

Logical reasoning and programming, lab session I

(October 1, 2018)

I.1 Decide which of the following formulae are tautologies:

- (a) $((p \rightarrow q) \rightarrow q) \rightarrow q$,
- (b) $((p \rightarrow q) \rightarrow p) \rightarrow p$,
- (c) $(p \rightarrow q) \vee (q \rightarrow p)$,
- (d) $((p \rightarrow q) \wedge q) \rightarrow p$,
- (e) $\neg p \rightarrow \neg(p \vee (p \wedge q))$.

Although formulae contain only two variables, try to find a better method than truth tables if possible.

I.2 Assume that $\varphi \leftrightarrow \psi$ is false, meaning $v(\varphi \leftrightarrow \psi) = 0$ for all valuations v . What can we say about the validity of the following formulae:

- (a) $\varphi \wedge \psi$,
- (b) $\varphi \vee \psi$,
- (c) $\varphi \rightarrow \psi$.

I.3 If $\varphi \rightarrow \psi \in \text{TAUT}$ and $\psi \rightarrow \chi \in \text{TAUT}$, then $\varphi \rightarrow \chi \in \text{TAUT}$. Why? Does the claim still hold if we replace **TAUT** with **SAT** and why?

I.4 Let $\text{Cl}(\Gamma) = \{\varphi : \Gamma \models \varphi\}$. Decide whether to any two sets of formulae Γ and Δ hold:

- (a) $\Gamma \subseteq \text{Cl}(\Gamma)$,
- (b) $\text{Cl}(\text{Cl}(\Gamma)) = \text{Cl}(\Gamma)$,
- (c) $\text{Cl}(\Gamma \cup \Delta) = \text{Cl}(\Gamma) \cup \text{Cl}(\Delta)$.

If the equality does not hold in (b) or (c), does at least one of the inclusions hold?

I.5 Recall that a Boolean function of n -variables is a function $f: \{0, 1\}^n \rightarrow \{0, 1\}$. Describe all functions of one variable. How many distinct Boolean functions of n variables exist? Try $n = 2$ first. Do you know, why are functions **NAND** (\uparrow) and **NOR** (\downarrow) interesting?¹

I.6 Produce a formula in CNF which is equivalent to $\varphi = (a \rightarrow (c \wedge d)) \vee (b \rightarrow (c \wedge e))$. Then use the Tseytin transformation to produce a formula in CNF which is equisatisfiable to φ .

¹These connectives are also called Sheffer stroke and Peirce arrow, respectively. They are defined as $x \uparrow y := \neg(x \wedge y)$ and $x \downarrow y := \neg(x \vee y)$.