Inference in Description Logic $\mathcal{ALC}$

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1 Inference Procedures

Ex. 1 — Why inconsistency of an ontology is a problem ? What is its consequence ?

Ex. 2 — Show that disjointness of two concepts can be reduced to unsatisfiability of a single concept.

Ex. 3 — A concept C is satisfiable w.r.t. $\mathcal{K}$ iff it is interpreted as a non-empty set in at least one model of $\mathcal{K}$. Is it possible to find out that C is interpreted as a non-empty set in all models of $\mathcal{K}$?

2 Tableaux Algorithm for $\mathcal{ALC}$

Ex. 4 — Decide, whether the $\mathcal{ALC}$ concept $\exists hasChild \cdot (Student \sqcap Employee) \sqcap \neg(\exists hasChild \cdot Student \sqcap \exists hasChild \cdot Employee)$ is satisfiable (w.r.t. an empty TBox). Show the run of the tableau algorithm in detail.

Ex. 5 — Decide, whether the theory/ontology $\mathcal{K} = (\mathcal{T}, \mathcal{A})$ is consistent. Show the run of the tableau algorithm in detail.

- $\mathcal{T} = \{\exists hasChild \cdot \top \equiv Parent\}$
- $\mathcal{A} = \{hasChild(JOHN, MARY), Woman(MARY)\}$

Ex. 6 — Decide and show, whether the ontology

$$\mathcal{K}_1 = (\mathcal{T} \cup \{Parent \sqsubseteq \forall hasChild \cdot \neg Woman\}, \mathcal{A})$$

is consistent.

Ex. 7 — Decide and show, whether the ontology

$$\mathcal{K}_2 = (\mathcal{T} \cup \{Parent \sqsubseteq \exists hasChild \cdot Parent\}, \mathcal{A})$$

is consistent.

3 Practically in Protégé

Ex. 8 — Model the previous ontology in Protégé and check (using the Pellet/HermiT reasoner) whether your solutions in the previous tasks were correct.

Ex. 9 — Adjust the Pizza ontology introduced in the previous seminar, so that the class $IceCream$ and $CheeseyVegetableTopping$ become satisfiable.

Ex. 10 — Explain, why the Pizza ontology is consistent, although it contains unsatisfiable classes.