### Randomized Sampling-based Motion Planning Methods

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Lecture 05

B4M36UIR – Artificial Intelligence in Robotics



### Overview of the Lecture

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### Part I

## [Part 1 – Sampling-based Motion Planning](#page-2-0)



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### (Randomized) Sampling-based Motion Planning

- It uses an explicit representation of the obstacles in  $C$ -space.
	- A "black-box" function is used to evaluate if a configuration  $q$  is a collision-free, e.g.,
	- **Based on geometrical models and testing** collisions of the models.
	- 2D or 3D shapes of the robot and environment can be represented as sets of triangles, i.e., tesselated models.
	- Collision test is then a test of for the intersection of the triangles.



E.g., using RAPID library <http://gamma.cs.unc.edu/OBB/>

- **Creates a discrete representation of**  $C_{free}$ .
- **Configurations in**  $C_{free}$  **are sampled randomly and connected to a** roadmap (probabilistic roadmap).
- Rather than the full completeness they provide **probabilistic com**pleteness or resolution completeness.

Probabilistic complete algorithms: with increasing number of samples an admissible solution would be found (if exists).



#### Probabilistic Roadmaps

A discrete representation of the continuous C-space generated by randomly sampled configurations in  $C_{free}$  that are connected into a graph.

- Nodes of the graph represent admissible configurations of the robot.
- Edges represent a feasible path (trajectory) between the particular configurations.

Probabilistic complete algorithms: with increasing number of samples an admissible solution would be found (if exists).



Having the graph, the final path (trajectory) can be found by a graph search technique.



### Incremental Sampling and Searching

- Single query sampling-based algorithms incrementally create a search graph (roadmap).
	- 1. Initialization  $G(V, E)$  an undirected search graph, V may contain  $q_{start}$ ,  $q_{goal}$  and/or other points in  $C_{free}$ .
	- 2. Vertex selection method choose a vertex  $q_{cur} \in V$  for the expansion.
	- 3. Local planning method for some  $q_{new} \in \mathcal{C}_{free}$ , attempt to construct a path  $\tau : [0,1] \to \mathcal{C}_{free}$  such that  $\tau(0) = q_{cur}$  and  $\tau(1) =$  $q_{new}$ ,  $\tau$  must be checked to ensure it is collision free.

<span id="page-6-0"></span>If  $\tau$  is not a collision-free, go to Step [2.](#page-6-0)

- 4. Insert an edge in the graph Insert  $\tau$  into E as an edge from  $q_{\text{cur}}$  to  $q_{\text{new}}$  and insert  $q_{\text{new}}$  to V if  $q_{\text{new}} \notin V$ . How to test  $q_{\text{new}}$  is in V?
- 5. Check for a solution Determine if G encodes a solution, e.g., using a single search tree or graph search technique.
- 6. Repeat Step [2](#page-6-0) iterate unless a solution has been found or a termination condition is satisfied.

LaValle, S. M.: Planning Algorithms (2006), Chapter 5.4



#### Probabilistic Roadmap Strategies

Multi-Query strategy is roadmap based.

- Generate a single roadmap that is then used for repeated planning queries.
- An representative technique is Probabilistic RoadMap (PRM).

Kavraki, L., Svestka, P., Latombe, J.-C., Overmars, M. H.B (1996): Probabilistic Roadmaps for Path Planning in High Dimensional Configuration Spaces. T-RO.

Single-Query strategy is an incremental approach.

- **For each planning problem, it constructs a new roadmap to char**acterize the subspace of C-space that is relevant to the problem.
	- **Rapidly-exploring Random Tree RRT**; LaValle, 1998
	- Expansive-Space Tree  $EST$ ;  $H_{\text{su et al.}}$  1997
	- Sampling-based Roadmap of Trees SRT.

A combination of multiple–query and single–query approaches. Plaku et al., 2005

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### Multi-Query Strategy

Build a roadmap (graph) representing the environment.

1. Learning phase

- 1.1 Sample *n* points in  $C_{free}$ .
- 1.2 Connect the random configurations using a local planner.
- 2. Query phase
	- 2.1 Connect start and goal configurations with the PRM.

E.g., using a local planner.

2.2 Use the graph search to find the path.

Probabilistic Roadmaps for Path Planning in High Dimensional Configuration Spaces Lydia E. Kavraki and Petr Svestka and Jean-Claude Latombe and Mark H. Overmars, IEEE Transactions on Robotics and Automation, 12(4):566–580, 1996.

> First planner that demonstrates ability to solve general planning problems in more than 4-5 dimensions.



#### Given problem domain:





#### Random configuration





Connecting random samples:





#### Connected roadmap:





#### Query configurations:





Final found path:





#### Practical PRM

- Incremental construction.
- **Connect nodes in a radius**  $\rho$ .
- Local planner tests collisions up to selected resolution  $\delta$ .
- Path can be found by Dijkstra's algorithm.



We need a couple of more formalisms.



#### Practical PRM

- **Incremental construction.**
- **Connect nodes in a radius**  $\rho$ .
- Local planner tests collisions up to selected resolution  $\delta$ .
- Path can be found by Dijkstra's algorithm.



#### What are the properties of the PRM algorithm?

We need a couple of more formalisms.



#### Practical PRM

- **Incremental construction.**
- **Connect nodes in a radius**  $\rho$ .
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#### What are the properties of the PRM algorithm?

We need a couple of more formalisms.



#### Path Planning Problem Formulation

**Path planning problem is defined by a triplet** 

 $P = (C_{free}, q_{init}, Q_{goal})$ , where

- $\mathcal{C}_{\text{free}} = \text{cl}(\mathcal{C} \setminus \mathcal{C}_{\text{obs}}), \; \mathcal{C} = (0, 1)^d, \; \text{for} \; d \in \mathbb{N}, \; d \geq 2;$  (scaling)
- q<sub>init</sub>  $\in \mathcal{C}_{free}$  is the initial configuration (condition);
- $\Box$   $\mathcal{Q}_{goal}$  is the goal region defined as an open subspace of  $\mathcal{C}_{free}$ .

#### **Function**  $\pi : [0,1] \to \mathbb{R}^d$  of *bounded variation* is called:

- $\blacksquare$  path if it is continuous;
- collision-free path if it is a path and  $\pi(\tau) \in \mathcal{C}_{\text{free}}$  for  $\tau \in [0,1]$ ;

**feasible** if it is a collision-free path, and  $\pi(0) = q_{init}$  and  $\pi(1) \in$  $cl(Q_{goal}).$ 

A function  $\pi$  with the total variation  $TV(\pi)<\infty$  is said to have bounded variation, where  $TV(\pi)$  is the total variation  $TVI(\underline{\hspace{0.3cm}})$  = sup

$$
\text{TV}(\pi) = \sup_{\{n \in \mathbb{N}, 0 = \tau_0 < \tau_1 < \ldots < \tau_n = s\}} \sum_{i=1}^n |\pi(\tau_i) - \pi(\tau_{i-1})|.
$$

The total variation  $TV(\pi)$  is de facto a path length.



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#### Path Planning Problem

#### Feasible path planning

For a path planning problem  $(\mathcal{C}_{free}, q_{init}, \mathcal{Q}_{goal})$ :

- Find a feasible path  $\pi : [0,1] \rightarrow \mathcal{C}_{free}$  such that  $\pi(0) = q_{init}$  and  $\pi(1) \in \text{cl}(\mathcal{Q}_{\text{goal}})$ , if such path exists;
- Report failure if no such path exists.

#### ■ Optimal path planning

For  $(\mathcal{C}_{free}, q_{init}, \mathcal{Q}_{goal})$  and a cost function  $c : \Sigma \to \mathbb{R}_{\geq 0}$ :

■ Find a feasible path  $\pi^*$  such that  $c(\pi^*) = \min\{c(\pi) : \pi$  is feasible};

Report failure if no such path exists.

The cost function is assumed to be monotonic and bounded, i.e., there exists  $k_c$  such that  $c(\pi) \leq k_c \text{TV}(\pi)$ 



#### Path Planning Problem

#### Feasible path planning

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- Find a feasible path  $\pi : [0,1] \rightarrow \mathcal{C}_{free}$  such that  $\pi(0) = q_{init}$  and  $\pi(1) \in \text{cl}(\mathcal{Q}_{\text{goal}})$ , if such path exists;
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#### ■ Optimal path planning

The optimality problem asks for a feasible path with the minimum cost.

For  $(\mathcal{C}_{free}, q_{init}, \mathcal{Q}_{goal})$  and a cost function  $c : \Sigma \to \mathbb{R}_{\geq 0}$ :

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### Probabilistic Completeness 1/2

First, we need robustly feasible path planning problem  $(\mathcal{C}_{free}, q_{init}, \mathcal{Q}_{goal})$ .

**q**  $\in$  C<sub>free</sub> is  $\delta$ -interior state of  $\mathcal{C}_{free}$  if the closed ball of radius  $\delta$  centered at q lies entirely inside  $C_{free}$ .



 $\blacksquare$  δ-interior of  $\mathcal{C}_{\text{free}}$  is  $\text{int}_{\delta}(\mathcal{C}_{\text{free}}) = \{q \in \mathcal{C}_{\text{free}} | \mathcal{B}_{\ell,\delta} \subseteq \mathcal{C}_{\text{free}} \}.$ 

A collection of all δ-interior states.

- A collision free path  $\pi$  has strong  $\delta$ -clearance, if  $\pi$  lies entirely inside int<sub> $\delta$ </sub> $(C_{free})$ .
- $(C_{free}, q_{init}, Q_{goal})$  is robustly feasible if a solution exists and it is a feasible path with *strong*  $\delta$ *-clearance*, for  $\delta > 0$ .

Probabilistic Completeness 2/2

An algorithm  $ALG$  is probabilistically complete if, for any robustly feasible path planning problem  $P = (C_{free}, q_{init}, Q_{goal})$ ,

 $\lim_{n\to\infty} Pr(\mathcal{ALG})$  returns a solution to  $\mathcal{P})=1.$ 

It is a "relaxed" notion of the completeness.

Applicable only to problems with a robust solution.



We need some space, where random configurations can be sampled.



## Asymptotic Optimality 1/4 **Homotopy**

Asymptotic optimality relies on a notion of weak  $\delta$ -clearance.

Notice, we use strong  $\delta$ -clearance for probabilistic completeness.

- $\blacksquare$  We need to describe possibly improving paths (during the planning).
- **Function**  $\psi : [0, 1] \rightarrow \mathcal{C}_{free}$  is called **homotopy**, if  $\psi(0) = \pi_1$  and  $\psi(1) = \pi_2$  and  $\psi(\tau)$  is collision-free path for all  $\tau \in [0,1]$ .
- A collision-free path  $\pi_1$  is **homotopic** to  $\pi_2$  if there exists homotopy function  $\psi$ .

A path homotopic to  $\pi$  can be continuously transformed to  $\pi$  through  $C_{free}$ .



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- A collision-free path  $\pi_1$  is homotopic to  $\pi_2$  if there exists homotopy function  $\psi$ .

A path homotopic to  $\pi$  can be continuously transformed to  $\pi$  through  $C_{free}$ .

### Asymptotic Optimality 2/4 Weak δ-clearance

A collision-free path  $\pi : [0, s] \rightarrow C_{free}$  has weak  $\delta$ -clearance if there exists a path  $\pi'$  that has strong  $\delta\text{-}$ clearance and homotopy  $\psi$  with  $\psi(0)=\pi$ ,  $\psi(1)=\pi'$ , and for all  $\alpha\in(0,1]$  there exists  $\delta_{\alpha} > 0$  such that  $\psi(\alpha)$  has strong  $\delta$ -clearance.

> Weak δ-clearance does not require points along a path to be at least a distance  $\delta$  away from obstacles.



- **C** A path  $\pi$  with a weak  $\delta$ -clearance.
- $\blacksquare$   $\pi'$  lies in int $_\delta(\mathcal{C}_{\mathit{free}})$  and it is the same homotopy class as  $\pi$ .



### Asymptotic Optimality 3/4 Robust Optimal Solution

- It is applicable with a robust optimal solution that can be obtained as a limit of robust (non-optimal) solutions.
- A collision-free path  $\pi^*$  is robustly optimal solution if it has weak δ-clearance and for any sequence of collision free paths  $\{\pi_n\}_{n\in\mathbb{N}}$ ,  $\pi_n \in \mathcal{C}_{\text{free}}$  such that  $\lim_{n \to \infty} \pi_n = \pi^*$ ,

$$
\lim_{n\to\infty}c(\pi_n)=c(\pi^*).
$$

There exists a path with strong  $\delta$ -clearance, and  $\pi^*$  is homotopic to such path and  $\pi^*$  is of the lower cost.

 $\blacksquare$  Weak  $\delta$ -clearance implies robustly feasible solution problem.

Thus, it implies the probabilistic completeness.



### Asymptotic Optimality 4/4 Asymptotically optimal algorithm

An algorithm  $\mathcal{ALG}$  is asymptotically optimal if, for any path planning problem  $P = (C_{free}, q_{init}, Q_{goal})$  and cost function c that admit a robust optimal solution with the finite cost  $c^*$ 

$$
Pr\left(\left\{\lim_{i\to\infty}Y_i^{\mathcal{ALG}}=c^*\right\}\right)=1.
$$

 $\blacksquare$   $Y_i^{\mathcal{ALG}}$  $\zeta^{ALG}_{i}$  is the extended random variable corresponding to the minimumcost solution included in the graph returned by  $ALG$  at the end of the iteration i.



### Properties of the PRM Algorithm

- Completeness for the standard PRM has not been provided when it was introduced.
- A simplified version of the PRM (called sPRM) has been mostly studied.
- **sPRM** is probabilistically complete.

What are the differences between PRM and sPRM?



### PRM vs simplified PRM (sPRM)

#### Algorithm 1: PRM

**Input**:  $q_{init}$ , number of samples *n*, radius  $\rho$ Output:  $PRM - G = (V, E)$ 

$$
V \leftarrow \emptyset; E \leftarrow \emptyset;
$$
\nfor  $i = 0, ..., n$  do

\n
$$
\begin{array}{c}\nq_{rand} \leftarrow \text{SampleFree}; \\
q_{rand} \leftarrow \text{SampleFree}; \\
U \leftarrow \text{Near}(G = (V, E), q_{rand}, \rho); \\
V \leftarrow V \cup \{q_{rand}\}; \\
\text{forecast } u \in U, \text{ with increasing} \\
|u - q_r|| \text{ do} \\
\qquad \qquad \text{if } q_{rand} \text{ and } u \text{ are not in the} \\
q_{rand} \text{ some connected component of} \\
G = (V, E) \text{ then} \\
q_{rand}, u) \\
q_{end} \text{ then} \\
E \leftarrow E \cup \\
q_{rand}, u), (u, q_{rand})\}; \\
\text{return } G = (V, E); \\
\end{array}
$$

Algorithm 2: sPRM

**Input:**  $q_{init}$ , number of samples *n*, radius ρ **Output:** PRM –  $G = (V, E)$  $V \leftarrow \{q_{init}\}$  ∪  ${SampleFree_i}_{i=1,\ldots,n-1}$ ;  $E \leftarrow \emptyset$ ; foreach  $v \in V$  do  $U \leftarrow \text{Near}(G = (V, E), v, \rho) \setminus \{v\};$ foreach  $u \in U$  do **if** CollisionFree $(v, u)$  then  $E \leftarrow E \cup \{ (v, u), (u, v) \};$ return  $G = (V, E)$ ;

There are several ways for the set  $U$  of vertices to connect them:

- $\blacksquare$  k-nearest neighbors to v;
- variable connection radius  $\rho$  as a function of n.



#### PRM – Properties

#### **SPRM** (simplified PRM):

- **Probabilistically complete and asymptotically optimal.**
- Processing complexity can be bounded by  $O(n^2)$ .
- Query complexity can be bounded by  $O(n^2)$ .
- Space complexity can be bounded by  $O(n^2)$ .
- $\blacksquare$  Heuristics practically used are usually not probabilistic complete.
	- $\blacksquare$  k-nearest sPRM is not probabilistically complete.
	- Variable radius sPRM is not probabilistically complete.

Based on analysis of Karaman and Frazzoli

#### PRM algorithm

- $+$  It has very simple implementation.
- $+$  It provides completeness (for sPRM).
- − Differential constraints (car-like vehicles) are not straightforward.



#### Comments about Random Sampling 1/2

■ Different sampling strategies (distributions) may be applied.



- Notice, one of the main issue of the randomized sampling-based approaches is the narrow passage.
- Several modifications of sampling based strategies have been proposed in the last decades.



### Comments about Random Sampling 2/2

A solution can be found using only a few samples.

Do you know the Oraculum? (from Alice in Wonderland)

- Sampling strategies are important:
	- Near obstacles:
	- Narrow passages;
	- Grid-based:
	- **Uniform sampling must be carefully considered.**

James J. Kuffner (2004): Effective Sampling and Distance Metrics for 3D Rigid Body Path Planning. ICRA.







Naïve sampling Uniform sampling of SO(3) using Euler angles

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### Rapidly Exploring Random Tree (RRT)

#### Single–Query algorithm

- If incrementally builds a graph (tree) towards the goal area. It does not guarantee precise path to the goal configuration.
- 1. Start with the initial configuration  $q_0$ , which is a root of the constructed graph (tree).
- 2. Generate a new random configuration  $q_{new}$  in  $C_{free}$ .
- 3. Find the closest node  $q_{near}$  to  $q_{new}$  in the tree.

E.g., using KD-tree implementation like ANN or FLANN libraries.

4. Extend  $q_{near}$  towards  $q_{new}$ .

Extend the tree by a small step, but often a direct control  $u \in U$  that will move robot the position closest to q<sub>new</sub> is selected (applied for  $\delta t$ ).

5. Go to Step 2, until the tree is within a sufficient distance from the goal configuration.

Or terminates after dedicated running time.



#### RRT Construction



### RRT Algorithm

- Motivation is a single query and control-based path finding
- If incrementally builds a graph (tree) towards the goal area

Algorithm 3: Rapidly Exploring Random Tree (RRT)

```
\nInput: 
$$
q_{init}
$$
, number of samples n\nOutput: Roadmap  $G = (V, E)$ \n $V \leftarrow \{q_{init}\}; E \leftarrow \emptyset;$ \n\nfor  $i = 1, \ldots, n$  do\n\n $q_{rand} \leftarrow \text{SampleFree};$ \n $q_{nearest} \leftarrow \text{Nearest}(G = (V, E), q_{rand});$ \n $q_{new} \leftarrow \text{Steve}(q_{nearest}, q_{new})$ \n\nif CollisionFree(q_{nearest}, q_{new}) then\n\n $\begin{bmatrix}\nV \leftarrow V \cup \{x_{new}\}; E \leftarrow E \cup \{(x_{nearest}, x_{new})\}; \\
\end{bmatrix}$ \n
```

return  $G = (V, E)$ ;

Extend the tree by a small step, but often a direct control  $u \in U$  that will move robot to the position closest to  $q_{new}$  is selected (applied for dt)



#### Properties of RRT Algorithms

■ The RRT algorithm rapidly explores the space.

qnew will more likely be generated in large not yet covered parts.

- Allows considering kinodynamic/dynamic constraints (during the expansion).
- Can provide trajectory or a sequence of direct control commands for robot controllers.
- A collision detection test is usually used as a "black-box".

E.g., RAPID, Bullet libraries.

- **Similarly to PRM, RRT algorithms have poor performance in narrow** passage problems.
- RRT algorithms provides feasible paths.

It can be relatively far from optimal solution, e.g., according to the length of the path.

Many variants of RRT have been proposed.



### RRT – Examples 1/2







Alpha puzzle benchmark Apply rotations to reach the goal



Bugtrap benchmark Variants of RRT algorithms

Courtesy of V. Vonásek and P. Vaněk



#### RRT – Examples 2/2



Planning for a car-like robot





#### Car-Like Robot

■ Configuration

$$
\vec{x} = \begin{pmatrix} x \\ y \\ \phi \end{pmatrix}
$$

position and orientation.

■ Controls

$$
\vec{u} = \begin{pmatrix} v \\ \varphi \end{pmatrix}
$$

forward velocity, steering angle.

System equation

$$
\dot{x} = v \cos \phi
$$
\n
$$
\dot{y} = v \sin \phi
$$
\n
$$
\dot{\phi} = \frac{v}{L} \tan \varphi
$$



Differential constraints on possible  $\dot{q}$ :

 $\dot{x}$ sin $(\phi) - \dot{y}$  cos $(\phi) = 0$ .



Jan Faigl, 2018

#### Control-Based Sampling

- Select a configuration q from the tree T of the current configurations.
- Pick a control input  $\vec{u} = (v, \varphi)$  and the integrate system (motion) equation over a short period  $\Delta t$ :

$$
\left(\begin{array}{c}\Delta x\\ \Delta y\\ \Delta \varphi\end{array}\right)=\int_{t}^{t+\Delta t}\left(\begin{array}{c}\mathsf{v}\cos\phi\\ \mathsf{v}\sin\phi\\ \frac{\mathsf{v}}{L}\tan\varphi\end{array}\right)dt.
$$



If the motion is collision-free, add the endpoint to the tree.

E.g., considering k configurations for  $k\delta t = dt$ .



## Part II

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**[Optimal Motion Planners](#page-47-0)** and the state of the [Rapidly-exploring Random Graph \(RRG\)](#page-51-0)

#### **Outline**

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### Sampling-Based Motion Planning

- PRM and RRT are theoretically probabilistic complete.
- They provide a feasible solution without quality guarantee.

Despite that, they are successfully used in many practical applications.

In 2011, a systematical study of the asymptotic behavior of randomized sampling-based planners has been published.

> It shows, that in some cases, they converge to a non-optimal value with a probability 1.

Based on the study, new algorithms have been proposed: RRG and optimal RRT (RRT<sup>∗</sup> ).

Karaman, S., Frazzoli, E. (2011):Sampling-based algorithms for optimal motion planning. IJRR.







<http://sertac.scripts.mit.edu/rrtstar>

#### RRT and Quality of Solution 1/2

- $\blacksquare$  Let  $Y_i^{RRT}$  be the cost of the best path in the RRT at the end of the iteration i.
- $\blacksquare$   $Y_i^{RRT}$  converges to a random variable

$$
\lim_{i \to \infty} Y_i^{RRT} = Y_{\infty}^{RRT}.
$$

 $\;\blacksquare\;$  The random variable  $Y^{RRT}_\infty$  is sampled from a distribution with zero mass at the optimum, and

$$
Pr[Y_{\infty}^{RRT} > c^*] = 1.
$$

Karaman and Frazzoli, 2011

■ The best path in the RRT converges to a sub-optimal solution almost surely.

#### RRT and Quality of Solution 2/2

- RRT does not satisfy a necessary condition for the asymptotic optimality.
	- For  $0 < R < \inf_{q \in \mathcal{Q}_{goal}} ||q q_{init}||$ , the event  $\{\lim_{n \to \infty} Y_n^{RTT} = c^*\}$ occurs only if the k-th branch of the RRT contains vertices outside the R-ball centered at  $q_{init}$  for infinitely many k.

See Appendix B in Karaman and Frazzoli, 2011

 It is required the root node will have infinitely many subtrees that extend at least a distance  $\epsilon$  away from  $q_{init}$ .

> The sub-optimality is caused by disallowing new better paths to be discovered.



#### **Outline**

<span id="page-51-0"></span>[Optimal Motion Planners](#page-47-0)

[Rapidly-exploring Random Graph \(RRG\)](#page-51-0)



### Rapidly-exploring Random Graph (RRG)

Algorithm 4: Rapidly-exploring Random Graph (RRG) Input:  $q_{init}$ , number of samples n Output:  $G = (V, E)$  $V \leftarrow \emptyset: E \leftarrow \emptyset$ ; for  $i = 0, \ldots, n$  do  $q_{rand} \leftarrow$  SampleFree;  $q_{nearest} \leftarrow$  Nearest $(G = (V, E), q_{rand})$ ;  $q_{new} \leftarrow \textsf{Steer}(q_{nearest}, q_{rand});$ if CollisionFree( $q_{nearest}, q_{new}$ ) then  $\mathcal{Q}_\mathit{near} \leftarrow \mathsf{Near}(\mathit{G} = \emptyset)$  $(V, E), q_{new}$ , min $\{\gamma_{RRG}(\log(\text{card}(V))/\text{card}(V))^{1/d}, \eta\});$  $V \leftarrow V \cup \{q_{new}\};$  $E \leftarrow E \cup \{(q_{nearest}, q_{new}), (q_{new}, q_{nearest})\};$ foreach  $q_{near} \in \mathcal{Q}_{near}$  do if CollisionFree( $q_{near}, q_{new}$ ) then<br>  $\mid E \leftarrow E \cup \{(q_{rand}, u), (u, q_{rand})\};$ return  $G = (V, E)$ ;

Proposed by Karaman and Frazzoli (2011). Theoretical results are related to properties of Random Geometric Graphs (RGG) introduced by Gilbert (1961) and further studied by Penrose (1999).



#### RRG Expansions

- At each iteration, RRG tries to connect new sample to all vertices in the  $r_n$  ball centered at it.
- The ball of radius

$$
r(\text{card}(V)) = \min \left\{ \gamma_{RRG} \left( \frac{\log (\text{card}(V))}{\text{card}(V)} \right)^{1/d}, \eta \right\},\
$$

where

- $\blacksquare$   $\eta$  is the constant of the local steering function;
- $\gamma_{RRG} > \gamma_{RRG}^{*} = 2(1 + 1/d)^{1/d} (\mu(\mathcal{C}_{free})/\xi_d)^{1/d};$ 
	- $d$  dimension of the space;
	- $\mu(\mathcal{C}_{free})$  Lebesgue measure of the obstacle–free space;
	- $\xi_d$  volume of the unit ball in d-dimensional Euclidean space.
- $\blacksquare$  The connection radius decreases with n.
- The rate of decay  $\approx$  the average number of connections attempted is proportional to  $log(n)$ .



### RRG Properties

- **Probabilistically complete;**
- Asymptotically optimal;
- Complexity is  $O(\log n)$ .

(per one sample)

- Computational efficiency and optimality:
	- If attempts a connection to  $\Theta(\log n)$  nodes at each iteration;

in average

- Reduce volume of the "connection" ball as  $log(n)/n$ ;
- Increase the number of connections as  $log(n)$ .



#### Other Variants of the Optimal Motion Planning

**PRM\*** follows the standard PRM algorithm where connections are attempted between roadmap vertices that are the within connection radius  $r$  as the function of  $n$ 

$$
r(n) = \gamma_{PRM}(\log(n)/n)^{1/d}.
$$

 $\blacksquare$  RRT\* is a modification of the RRG, where cycles are avoided.

It is a tree version of the RRG.

- A tree roadmap allows to consider non-holonomic dynamics and kinodynamic constraints.
- It is basically the RRG with "rerouting" the tree when a better path is discovered.

#### Example of Solution 1/3



Karaman & Frazzoli, 2011



#### Example of Solution 2/3





#### Example of Solution 3/3



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### Overview of Randomized Sampling-based Algorithms



Notice, k-nearest variants of RRG, PRM\*, and RRT\* are complete and optimal as well.



## Summary of the Lecture



### Summary

<span id="page-61-0"></span>Properties of the sampling-based motion planning algorithms

- Single and multi-query approaches
- Path, collision-free path, feasible path
- **Feasible path planning and optimal path planning**
- **Probabilistic completeness, strong**  $\delta$ **-clearance, robustly fea**sible path planning problem
- **Asymptotic optimality, homotopy, weak**  $\delta$ **-clearance, robust** optimal solution
- PRM, RRT, RRG, PRM\*, RRT\*



#### [Topics Discussed](#page-61-0)

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- Randomized Sampling-based Methods
- Probabilistic Road Map (PRM)
- Characteristics of path planning problems
- Random sampling
- Rapidly Exploring Random Tree (RRT)
- Optimal sampling-based motion planning
- Rapidly-exploring Random Graph (RRG)
- Next: Improved Sampling-based Motion Planning
- Next: Multi-Goal Motion Planning and Multi-Goal Path Planning



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