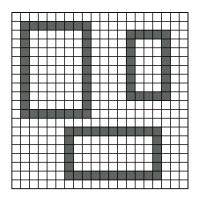
## GRAPHICAL MARKOV MODELS (WS2011) 6. SEMINAR

Assignment 1. Consider the language L of all (rectangular) b/w images containing an arbitrary number of non-overlapping rectangular frames (one pixel wide).

a) Prove that L is not expressible by a locally conjunctive predicate i.e. a conjunction of predicates, each one defined on an image fragment of a fixed size.

**b**) Show that L can be expressed by introducing non-terminal symbols and a locally conjunctive predicate for these.



Assignment 2. Transform the *Travelling Salesman Problem* into a  $(\min, +)$ -problem.

Assignment 3. Consider a CSP for K-valued labellings s of the vertices of a graph  $\mathcal{G} = (V, E)$ 

$$G(s) = \left[\bigwedge_{i \in V} g_i(s_i)\right] \land \left[\bigwedge_{ij \in E} g_{ij}(s_i, s_j)\right],$$

where  $g_{ij}$  and  $g_i$  are predicates of arity 2 and 1 respectively. The task is to calculate  $c = \bigvee_{s \in K^V} G(s)$ .

Mister X proposes the following algorithm. Edges and vertices are repeatedly visited, each time updating the functions  $g_{ij}$  and  $g_i$  respectively by

$$g_{ij}(s_i, s_j) := g_i(s_i) \land g_{ij}(s_i, s_j) \land g_j(s_j)$$
$$g_i(s_i) := g_i(s_i) \land \left[\bigwedge_{j \in N_i} \left[\bigvee_{k \in K} g_{ij}(s_i, k)\right]\right].$$

The algorithm stops in a fix-point  $g_i^*$ ,  $g_{ij}^*$ . If all these functions are identically equal to zero then c = 0 is assumed. Otherwise c is assumed to be equal to 1.

a) Prove that the algorithm is not correct in general. Construct a counterexample.

**b**<sup>\*\*</sup>) Prove that the algorithm is indeed correct if the arity 2 predicates  $g_{ij}$  are supermodular w.r.t. an ordering of the label set K, i.e.

$$g\left(\max(k_1, k_1'), \max(k_2, k_2')\right) \land g\left(\min(k_1, k_1'), \min(k_2, k_2')\right) \ge g(k_1, k_2) \land g(k_1', k_2'),$$

where the operations min and max return the greater and lower label respectively (w.r.t. the ordering of K).

Hints

- Prove, that the update rules preserve supermodularity. - Consider  $k_i^* = \max_k \{k \in K \mid g_i^*(k) = 1\}$ . Prove that  $s^*$  defined by  $s_i^* = k_i^*$  is a solution, i.e.  $G(s^*) = 1$ .