Merge** actions in  $O'$ : for each lit  $L$  and merge  $m \in M$  with  $m = \{t_1, \dots, t_n\}$

$$\text{merge}_{L,m} : KL/t_1 \wedge \dots \wedge KL/t_n \rightarrow KL$$

# Properties of Translation $K_{T,M}$

- ▶ If  $T$  contains only the empty tag,  $K_{T,M}(\Pi)$  reduces to  $K_0(\Pi)$
- ▶  $K_{T,M}(\Pi)$  is always **sound**

We will see that...

- ▶ For suitable choices of  $T, M$  translation is **complete**
- ▶ ...and sometimes **polynomial** as well

## Intuition of soundness

- ▶ Idea:
  - ▶ if sequence of actions  $\pi$  makes  $KL/t$  true in  $K_{T,M}(\Pi)$
  - ▶  $\pi$  makes  $L$  true in  $\Pi$  over all **trajectories** starting at initial states satisfying  $t$

### Theorem (Soundness $K_{T,M}(\Pi)$ )

*If  $\pi$  is a **plan that solves the classical planning problem**  $K_{T,M}(\Pi)$ , then the action sequence  $\pi'$  that results from  $\pi$  by dropping the merge actions is a **plan that solves the conformant planning problem**  $\Pi$ .*

## A complete but exponential instance of $K_{T,M}(\Pi)$ : $K_{s_0}$

If possible initial states are  $s_0^1, \dots, s_0^n$ , scheme  $K_{s_0}$  is the instance of  $K_{T,M}(\Pi)$  with

- ▶  $T = \{ \{\}, s_0^1, \dots, s_0^n \}$
- ▶  $M = \{ \{s_0^1, \dots, s_0^n\} \}$   
*i.e.*, only **one merge** for the disjunction of possible initial states
- ▶ **Intuition**: applying actions in  $K_{s_0}$  keeps track of each fluent for each possible initial states
- ▶ This instance is **complete**, but exponential in the number of fluents
  - ▶ ...although not a bad conformant planner

Performance of  $K_{S_0} + FF$ 

Problem	# $S_0$	Planners exec time (s)			
		$K_{S_0}$	$KP$	POND	CFF
Bomb-10-1	1k	648,9	0	1	0
Bomb-10-5	1k	2795,4	0,1	3	0
Bomb-10-10	1k	5568,4	0,1	8	0
Bomb-20-1	1M	> 1.8G	0,1	4139	0
Sqr-4-16	4	0,3	fail	1131	13,1
Sqr-4-24	4	1,6	fail	> 2h	321
Sqr-4-48	4	57,5	fail	> 2h	> 2h
Sortnet-6	64	2,2	fail	2,1	fail
Sortnet-7	128	27,9	fail	17,98	fail
Sortnet-8	256	> 1.8G	fail	907,1	fail

Translation time included in all tables.