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**DCGI**

**DEPARTMENT OF COMPUTER GRAPHICS AND INTERACTION**

# CONVEX HULLS

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<https://cw.felk.cvut.cz/doku.php/courses/a4m39vg/start>

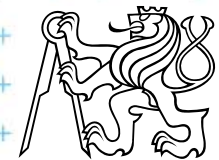
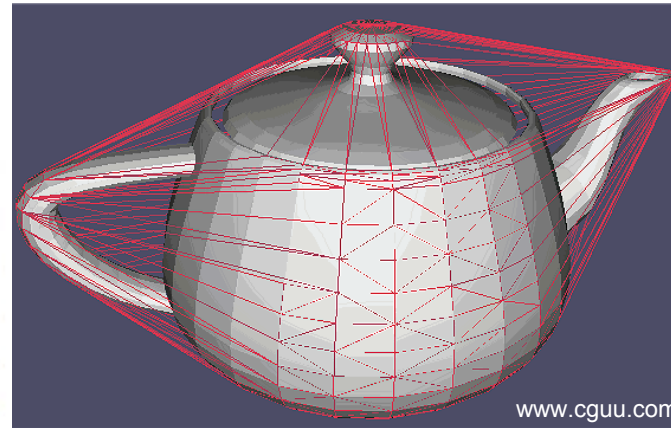
Based on [Berg] and [Mount]

Version from 8.11.2012

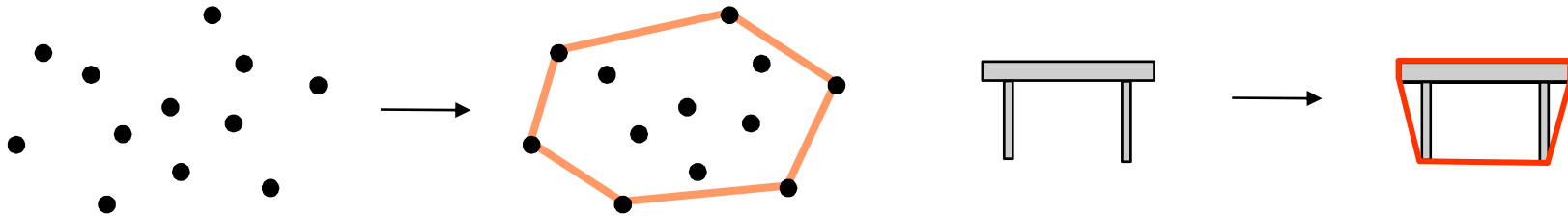
# Talk overview

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- Motivation and Definitions
- Graham's scan – incremental algorithm
- Divide & Conquer
- Quick hull
- Jarvis's March – selection by gift wrapping
- Chan's algorithm – optimal algorithm

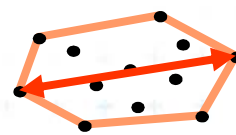


# Convex hull (CH) – why to deal with it?

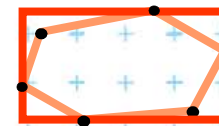
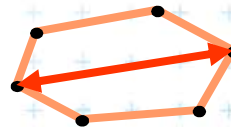


- *Shape approximation* of a point set or complex shapes (other common approximations include: minimal area enclosing rectangle, circle, and ellipse,...) – e.g., for collision detection
- *Initial stage* of many algorithms to filter out irrelevant points, e.g.:

– diameter of a point set



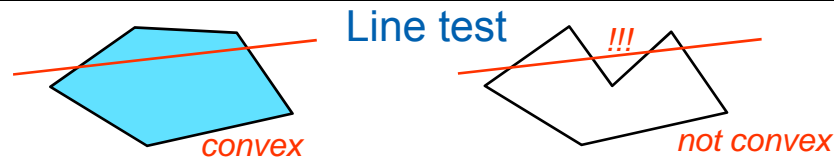
– minimum enclosing convex shapes (such as rectangle, circle, and ellipse) depend only on points on CH



# Convexity

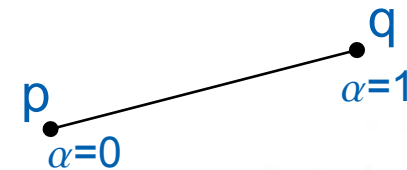
- A set  $S$  is *convex*

- if for any points  $p, q \in S$  the line segment  $\overline{pq} \subseteq S$ , or
- if any convex combination of  $p$  and  $q$  is in  $S$



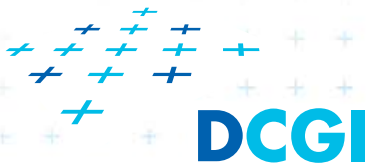
- *Convex combination* of points  $p, q$  is any point that can be expressed as

$(1 - \alpha) p + \alpha q$ , where  $0 \leq \alpha \leq 1$



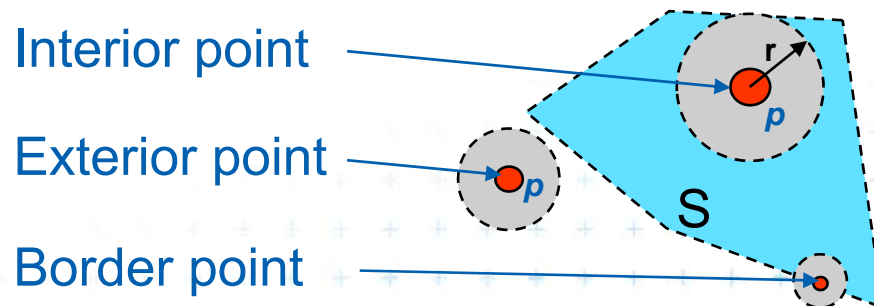
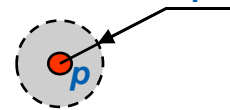
- *Convex hull*  $CH(S)$  of set  $S$  – is (similar definitions)

- the smallest set that contains  $S$
- or: intersection of all convex sets that contain  $S$
- Or in 2D for points: the smallest convex polygon containing all given points


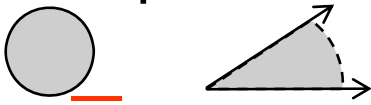




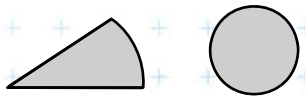
# Definitions from topology in metric spaces

- **Metric space** – each two of points have defined a **distance**  $r$
- **$r$ -neighborhood** of a point  $p$  and radius  $r > 0$   
= set of points whose **distance** to  $p$  is strictly less than  $r$   
(open ball of diameter  $r$  centered about  $p$ )
- Given set  $S$ , point  $p$  is
  - **Interior point** of  $S$  – if  $(r$ -neighborhood about  $p$  of radius  $r) \subset S$
  - **Exterior point** – if it lies in interior of the complement of  $S$
  - **Border point** – is neither interior neither exterior



# Definitions from topology in metric spaces

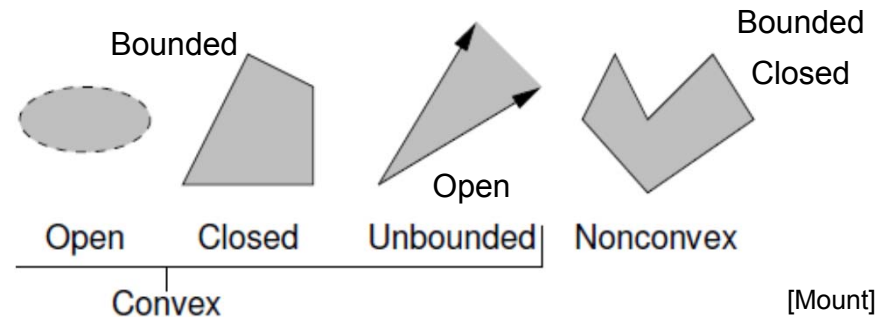
- Set S is Open** (otevřená) 
  - $\forall p \in S \exists (r\text{-neighborhood about } p \text{ of radius } r) \subseteq S$
  - it contains only interior points, none of its border points
- Closed** (uzavřená) 
  - If it is equal to its **closure**  $\overline{S}$  (uzávěr = smallest closed set containing S in topol. space)
  - $\forall (r\text{-neighborhood about } p \text{ of radius } r) \cap S \neq \emptyset$
- Clopen** (otevřená i uzavřená) – Ex. Empty set  $\emptyset$ , finite set of disjoint components
  - if it is both **closed** and **open**

space  $Q =$  rational numbers  
 (S= all positive rational numbers whose square is bigger than 2)  $S = (\sqrt{2}, \infty)$  in  $Q$ ,  $\sqrt{2} \notin Q$ ,  $S = \overline{S}$
- Bounded** (ohraničená)  **Unbounded** 
  - if it can be enclosed in a ball of finite radius
- Compact** (kompaktní) 
  - if it is both closed and bounded



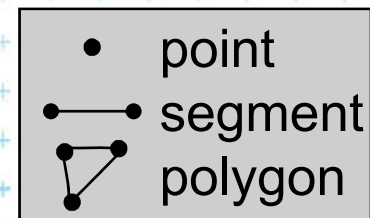
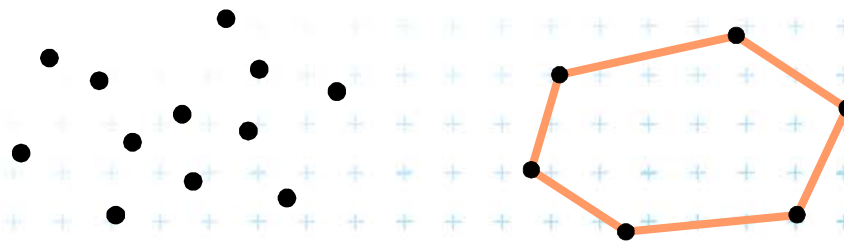
# Definitions from topology in metric spaces

- *Convex set  $S$  may be bounded or unbounded*



- *Convex hull  $CH(S)$  of a finite set  $S$  of points in the plane*

= Bounded, closed, (= compact) convex polygon

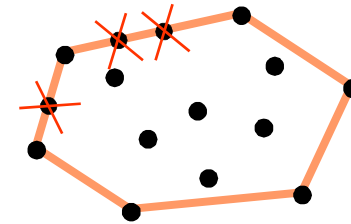




# Convex hull representation

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- CCW enumeration of vertices
- Contains only the extreme points (“endpoints” of collinear points)



- Simplification for this semester

Assume the input points are in **general position**,

- no two points have the same  $x$ -coordinates and
- no three points are collinear

-> We avoid problem with non-extreme points on  $x$   
(solution may be simple – e.g. lexicographic ordering)



# Online x offline algorithms


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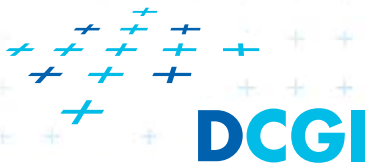
- **Incremental algorithm**
  - Proceeds one element at a time (step-by-step)
- **Online algorithm** (must be incremental)
  - is started on a partial (or empty) input and
  - continues its processing as additional input data becomes available (comes online, thus the name).
  - Ex.: insertion sort
- **Offline algorithm** (may be incremental)
  - requires the entire input data from the beginning
  - than it can start
  - Ex.: selection sort



# Graham's scan

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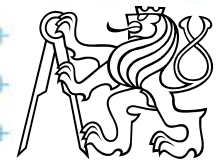
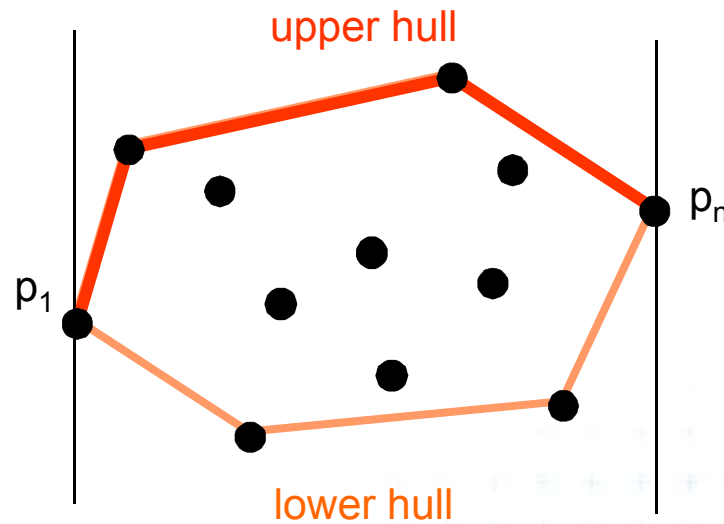
- Incremental  $O(n \log n)$  algorithm
- Objects (points) are added one at a time
- Order of insertion is important
  - Random insertion
    - we need to test: *is-point-inside-the-hull*( $p$ ) 
  - **Ordered insertion**  
**Sort points according to  $x$  and add them left to right** - it guarantees, that just added point is outside current hull
    - Original algorithm sorted the *angles* around the point with minimal  $y$
    - Sorting  *$x$ -coordinates* is simpler to implement than sorting of angles



# Graham's scan

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- $O(n \log n)$  for unsorted points,  $O(n)$  for sorted pts.
- Upper hull, then lower hull. Merge.
- Minimum and maximum on  $x$  belong to CH

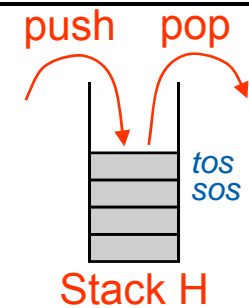


# Graham's scan – incremental algorithm

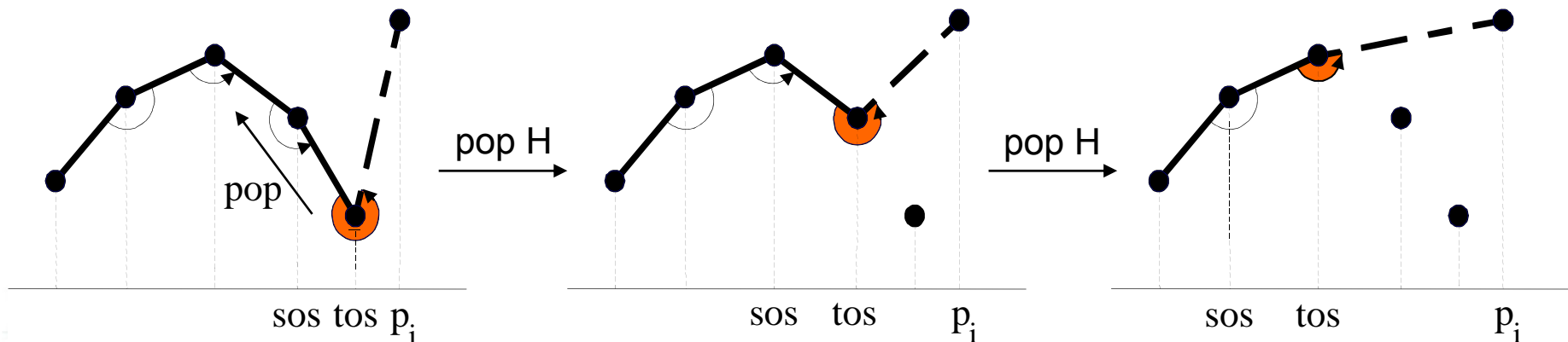
GrahamsScan(points p)

Input: points p

Output: CCW points on the convex hull



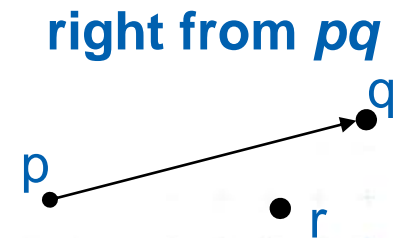
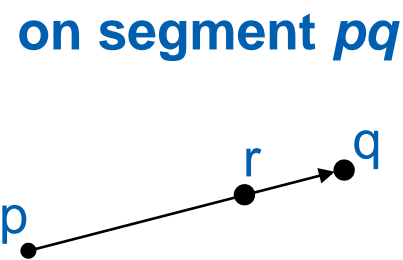
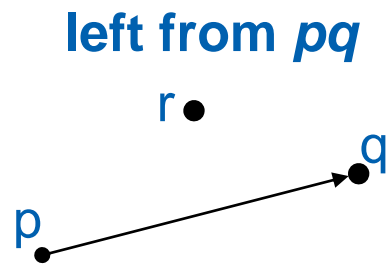
1. sort points according to increasing x-coord  $\rightarrow \{p_1, p_2, \dots, p_n\}$
2. push(  $p_1$ , H ), push(  $p_2$ , H )
3. **for**  $i = 3$  **to**  $n$  **do**
4.     **while**( size(H)  $\geq 2$  and orient(  $sos$ ,  $tos$ ,  $p_i$  )  $\geq 0$  ) // skip left turns
5.     pop H // (back-tracking)
6.     push(  $p_i$ , H ) // store right turn
7. store H to the output (in reverse order) // upper hull
8. Symmetrically the **lower hull**



# Position of point in relation to segment

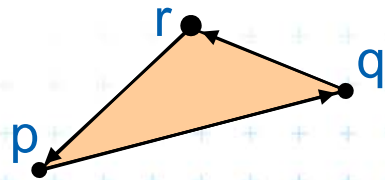
$\text{orient}(p, q, r) \begin{cases} > 0 & r \text{ is left from } pq, \text{ CCW orient} \\ = 0 & \text{if } (p, q, r) \text{ are collinear} \\ < 0 & r \text{ is right from } pq, \text{ CW orient} \end{cases}$

Point  $r$  is:



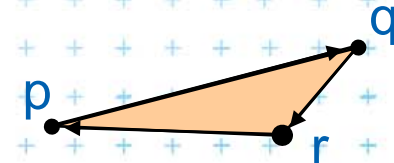
Convex polygon with edges  $pq$  and  $qr$  or

Triangle  $pqr$ : is CCW oriented



degenerated to line

is CW oriented



# Geometric meaning: Area of Triangle ABC

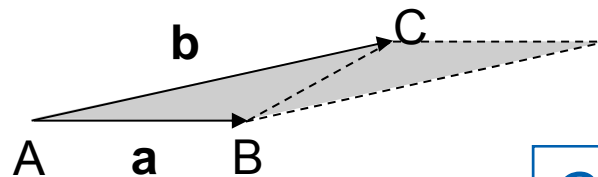
- Position of point C in relation to segment AB is given by the sign of the triangle ABC area 2x Oriented area

- $T = \frac{1}{2} |\vec{AB} \times \vec{AC}|$

- $\mathbf{a} = \mathbf{B} - \mathbf{A}$

- $\mathbf{b} = \mathbf{C} - \mathbf{A}$

- $T = \frac{1}{2} (\mathbf{a}_x \mathbf{b}_y - \mathbf{a}_y \mathbf{b}_x)$



Can be computed as vector product

$$\Rightarrow 2T = A_x B_y + B_x C_y + C_x A_y - A_x C_y - B_x A_y - C_x B_y$$

Or directly as determinant

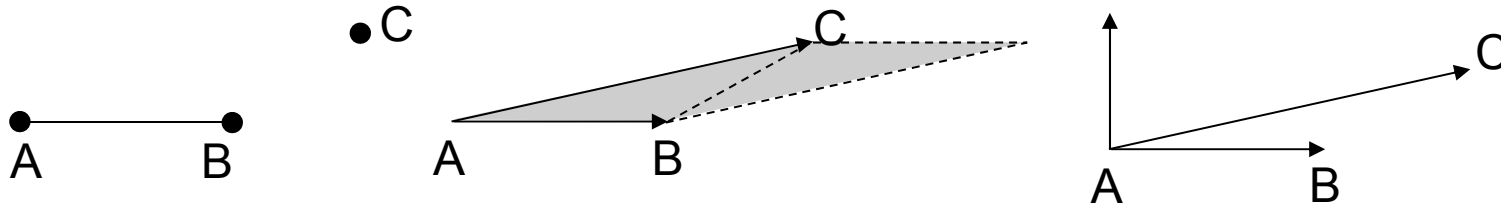
$$2T = \begin{vmatrix} A_x & A_y & 1 \\ B_x & B_y & 1 \\ C_x & C_y & 1 \end{vmatrix} = A_x B_y + B_x C_y + C_x A_y - A_x C_y - B_x A_y - C_x B_y$$





# Geometric meaning: Area of Triangle ABC

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Equal to size of **Vector product of vectors  $AB \times AC$**

= Vector perpendicular to both vectors  $AB$  and  $AC$

- If vectors in plane

- it is perpendicular to the **plane** (*normal vector of the plane*)

- only z-coordinate is non-zero

- $|\vec{AB} \times \vec{AC}|$  = z-coordinate of the normal vector

- = area of parallelogram

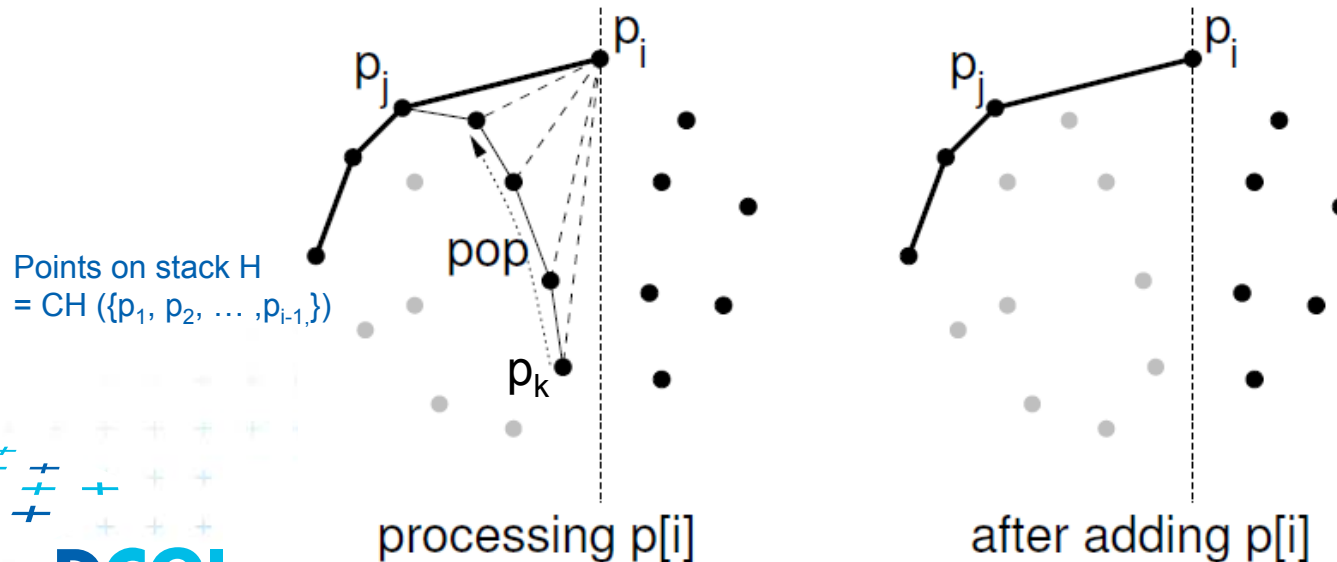
- = 2x area  $T$  of triangle  $ABC$





# Is Graham's scan correct?

- Stack H at any stage contains upper hull of the points  $\{p_1, \dots, p_j, p_i\}$ , processed so far
  - For induction basis  $H = \{p_1, p_2\}$  ... true
  - $p_i$  = last added point to CH,  $p_j$  = its predecessor on CH
  - Each point  $p_k$  that lies between  $p_j$  and  $p_i$  lies below  $p_j p_i$  and should not be part of UH after addition of  $p_i \Rightarrow$  is removed before push  $p_i$ .  
[  $\text{orient}(p_j, p_k, p_i) > 0$ ,  $p_i$  is left from  $p_j p_k \Rightarrow p_k$  is removed from UH ]
  - Stop if 2 points in the stack or after construction of the upper hull



[Mount]



# Complexity of Graham's scan

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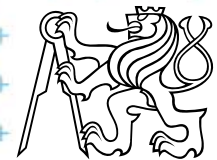
- Sorting according  $x$  –  $O(n \log n)$
- Each point pushed once –  $O(n)$
- Some ( $d_i \leq n$ ) points deleted while processing  $p_i$   
–  $O(n)$
- The same for lower hull –  $O(n)$
  
- Total  $O(n \log n)$  for unsorted points  
 $O(n)$  for sorted points



# Divide & Conquer

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- $\Theta(n \log(n))$  algorithm
- Extension of mergesort
- Principle
  - Sort points according to  $x$ -coordinate,
  - recursively partition the points and solve CH.



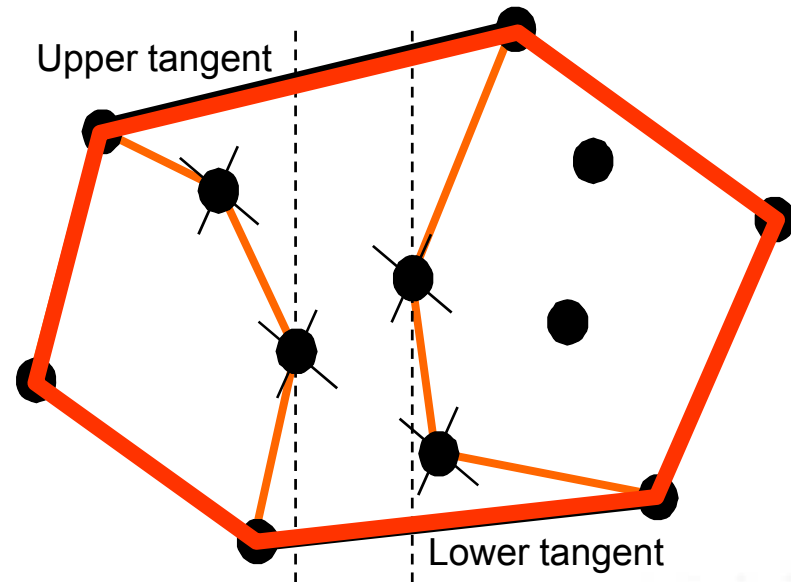
# Convex hull by D&C

## ConvexHullD&C( points P )

*Input:* points p

*Output:* CCW points on the convex hull

- Sort points P according to x
- return hull( P )
- hull( points P )**
- if  $|P| \leq 3$  then
- compute CH by brute force,
- return
- Partition P into two sets L and R (with lower & higher coords x)
- Recursively compute  $H_L = \text{hull}(L)$ ,  $H_R = \text{hull}(R)$
- $H = \text{Merge hulls}(H_L, H_R)$  by computing
- Upper\_tangent(  $H_L, H_R$  ) // find nearest points,  $H_L$  CCW,  $H_R$  CW
- Lower\_tangent(  $H_L, H_R$  ) // ( $H_L$  CW,  $H_R$  CCW)
- discard points between these two tangents
- return H



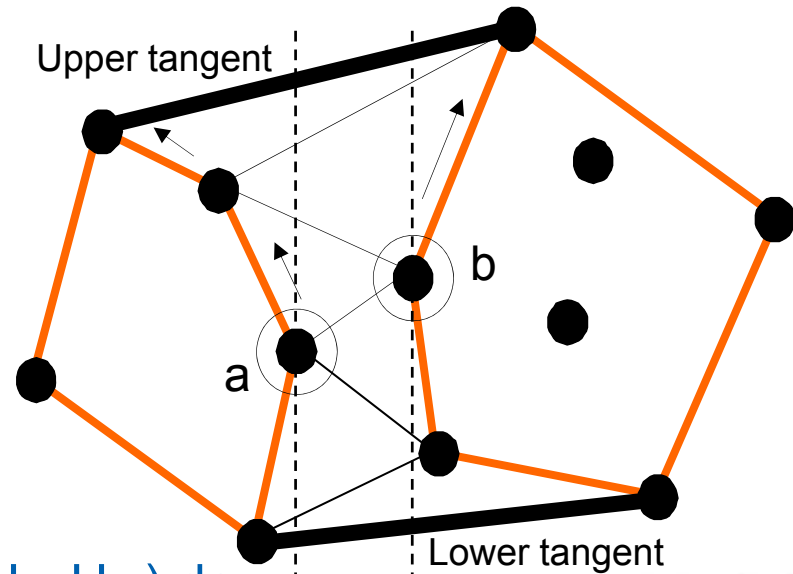
# Search for upper tangent (lower is symmetrical)

**Upper\_tangent**(  $H_L, H_R$  )

*Input:* two non-overlapping CH's

*Output:* upper tangent  $ab$

1.  $a = \text{rightmost } H_L$
2.  $b = \text{leftmost } H_R$
3. while(  $ab$  is not the upper tangent for  $H_L, H_R$  ) do
4.     while(  $ab$  is not the upper tangent for  $H_L$  )  $a = a.succ$  // move CCW
5.     while(  $ab$  is not the upper tangent for  $H_R$  )  $b = b.pred$  // move CW
6. Return  $ab$



Where: (  $ab$  is not the upper tangent for  $H_L$  )  $\Rightarrow \text{orient}(a, b, a.succ) \geq 0$   
 which means  $a.succ$  is **left from line  $ab$**

$$m = |H_L| + |H_R| \leq |L| + |R| \Rightarrow \text{Upper Tangent: } O(m) = O(n)$$



# Convex hull by D&C complexity

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- Initial sort  $O(n \log(n))$
- Function hull()
  - Upper and lower tangent  $O(n)$
  - Merge hulls  $O(1)$
  - Discard points between tangents  $O(n)$

- Overall complexity

- Recursion

$$T(n) = \begin{cases} 1 & \dots \text{ if } n \leq 3 \\ 2T(n/2) + O(n) & \dots \text{ otherwise} \end{cases}$$

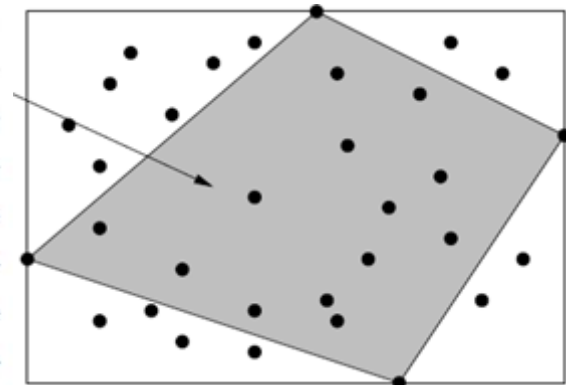
- Overall complexity of CH by D&C:  $\Rightarrow O(n \log(n))$



# Quick hull

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- A variant of Quick Sort
- $O(n \log n)$  expected time, max  $O(n^2)$
- Principle
  - in praxis, most of the points lie in the interior of CH
  - E.g., for uniformly distributed points in unit square, we expect only  $O(\log n)$  points on CH
- Find extreme points (parts of CH) quadrilateral, discard inner points
  - Add 4 edges to temp hull T
  - Process points outside 4 edges



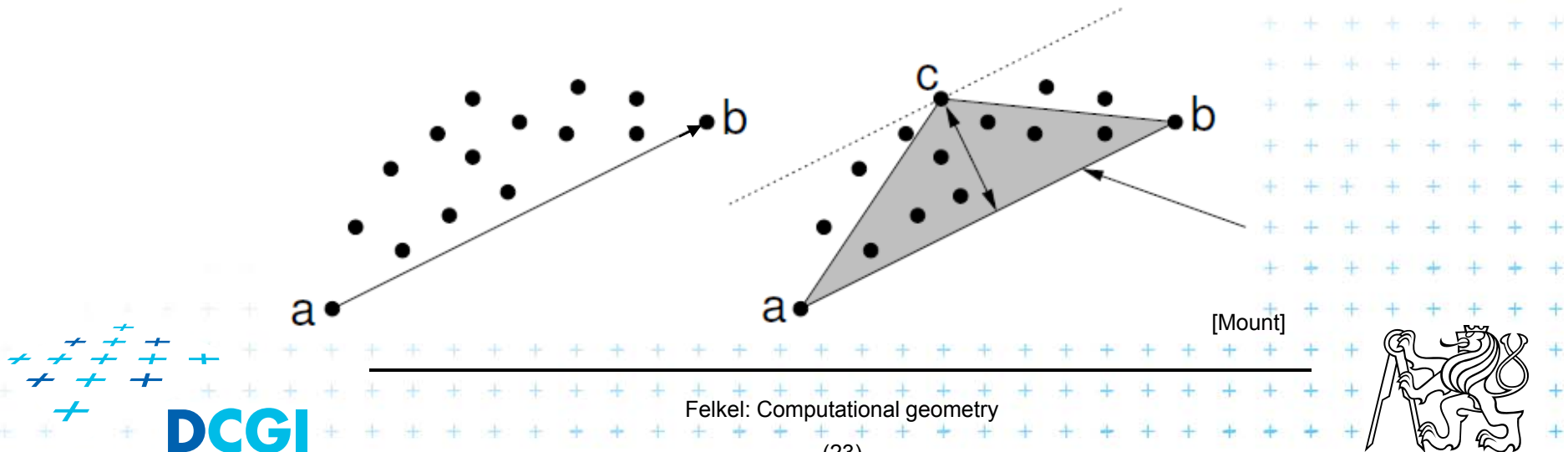
[Mount]





# Process each of four groups of points outside

- For points outside  $ab$  (left from  $ab$ )
  - Find point  $c$  on the hull – max. perpend. distance to  $ab$
  - Discard points inside triangle  $abc$  (right from the edges)
  - Split points into two subsets
    - outside  $ac$  (left from  $ac$ ) and outside  $cb$  (left from  $cb$ )
  - Process points outside  $ac$  and  $cb$  recursively
  - Replace edge  $ab$  in  $T$  by edges  $ac$  and  $cb$



[Mount]





# Quick hull complexity

---

- $n$  points remain outside the hull
- $T(n)$  = running time for such  $n$  points outside
  - $O(n)$  - selection of splitting point  $c$
  - $O(n)$  - point classification to inside &  $(n_1+n_2)$  outside
  - $n_1+n_2 \leq n$

– The running time is given by recurrence

$$T(n) = \begin{cases} 1 & \text{if } n = 1 \\ T(n_1) + T(n_2) & \text{where } n_1+n_2 \leq n \end{cases}$$

- If evenly distributed that  $\max(n_1, n_2) \leq \alpha n$ ,  $0 \leq \alpha \leq 1$  then solves as QuickSort to  $O(cn \log n)$  where  $c=f(\alpha)$  else  $O(n^2)$  for unbalanced splits



# Jarvis's March – selection by gift wrapping

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- Variant of  $O(n^2)$  selection sort
- Output sensitive algorithm
- $O(nh)$  ...  $h = \text{number of points on convex hull}$

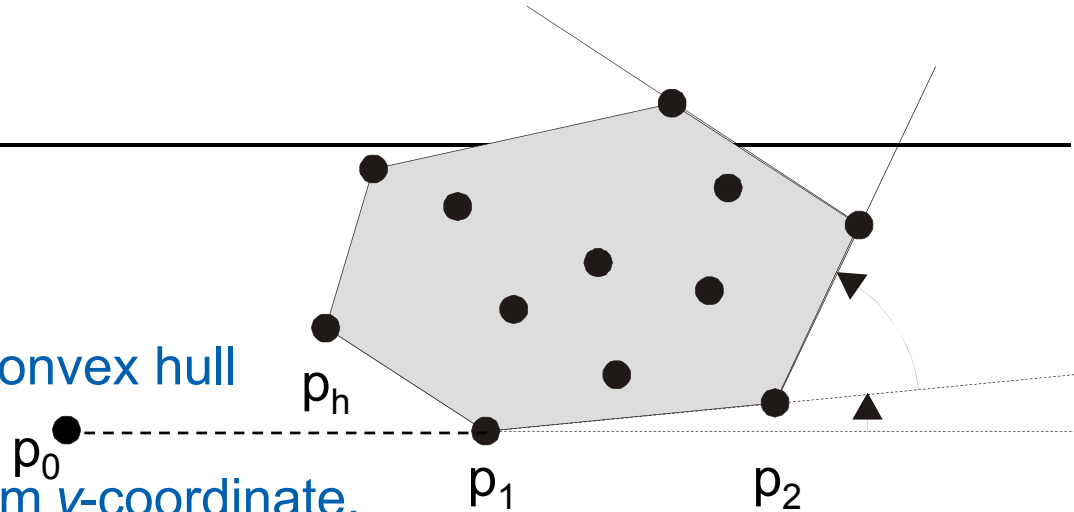


# Jarvis's March

**JarvisCH**(points P)

*Input:* points p

*Output:* CCW points on the convex hull

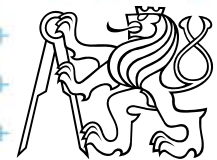


1. Take point  $p_1$  with minimum  $y$ -coordinate,  
//  $p_1$  will be the first point in the hull
2. Take a horizontal line, i.e., create temporary point  $p_0 = (-\infty, p_1.y)$
3.  $i = 1$
4. repeat
5. | Rotate the line around  $p_i$  until bounces to the nearest point  $q$   
| // compute the smallest angle by the smallest  $\text{orient}(p_{i-1}, p_i, q)$
6. |  $i++$   
|  $p_i =$  the bounced nearest point  $q$
7. until ( $q \neq p_1$ )

Output sensitive algorithm

Complexity:  $O(n) + O(n) * h \Rightarrow O(h * n)$

good for low number of points on convex hull



# Output sensitive algorithm

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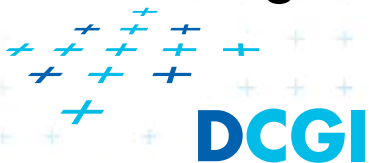
- Worst case complexity analysis analyzes the worst case data
  - Presumes, that **all (const fraction of) points lie on** the CH
  - The points are ordered along CH
    - => We need sorting =>  $\Omega(n \log n)$  of CH algorithm
- Such assumption is rare
  - usually only **much less of points are on CH**
- Output sensitive algorithms
  - Depend on: input size  $n$  and the size of the output  $h$
  - Are more efficient for small output sizes
  - Reasonable time for CH is  **$O(n \log h)$**



# Chan's algorithm

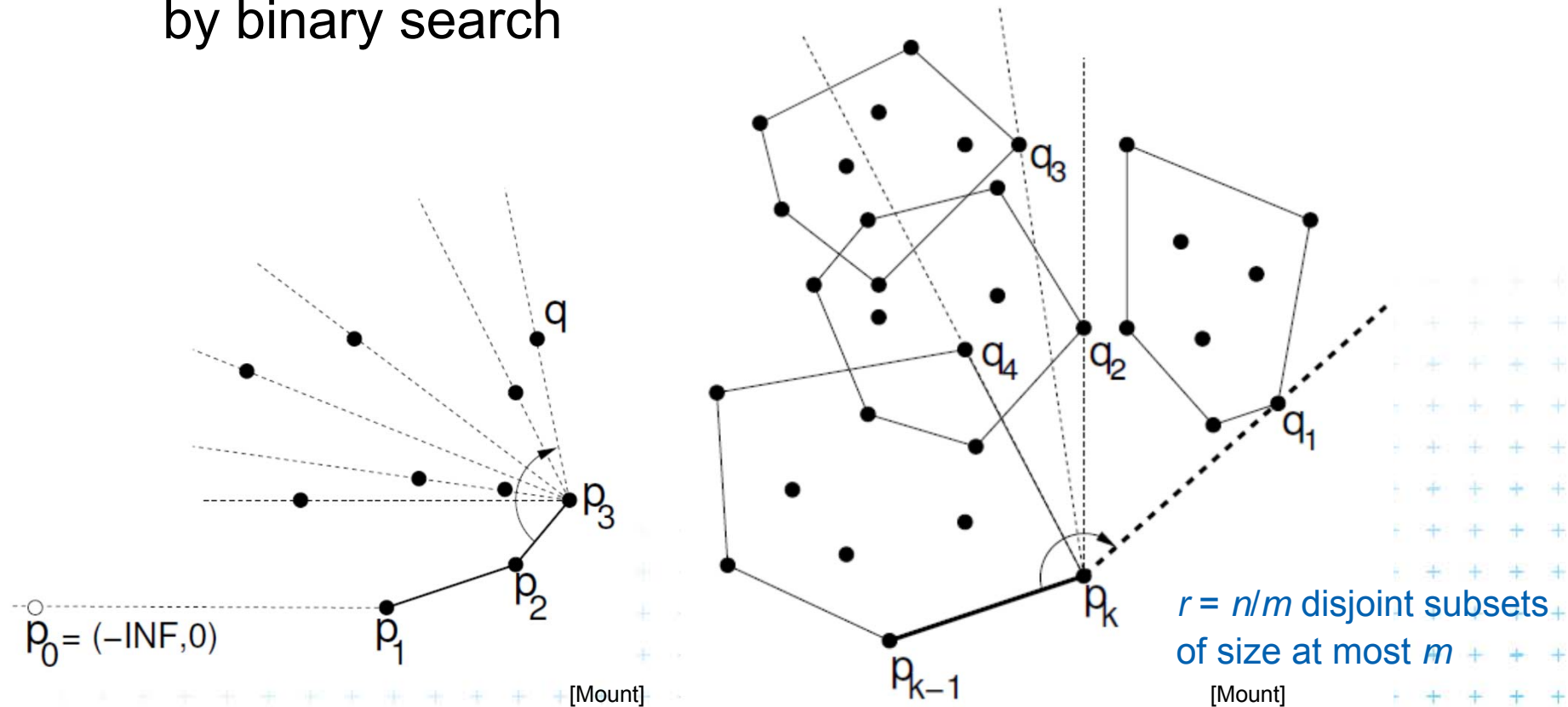
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- Cleverly combines Graham's scan and Jarvis's march algorithms
- Goal is  $O(n \log h)$  running time
  - We cannot afford sorting of all points -  $\Omega(n \log n)$
  - => Idea: limit the set sizes to polynomial  $h^c$   
the complexity does not change =>  $\log h^c = \log h$
  - $h$  is unknown – we get the estimation later
  - Use estimation  $m$ , better not too high =>  $h \leq m \leq h^2$
- Partition points  $P$  into  $r$ -groups of size  $m$ ,  $r = n/m$ 
  - Each group take  $O(m \log m)$  time - sort + Graham
  - $r$ -groups take  $O(r m \log m) = O(n \log m)$  - Jarvis



# Merging of $m$ parts in Chan's algorithm

- Merge  $r$ -group CHs as “fat points”
  - Tangents to convex  $m$ -gon can be found in  $O(\log m)$  by binary search



# Chan's algorithm complexity

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- $h$  points on the final convex hull

- ⇒ at most  $h$  steps in the Jarvis march algorithm
  - each step computes  $r$ -tangents,  $O(\log m)$  each
  - merging together  $O(hr \log m)$

$r$ -groups of size  $m$ ,  $r = n/m$

- Complete algorithm  $O(n \log h)$

- Graham's scan on partitions  $O(r \cdot m \log m) = O(n \log m)$
- Jarvis Merging:  $O(hr \log m) = O(h n/m \log m)$ , ...4a)  
 $h \leq m \leq h^2$   $= O(n \log m)$
- How to guess  $m$ ? *Wait!*



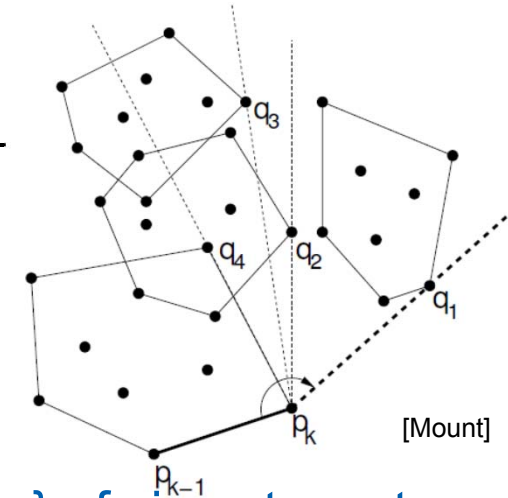


# Chan's algorithm for known $m$

PartialHull(  $P, m$  )

Input: points  $P$

Output: group of size  $m$



1. Partition  $P$  into  $r = \lceil n/m \rceil$  disjoint subsets  $\{p_1, p_2, \dots, p_r\}$  of size at most  $m$
2. for  $i=1$  to  $r$  do
  - a) Convex hull by GrahamsScan( $P_i$ ), store vertices in ordered array
3. let  $p_1$  = the bottom most point of  $P$  and  $p_0 = (-\infty, p_1.y)$
4. for  $k = 1$  to  $m$  do // compute merged hull points
  - a) for  $i = 1$  to  $r$  do // angle to all  $r$  subsets
 

Compute the point  $q_i \in P$  that maximizes the angle  $\angle p_{k-1}, p_k, q_i$
  - b) let  $p_{k+1}$  be the point  $q \in \{q_1, q_2, \dots, q_r\}$  that maximizes  $\angle p_{k-1}, p_k, q$   
 ( $p_{k+1}$  is the new point in CH)
  - c) if  $p_{k+1} = p_1$  then return  $\{p_1, p_2, \dots, p_k\}$
5. return "Fail,  $m$  was too small"

$O(\log m)$





# Chan's algorithm – estimation of $m$

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ChansHull

*Input:* points  $P$

*Output:* convex hull  $p_1 \dots p_k$

1. for  $t = 1, 2, \dots$  do {
  - a) let  $m = \min(2^{2^t}, n)$
  - b)  $L = \text{PartialHull}(P, m)$
  - c) if  $L \neq \text{"Fail, } m \text{ was too small"}$  then return  $L$}

Sequence of choices of  $m$  are  $\{ 4, 16, 256, \dots, 2^{2^t}, \dots, n \}$  ... squares

Example: for  $h = 23$  points on convex hull of  $n = 57$  points, the algorithm will try this sequence of choices of  $m$   $\{ 4, 16, 57 \}$

1. 4 and 16 will fail
2. 256 will be replaced by  $n$



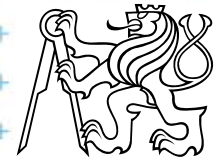
# Complexity of Chan's Convex Hull?

- The worst case: Compute all iterations
- $t^{\text{th}}$  iteration takes  $O(n \log 2^{2^t}) = O(n 2^t)$
- Algorithm stops when  $2^{2^t} \geq h \Rightarrow t = \lceil \lg \lg h \rceil$
- All  $t = \lceil \lg \lg h \rceil$  iterations take:

Using the fact that  $\sum_{i=0}^k 2^i = 2^{k+1} - 1$

$$\sum_{t=1}^{\lg \lg h} n 2^t = n \sum_{t=1}^{\lg \lg h} 2^t \leq n 2^{1+\lg \lg h} = 2n \lg h = O(n \log h)$$

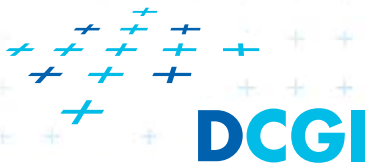
2x more work in the worst case



# Conclusion in 2D

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- Graham's scan:  $O(n \log n)$ ,  $O(n)$  for sorted pts
- Divide & Conquer:  $O(n \log n)$
- Quick hull:  $O(n \log n)$ , max  $O(n^2) \sim$  distrib.
- Jarvis's march:  $O(hn)$ , max  $O(n^2) \sim$  pts on CH
- Chan's alg.:  $O(n \log h) \sim$  pts on CH



# References

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