

## Logical reasoning and programming, lab session XII

(December 17, 2018)

**XII.1** Prove the following formulae using tableaux

- (a)  $\exists X \forall Y r(X, Y) \rightarrow \forall Y \exists X r(X, Y)$ ,
- (b)  $\exists X (p(X) \rightarrow \forall X p(X))$ ,
- (c)  $\forall X \forall Y (p(X) \wedge p(Y)) \rightarrow \exists X \exists Y (p(X) \vee p(Y))$ ,
- (d)  $\forall X \forall Y (p(X) \wedge p(Y)) \rightarrow \forall X \forall Y (p(X) \vee p(Y))$ ,
- (e)  $\forall X \exists Y \forall Z \exists W (r(X, Y) \vee \neg r(W, Z))$ ,
- (f)  $\exists X (\forall Y \forall Z p(Y, f(X, Y, Z)) \rightarrow (\forall Y p(Y, f(X, Y, X)) \wedge \forall Y \exists Z p(g(Y), Z))$ .

**XII.2** Try all the examples in the SMT-LIB Examples. You can use an online version of Z3, or you can install Z3 or CVC4 yourself. Another option is to use pySMT, a convenient way how to experiment with various SMT solvers in Python. If you want to learn a bit more about the Z3 prover, you should start with the tutorial.

**XII.3** If we want to combine theories in SMT using the Nelson–Oppen method, we require that both of them are stably infinite. Assume two theories  $\mathcal{T}_1$  with the language  $\{f\}$  and  $\mathcal{T}_2$  with the language  $\{g\}$ , where  $f$  and  $g$  are uninterpreted unary function symbols. Moreover,  $\mathcal{T}_1$  has only models of size at most 2 (for example, it contains  $\forall X \forall Y \forall Z (X = Y \vee X = Z)$  as an axiom). Show that the Nelson–Oppen method says that

$$f(x_1) \neq f(x_2) \wedge g(x_2) \neq g(x_3) \wedge g(x_1) \neq g(x_3).$$

is satisfiable in the union of  $\mathcal{T}_1$  and  $\mathcal{T}_2$ , but this is clearly incorrect.