## **Inference in Description Logics**

Petr Křemen

## **1** Inference Procedures

- 1. Why inconsistency of an ontology is a problem ? What is its consequence ?
- 2. Show that disjointness of two concepts can be reduced to unsatisfiability of a single concept.
- 3. A satisfiable concept is interpreted as a non-empty set in at least one model of the theory/ontology. How to check that a concept is interpreted as a non-empty set in all models ?

## **2** Tableaux Algorithm for $\mathcal{ALC}$

- 1. Decide, whether the  $\mathcal{ALC}$  concept  $\exists hasChild \cdot (Student \sqcap Employee) \sqcap \neg (\exists hasChild \cdot Student \sqcap \exists hasChild \cdot Employee)$  is satisfiable (w.r.t. an empty TBox). Show the run of the tableau algorithm in detail.
- 2. Decide, whether the theory/ontology  $\mathcal{K} = (\mathcal{T}, \mathcal{A})$  is consistent. Show the run of the tableau algorithm in detail.
  - $\mathcal{T} = \{ \exists hasChild \cdot \top \equiv Parent \}$
  - $\mathcal{A} = \{ hasChild(JOHN, MARY), Woman(MARY) \}$
- 3. Decide and show, whether the ontology

$$\mathcal{K}_1 = \{ \mathcal{T} \cup \{ Parent \sqsubseteq \forall hasChild \cdot \neg Woman \}, \mathcal{A} \}$$

is consistent.

4. Decide and show, whether the ontology

$$\mathcal{K}_1 = \{\mathcal{T} \cup \{Parent \sqsubseteq \exists hasChild \cdot Parent\}, \mathcal{A}\}$$

is consistent.

## 3 Practically in Protégé

- 1. Model the previous ontology in Protégé and check (using the Pellet/HermiT reasoner) whether your solutions in the previous tasks were correct.
- 2. Adjust the Pizza ontology introduced in the previous seminar, so that the class *IceCream* and *CheeseyVegetableTopping* become satisfiable.
- 3. Explain, why the Pizza ontology is consistent, although it contains unsatisfiable classes.