## Adversarial Search

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## Games, man vs. algorithm

- Deep Blue
- Alpha Go
- Deep Stack
- Why Games, actually?


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Games are interesting for AI because they are hard (to solve).

## More: Adversarial Learning



Video: Adversing visual segmentation
Vision for Robotics and Autonomous Systems, http://cyber.felk.cvut.cz/vras

## Elements of the game



Considering the notation, we are making slight transition from [1] to [2].

- Players: $P=\{1,2, \ldots, N\}$ (often just $N=2$ )
- Transition functions: $S \times A \rightarrow S$.
- Terminal utilities: $S \times P \rightarrow R .(R-$ as a Reward $)$

What are we loking for? A strategy/policy $S \rightarrow A$

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- Player(s). Which player has to move in $s$.

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- Result( $s, a$ ). Transition, result of a move.

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- terminal-test(s). Game over?
- terminal-Utility $(s, p)$. What is prize? Examples for some games ..

https://commons.wikimedia.org/ wiki/File:Tic-tac-toe_5.png

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- Zero-sum: players have opposite utilities (values)
- Zero-sum: playing against
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- Zero-sum: playing against
- General game: independent utilities
- General game: cooperations, competition,


## Game Tree(s)



## State Value $V(s)$

Think about the State Value. It is a theoretical construct, definition. Depending on the problem, there may be various computational algorithms In a game, what State Values are known? Usually, only terminal states.
Think, for a moment, you are the only player. You can control every step. How would you compute the $V(s)$ for a given state $s$ ?

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Two-ply game: max for me, min for the opponent.

I'm player that starts (state A) and want to decide what to play, actions/plays/moves $a_{1}, a_{2}, a_{3}$ are the options. $\mathrm{B}, \mathrm{C}, \mathrm{D}$ are the possible outcomes of my moves. Now the opponent is about to play. The numbers in terminal states denote my profit/utility.

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## Zero-Sum game: max for me, min for the opponent.

Max step: I want to maximize my outcome.
Min step: I want to minimize the outcome of the opponent.

MIN (o)

MAX (x)


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$\operatorname{MAX}(\mathbf{x})$

$\operatorname{MinimAX}(s)=$
$\operatorname{UTILITY}(s)$ if TERMINAL-TEST( $s$ )
max $\operatorname{MinimAx}(\operatorname{RESULT}(s, a))$ if $\operatorname{PLAYER}(s)=$ MAX $a \in \operatorname{ACTIONS}(s)$

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## Minimax algorithm

function MINIMAX(state) returns an action
function MIN-VALUE(state) returns a utility value $v$

## function MAX-VALUE(state) returns a utility value $v$

## Minimax algorithm

function minimax(state) returns an action
return argmax min-VALUE(Result(state, a))
$a \in$ Actions(s)
end function
function MIN-VALUE(state) returns a utility value $v$

## function MAX-VALUE(state) returns a utility value $v$

## Minimax algorithm

function MINIMAX(state) returns an action
return argmax min-VALUE(ReSult(state, a))
$a \in$ Actions(s)

## end function

function MIN-VALUE(state) returns a utility value $v$ if TERMINAL-TEST(state) then return UTILITY(state) end if
$v \leftarrow \infty$ for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m i n }}(v$, MAX-VALUE(RESULT$($ state,$a)))$ end for
end function
function MAX-VALUE(state) returns a utility value $v$

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function MINIMAX(state) returns an action
return $\begin{aligned} & \text { argmax } \operatorname{MIN-VALUE}(\operatorname{RESULT}(\text { state }, ~ a))\end{aligned}$
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## end function

```
function MAX-VALUE(state) returns a utility value v
    if TERMINAL-TEST(state) then return UTILITY(state)
    end if
    v}\leftarrow-
    for all ACTIONS(state) do
            v\leftarrow\boldsymbol{max}(v, MIN-VALUE(RESULT(state,a)))
    end for
end function
```

A two ply game, down to terminal and back again
function MINIMAX $(s)$ returns a $\operatorname{argmax} \operatorname{MINVAL}(\operatorname{RES}(s, a))$

$$
a \in \operatorname{Actions}(s)
$$

end function

$$
\text { function MINVAL(s) returns } v
$$

if TERMINAL( $s$ ) then UTIL( $s$ ) end if
$v \leftarrow \infty$
for all ACTIONS( $s$ ) do
$v \leftarrow \min (v, \operatorname{MAXVAL}(\operatorname{RES}(s, a)))$
end for
end function
function MAXVAL( $s$ ) returns $v$
if TERMINAL( $s$ ) then $\operatorname{UTIL}(s)$
end if
$v \leftarrow-\infty$
for all ACTIONS(s) do
$v \leftarrow \max (v, \operatorname{MiNVAL}(\operatorname{RES}(s, a)))$
end for end function

Before going to the animation on the next slide, try to follow the algorithm by a pencil and paper.

A two ply game, recursive run


Efficiency/complexity:

- Exhaustive DFS
- Time $O\left(b^{m}\right)$
- Space $O(b m)$

Chess $b \approx 35, m \approx 100 \ldots$

- We cannot go(dive) to the end
- Can we save something?


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What is the complexity? How many nodes to visit?

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## Can we do better? How?

Nodes (sub-trees) worth visiting

A

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$$
\langle-(-\infty, \infty\rangle
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## $\alpha-\beta$ prunning

$\alpha$ highest (best) value choice found so far for any choice along MAX $\beta$ lowest (best) value choice found so far for any choice along Min $v$ value of the state

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Functions scope: MAX-VALUE MIN-VALUE
In MAX nodes $\alpha$ is changing and $\beta$ is stopping, in MIN nodes $\beta$ is changing and $\alpha$ is stopping.
It is clear that ordering of child nodes matters. Draw tree of $\alpha-\beta$ search in case of perferct ordering. Effective branching factor becomes $\sqrt{b}$ instead of $b$ which effectively doubles the depth can be searched.

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$\alpha=-\infty, \beta=\infty, v=?$

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In MIN-VAL: $v \leftarrow 2$ $v \leq \alpha$ then: return $v$ !

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In MAX nodes $\alpha$ is changing and $\beta$ is stopping, in MIN nodes $\beta$ is changing and $\alpha$ is stopping.
It is clear that ordering of child nodes matters. Draw tree of $\alpha-\beta$ search in case of perferct ordering. Effective branching factor becomes $\sqrt{b}$ instead of $b$ which effectively doubles the depth can be searched.

Take the tree from the previous slide and try to go step-by-step, watch $\alpha$
function ALPHA-BETA-SEARCH(state) returns an action
$v \leftarrow \operatorname{MAX}-\operatorname{VALUE}($ state, $\alpha=-\infty, \beta=\infty)$
return action corresponding to $v$

## end function

function MAX-VALUE(state, $\alpha, \beta$ ) returns a utility value $v$
if TERMINAL-TEST(state) return UTILITY(state)
$v \leftarrow-\infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m a x }}(v, \operatorname{MIN}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a), \alpha, \beta))$
if $v \geq \beta$ return $v$
$\alpha \leftarrow \max (\alpha, v)$
end for end function

Take the tree from the previous slide and try to go step-by-step, watch $\alpha$, $\beta$ and $v$
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if $v \geq \beta$ return $v$
$\alpha \leftarrow \max (\alpha, v)$
end for
end function
function MIN-VALUE(state, $\alpha, \beta$ ) returns a utility value $v$
if TERMINAL-TEST(state) return UTILITY(state)
$v \leftarrow \infty$
for all ACTIONS(state) do
$v \leftarrow \boldsymbol{\operatorname { m i n }}(v, \operatorname{MAX}-\operatorname{VALUE}(\operatorname{RESULT}($ state,$a), \alpha, \beta))$
if $v \leq \alpha$ return $v$
$\beta \leftarrow \min (\beta, v)$
end for
end function

Take the tree from the previous slide and try to go step-by-step, watch $\alpha$ $\beta$ and $v$

[^0]```
H-MinimAX (s,d)=
    EVAL(s) if CUTOFF-TEST}(s,d
```

```
        H-Minimax (s,d) =
        EVAL(s) if CUTOFF-TEST(s,d)
        max H-MinimAx(RESULT}(s,a),d+1) if PlayER(s) = MAX
ma_ACTIONS(s)
```

```
        H-MinimAX (s,d)=
        EVAL(s) if CUTOFF-TEST(s,d)
    max ctions(s)}\operatorname{H-MINIMAX(RESULT}(s,a),d+1) if PLAYER(s)= MAX
min
a\inACTIONS(s)
```


## Replace

if TERMINAL-TEST(s) then return UTILITY(s)
with:
if Cutoff-TEst( $(\mathrm{s}, \mathrm{d})$ then return $\operatorname{EvaL}(\mathrm{s})$

EVAL(s) - Evaluation functions
(estimate of) State value for non-terminal states

(a) White to move

(b) White to move

For many problems it is not so easy to find/construct proper function. We may try more functions and combine them conveniently.
$f_{1}(s)=$ number of white pawns - number of black pawns
How to tune weights $w_{i}$ ?
or Deep Nets! Yeah!
How the get training data for supervised learning? More later.

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$$
\operatorname{EvaL}(s)=w_{1} f_{1}(s)+w_{2} f_{2}(s)+\cdots w_{n} f_{n}(s)
$$

## References

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[^0]:    h-minimax $(s, d)=$

