



**OPPA European Social Fund
Prague & EU: We invest in your future.**

PLÁNOVÁNÍ A HRY – CV 4

STATE-SPACE SEARCH

State-space Search

- **Forward Search**
- **Backward Search**
- **Heuristic Search**



Forward search

Forward Search

Forward-search(O, s_0, g)

$s \leftarrow s_0$

$\pi \leftarrow$ the empty plan

loop

if s satisfies g then return π

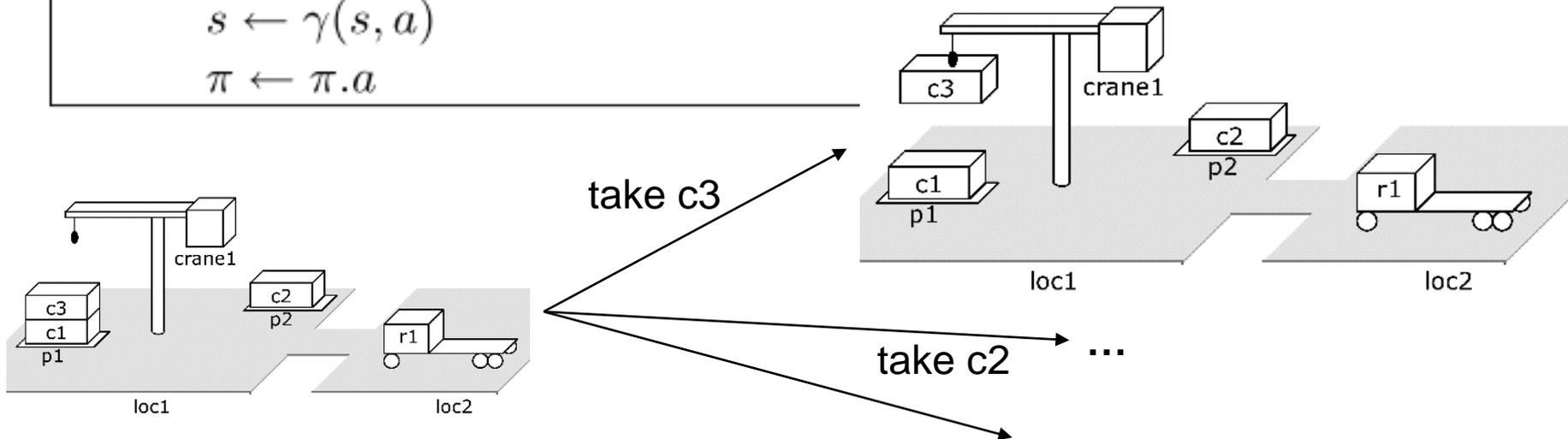
$E \leftarrow \{a \mid a \text{ is a ground instance an operator in } O,$
and $\text{precond}(a)$ is true in $s\}$

if $E = \emptyset$ then return failure

nondeterministically choose an action $a \in E$

$s \leftarrow \gamma(s, a)$

$\pi \leftarrow \pi.a$

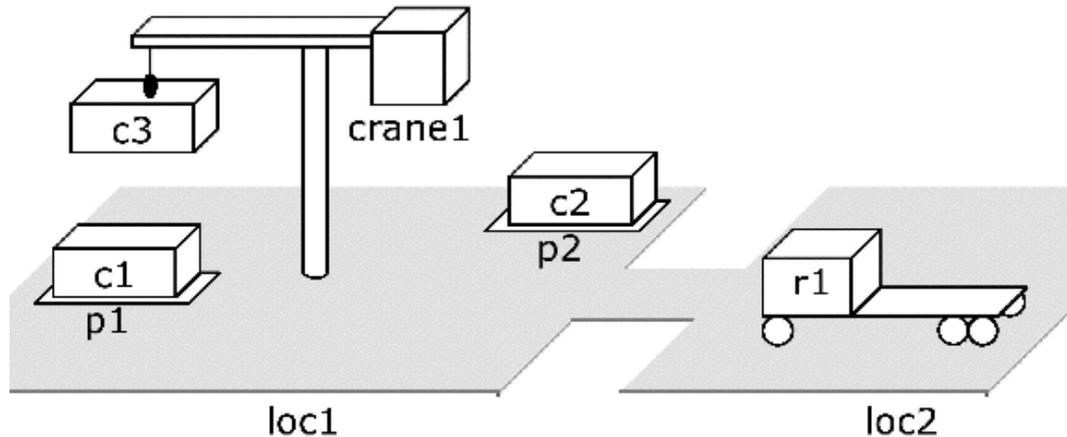


Forward Search Properties

- Forward-search *is sound*
 - ▣ for any plan returned by any of its nondeterministic traces, this plan is guaranteed to be a solution
- Forward-search *is also complete*
 - ▣ if a solution exists then at least one of Forward-search's nondeterministic traces will return a solution.

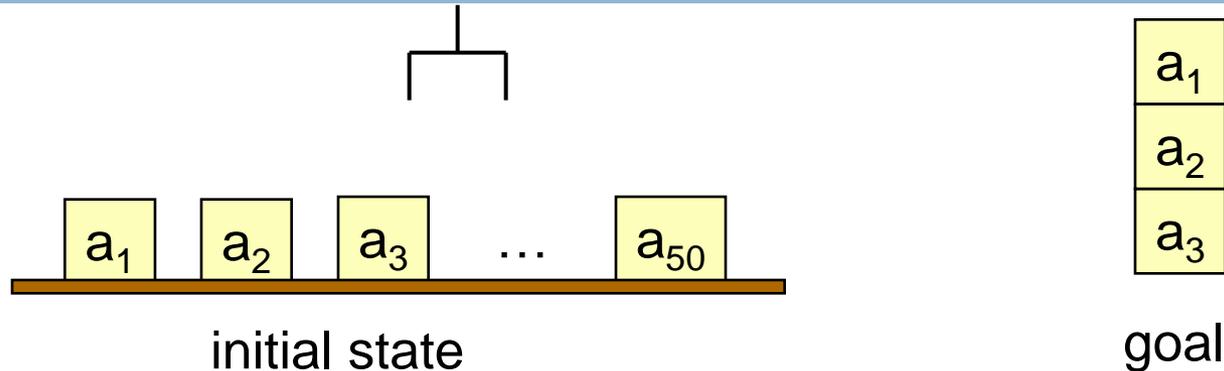
Task 1: DWR, find one finite and one infinite trace

□ S_0 :



□ g : {at(r1, loc1), loaded(r1, c3)}

Branching Factor of Forward Search



- Forward search can have a very large branching factor
 - ▣ E.g., many applicable actions that don't progress toward goal
- Why this is bad:
 - ▣ Deterministic implementations can waste time trying lots of irrelevant actions
- Need a good heuristic function and/or pruning procedure
- How to do pruning?



Backward search

Backward Search

- For forward search, we started at the initial state and computed state transitions
 - ▣ new state = $\gamma(s, a)$
- For backward search, we start at the goal and compute inverse state transitions
 - ▣ new set of subgoals = $\gamma^{-1}(g, a)$
- To define $\gamma^{-1}(g, a)$, must first define *relevance*:
 - ▣ An action a is relevant for a goal g if
 - a makes at least one of g 's literals true
 - $g \cap \text{effects}(a) \neq \emptyset$
 - a does not make any of g 's literals false
 - $g^+ \cap \text{effects}^-(a) = \emptyset$ and $g^- \cap \text{effects}^+(a) = \emptyset$

Inverse State Transitions

- If a is relevant for g , then
 - $\gamma^{-1}(g,a) = (g - \text{effects}(a)) \cup \text{precond}(a)$
- Otherwise $\gamma^{-1}(g,a)$ is undefined
- Example: suppose that
 - $g = \{\text{on}(b1,b2), \text{on}(b2,b3)\}$
 - $a = \text{stack}(b1,b2)$
- What is $\gamma^{-1}(g,a)$?

Backward Search

Backward-search(O, s_0, g)

$\pi \leftarrow$ the empty plan

loop

if s_0 satisfies g then return π

$A \leftarrow \{a \mid a \text{ is a ground instance of an operator in } O$
and $\gamma^{-1}(g, a)$ is defined}

if $A = \emptyset$ then return failure

nondeterministically choose an action $a \in A$

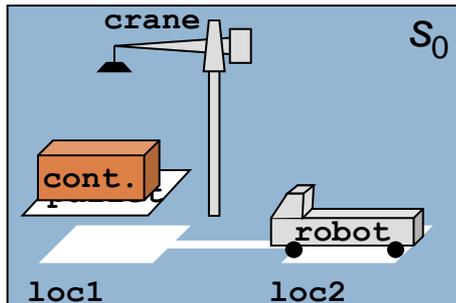
$\pi \leftarrow a.\pi$

$g \leftarrow \gamma^{-1}(g, a)$

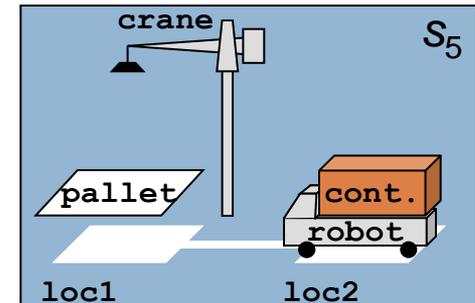
Task 2: DWR, backward search

- Solve the problem by the backward-search, trace the algorithm.
- **Actions:** load(crane, loc, cont, r), take(crane, loc, cont, pallet, pile), move(r, from, to)

initial state:



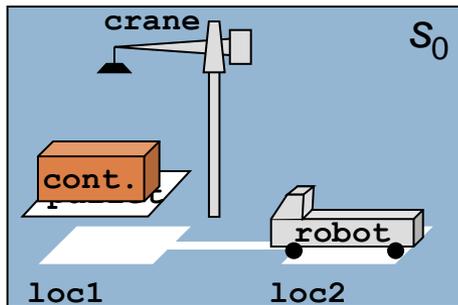
goal state:



Task 2: DWR, backward search

- Solve the problem by the backward-search, trace the algorithm.
- **Actions:** load(crane, loc, cont, r), take(crane, loc, cont, pallet, pile), move(r, from, to)

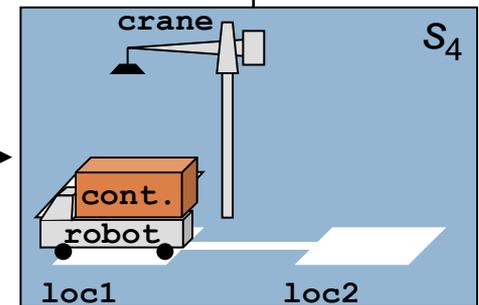
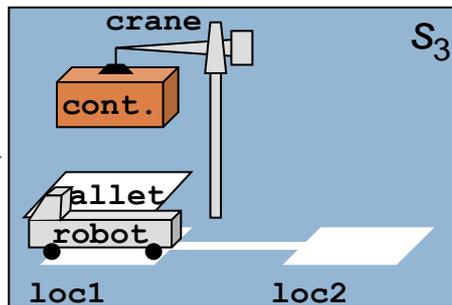
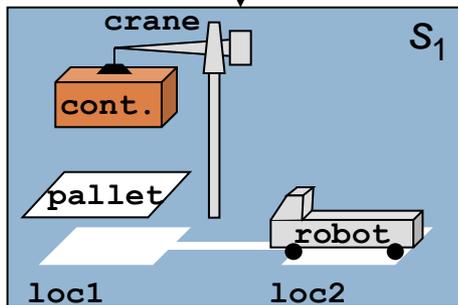
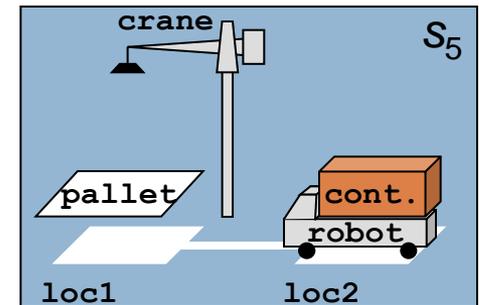
initial state:



plan =

```
take(crane,loc1,cont,pallet,pile)
move(robot,loc2,loc1)
load(crane,loc1,cont,robot)
move(robot,loc1,loc2)
```

goal state:

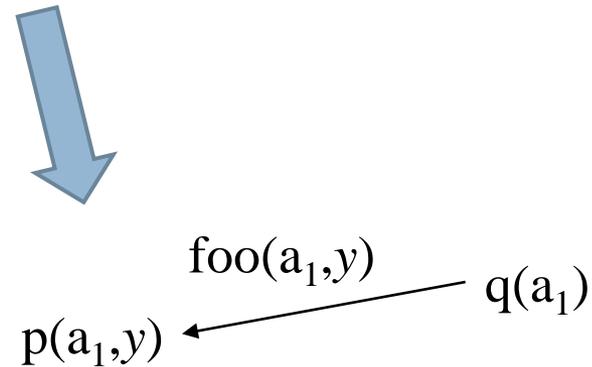
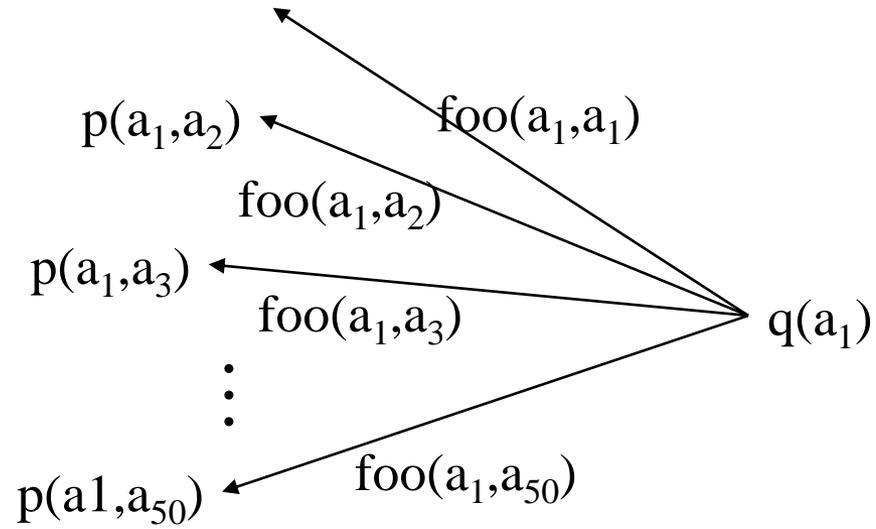


Lifting I.

- Backward search can *also* have a very large branching factor
 - E.g., an operator o that is relevant for g may have many ground instances a_1, a_2, \dots, a_n such that each a_i 's input state might be unreachable from the initial state
- Can reduce the branching factor of backward search if we *partially* instantiate the operators
 - this is called *lifting*
- Basic Idea: Delay grounding of operators until necessary in order to bind variables with those required to realize goal or subgoal

Lifting II.

$\text{foo}(x,y)$
precond: $p(x,y)$
effects: $q(x)$



Lifted Backward Search

- More complicated than Backward-search
 - Have to keep track of what substitutions were performed
- But it has a much smaller branching factor

Lifted-backward-search(O, s_0, g)

$\pi \leftarrow$ the empty plan

loop

if s_0 satisfies g then return π

$A \leftarrow \{(o, \theta) \mid o \text{ is a standardization of an operator in } O,$
 $\theta \text{ is an mgu for an atom of } g \text{ and an atom of effects}^+(o),$
 $\text{and } \gamma^{-1}(\theta(g), \theta(o)) \text{ is defined}\}$

if $A = \emptyset$ then return failure

nondeterministically choose a pair $(o, \theta) \in A$

$\pi \leftarrow$ the concatenation of $\theta(o)$ and $\theta(\pi)$

$g \leftarrow \gamma^{-1}(\theta(g), \theta(o))$



Heuristic search

Local heuristic search: Hill climbing

Hill-climbing

$\sigma := \text{make-root-node}(\text{init}())$

forever:

if $\text{is-goal}(\text{state}(\sigma))$:

return $\text{extract-solution}(\sigma)$

$\Sigma' := \{ \text{make-node}(\sigma, o, s) \mid \langle o, s \rangle \in \text{succ}(\text{state}(\sigma)) \}$

$\sigma :=$ an element of Σ' minimizing h (random tie breaking)

Enforced hill-climbing: procedure improve

```
def improve( $\sigma_0$ ):  
  queue := new fifo-queue  
  queue.push-back( $\sigma_0$ )  
  closed :=  $\emptyset$   
  while not queue.empty():  
     $\sigma$  = queue.pop-front()  
    if state( $\sigma$ )  $\notin$  closed:  
      closed := closed  $\cup$  {state( $\sigma$ )}  
      if h( $\sigma$ ) < h( $\sigma_0$ ):  
        return  $\sigma$   
      for each  $\langle o, s \rangle \in$  succ(state( $\sigma$ )):  
         $\sigma'$  := make-node( $\sigma, o, s$ )  
        queue.push-back( $\sigma'$ )  
fail
```

Enforced hill-climbing

```
 $\sigma$  := make-root-node(init())  
while not is-goal(state( $\sigma$ )):  
   $\sigma$  := improve( $\sigma$ )  
return extract-solution( $\sigma$ )
```

Systematic heuristic search: Greedy best-first search

Greedy best-first search (with duplicate detection)

```
open := new min-heap ordered by ( $\sigma \mapsto h(\sigma)$ )
open.insert(make-root-node(init()))
closed :=  $\emptyset$ 
while not open.empty():
     $\sigma = \textit{open}$ .pop-min()
    if  $\textit{state}(\sigma) \notin \textit{closed}$ :
         $\textit{closed} := \textit{closed} \cup \{\textit{state}(\sigma)\}$ 
        if is-goal( $\textit{state}(\sigma)$ ):
            return extract-solution( $\sigma$ )
        for each  $\langle o, s \rangle \in \textit{succ}(\textit{state}(\sigma))$ :
             $\sigma' := \textit{make-node}(\sigma, o, s)$ 
            if  $h(\sigma') < \infty$ :
                 $\textit{open}$ .insert( $\sigma'$ )
return unsolvable
```



**OPPA European Social Fund
Prague & EU: We invest in your future.**
