

Path Planning

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Lecture 03

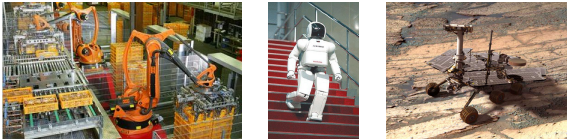
B4M36UIR – Artificial Intelligence in Robotics

Robot Motion Planning – Motivational problem

- How to transform high-level task specification (provided by humans) into a low-level description suitable for controlling the actuators?

To develop algorithms for such a transformation.

The motion planning algorithms provide transformations how to move a robot (object) considering all operational constraints.



It encompasses several disciplines, e.g., mathematics.

Real Mobile Robots

In a real deployment, the problem is a more complex.

- The world is changing
- Robots update the knowledge about the environment
- New decisions have to be made
- A feedback from the environment



Josef Štrunc, Bachelor thesis, CTU, 2009.

An example of robotic mission:

Multi-robot exploration of unknown environment

How to deal with real-world complexity?

Relaxing constraints and considering realistic assumptions.

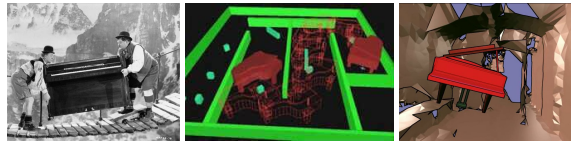
Overview of the Lecture

- Part 1 – Path Planning
 - Introduction to Path Planning
 - Notation and Terminology
 - Path Planning Methods

Piano Mover's Problem

A classical motion planning problem

Having a CAD model of the piano, model of the environment, the problem is how to move the piano from one place to another without hitting anything.



Basic motion planning algorithms are focused primarily on rotations and translations.

- We need notion of model representations and formal definition of the problem.
- Moreover, we also need a context about the problem and realistic assumptions.

The plans have to be admissible and feasible.

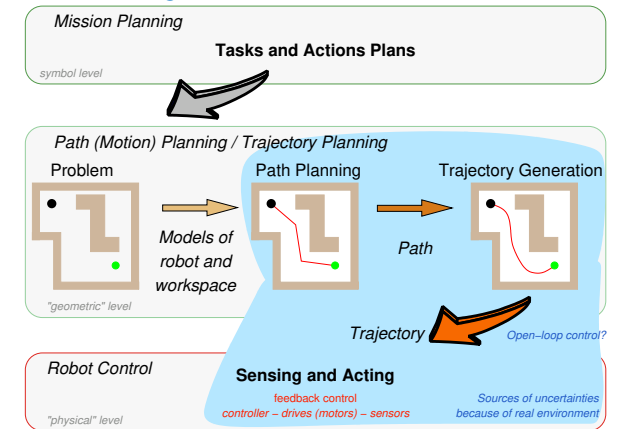
Notation

- \mathcal{W} – **World model** describes the robot workspace and its boundary determines the obstacles \mathcal{O}_i .
 $2D\ world, \mathcal{W} = \mathbb{R}^2$
- A **Robot** is defined by its geometry, parameters (kinematics) and it is controllable by the motion plan.
- \mathcal{C} – **Configuration space (C-space)**
A concept to describe possible configurations of the robot. The robot's configuration completely specify the robot location in \mathcal{W} including specification of all degrees of freedom.
E.g., a robot with rigid body in a plane $\mathcal{C} = \{x, y, \varphi\} = \mathbb{R}^2 \times S^1$.
 - Let \mathcal{A} be a subset of \mathcal{W} occupied by the robot, $\mathcal{A} = \mathcal{A}(q)$.
 - A subset of \mathcal{C} occupied by obstacles is
 $\mathcal{C}_{obs} = \{q \in \mathcal{C} : \mathcal{A}(q) \cap \mathcal{O}_i, \forall i\}$
 - Collision-free configurations** are
 $\mathcal{C}_{free} = \mathcal{C} \setminus \mathcal{C}_{obs}$.

Part I

Part 1 – Path and Motion Planning

Robotic Planning Context



Path / Motion Planning Problem

- Path** is a continuous mapping in \mathcal{C} -space such that
 $\pi : [0, 1] \rightarrow \mathcal{C}_{free}$, with $\pi(0) = q_0$, and $\pi(1) = q_f$,
- Trajectory** is a path with explicit parametrization of time, e.g., accompanied by a description of the motion laws ($\gamma : [0, 1] \rightarrow \mathcal{U}$, where \mathcal{U} is robot's action space).
It includes dynamics.
 $[T_0, T_f] \ni t \rightsquigarrow \tau \in [0, 1] : q(t) = \pi(\tau) \in \mathcal{C}_{free}$

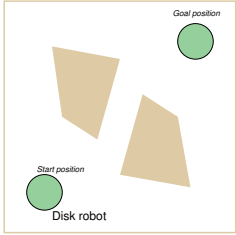
The path planning is the determination of the function $\pi(\cdot)$.

Additional requirements can be given:

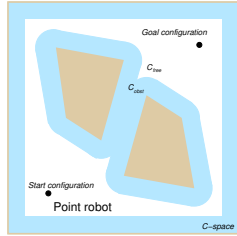
- Smoothness of the path
 - Kinodynamic constraints – e.g., considering friction forces
 - Optimality criterion – shortest vs fastest (length vs curvature)
- Path planning** – planning a collision-free path in \mathcal{C} -space
 - Motion planning** – planning collision-free motion in the **state space**

Planning in C -space

Robot motion planning robot for a disk robot with a radius ρ .



Motion planning problem in geometrical representation of \mathcal{W}

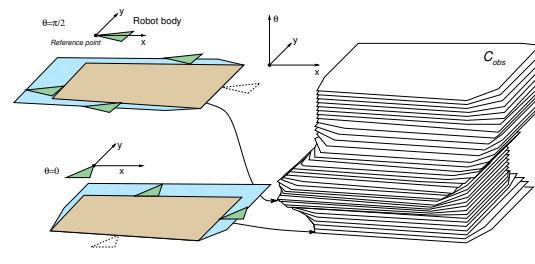


Motion planning problem in C -space representation

C -space has been obtained by enlarging obstacles by the disk \mathcal{A} with the radius ρ .

By applying Minkowski sum: $\mathcal{O} \oplus \mathcal{A} = \{x + y \mid x \in \mathcal{O}, y \in \mathcal{A}\}$.

Example of C_{obs} for a Robot with Rotation



A simple 2D obstacle \rightarrow has a complicated C_{obs}

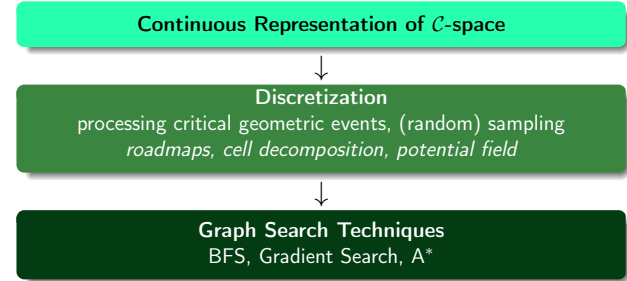
- Deterministic algorithms exist

Requires exponential time in C dimension,
J. Canny, PAMI, 8(2):200–209, 1986

- Explicit representation of C_{free} is impractical to compute.

Representation of C -space

How to deal with continuous representation of C -space?



Planning Methods - Overview

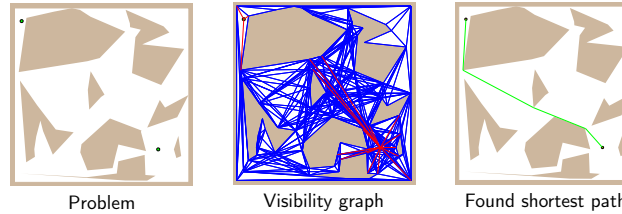
(selected approaches)

- **Point-to-point** path/motion planning
Multi-goal path/motion/trajectory planning later
- **Roadmap based methods** – Create a connectivity graph of the free space.
 - Visibility graph (complete but impractical)
 - Cell decomposition
 - Voronoi graph
- Discretization into a **grid-based** (or lattice-based) representation (resolution complete)
- **Potential field methods** (complete only for a “navigation function”, which is hard to compute in general)
Classic path planning algorithms
- **Randomized sampling-based methods**
 - Creates a roadmap from connected random samples in C_{free}
 - Probabilistic roadmaps
samples are drawn from some distribution
 - Very successful in practice

Visibility Graph

1. Compute visibility graph
2. Find the shortest path

E.g., by Dijkstra's algorithm

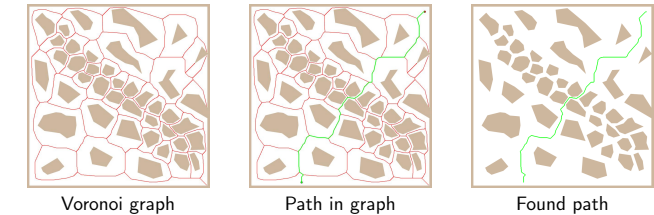


Constructions of the visibility graph:

- Naïve – all segments between n vertices of the map $O(n^3)$
- Using rotation trees for a set of segments – $O(n^2)$
M. H. Overmars and E. Welzl, 1988

Voronoi Graph

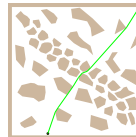
1. Roadmap is Voronoi graph that maximizes clearance from the obstacles
2. Start and goal positions are connected to the graph
3. Path is found using a graph search algorithm



Visibility Graph vs Voronoi Graph

Visibility graph

- Shortest path, but it is close to obstacles. We have to consider safety of the path
An error in plan execution can lead to a collision.
- Complicated in higher dimensions



Voronoi graph

- It maximizes clearance, which can provide conservative paths
- Small changes in obstacles can lead to large changes in the graph
- Complicated in higher dimensions



A combination is called Visibility-Voronoi – R. Wein, J. P. van den Berg, D. Halperin, 2004

For higher dimensions we need other roadmaps.

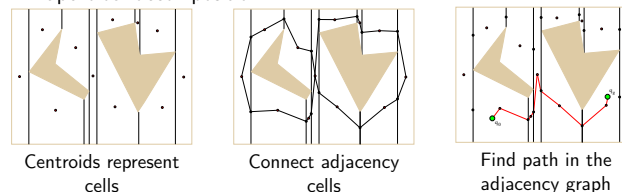
Cell Decomposition

1. Decompose free space into parts.

Any two points in a convex region can be directly connected by a segment.

2. Create an adjacency graph representing the connectivity of the free space.
3. Find a path in the graph.

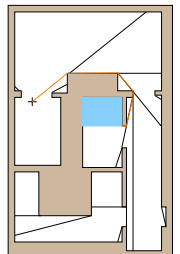
Trapezoidal decomposition



Other decomposition (e.g., triangulation) are possible.

Shortest Path Map (SPM)

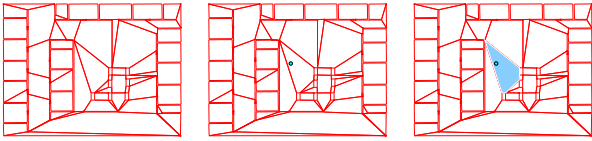
- Speedup computation of the shortest path towards a particular goal location p_g for a polygonal domain \mathcal{P} with n vertices
- A partitioning of the free space into cells with respect to the particular location p_g
- Each cell has a vertex on the shortest path to p_g
- Shortest path from any point p is found by determining the cell (in $O(\log n)$ using point location alg.) and then traversing the shortest path with up to k bends, i.e., it is found in $O(\log n + k)$
- Determining the SPM using “wavefront” propagation based on *continuous Dijkstra paradigm*
- SPM is a precompute structure for the given \mathcal{P} and p_g
 - single-point query



Joseph S. B. Mitchell: A new algorithm for shortest paths among obstacles in the plane, *Annals of Mathematics and Artificial Intelligence*, 3(1):83–105, 1991.

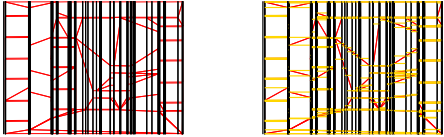
Point Location Problem

- For a given partitioning of the polygonal domain into a discrete set of cells, determine the cell for a given point p



Masato Edahiro, Iwao Kokubo and Takao Asano: *A new point-location algorithm and its practical efficiency: comparison with existing algorithms*, *ACM Trans. Graph.*, 3(2):86–109, 1984.

- It can be implemented using **interval trees** – slabs and slices

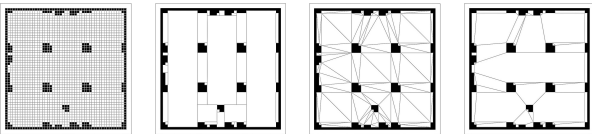


Point location problem, SPM and similarly problems are from the [Computational Geometry field](#)

Navigation Mesh

- In addition to robotic approaches, fast shortest path queries are studied in computer games
- There is a class of algorithms based on navigation mesh
 - A supporting structure representing the free space

It usually originated from the grid based maps, but it is represented as **CDT** – **Constrained Delaunay triangulation**



Grid mesh Merged grid mesh CDT mesh Merged CDT mesh

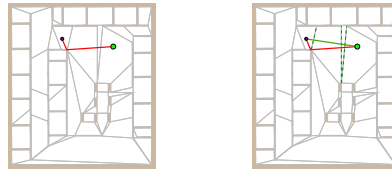
- E.g., **Polyanya** algorithm based on navigation mesh and best-first search
M. Cui, D. Harabor, A. Grastien: *Compromise-free Pathfinding on a Navigation Mesh*, *IJCAI 2017*, 496–502.
<https://bitbucket.org/dharabor/pathfinding>

Informative

Summary of the Lecture

Approximate Shortest Path and Navigation Mesh

- We can use any convex partitioning of the polygonal map to speed up shortest path queries
 - Precompute all shortest paths from map vertices to p_g using visibility graph
 - Then, an estimation of the shortest path from p to p_g is the shortest path among the one of the cell vertex



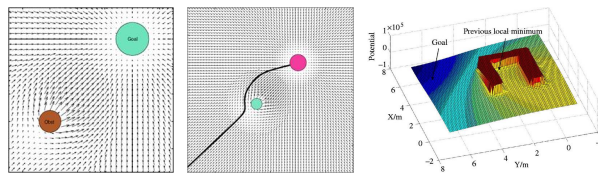
- The estimation can be further improve by “ray-shooting” technique combined with walking in triangulation (convex partitioning)

(Faigl, 2010)

Artificial Potential Field Method

- The idea is to create a function f that will provide a direction towards the goal for any configuration of the robot.
- Such a function is called **navigation function** and $-\nabla f(q)$ points to the goal.
- Create a **potential field** that will attract robot towards the goal q_f while obstacles will generate **repulsive potential** repelling the robot away from the obstacles.

The navigation function is a sum of potentials.



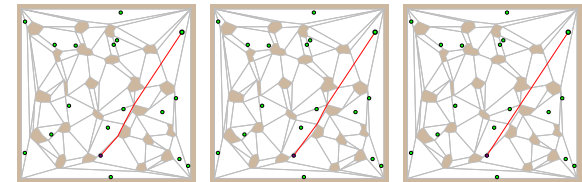
Such a potential function can have several local minima.

Topics Discussed

- Motion planning problem
- Path planning methods – overview
- Notation of configuration space
- Shortest-Path Roadmaps
- Voronoi diagram based planning
- Cell decomposition method
- Artificial potential field method
- Next: Grid-based path planning**

Path Refinement

- Testing collision of the point p with particular vertices of the estimation of the shortest path
 - Let the initial path estimation from p to p_g be a sequence of k vertices $(p, v_1, \dots, v_k, p_g)$
 - We can iteratively test if the segment (p, v_i) , $1 < i \leq k$ is collision free up to (p, p_g)



path over v_0

path over v_1

full refinement

With precomputed structures, it allows to estimate the shortest path in units of microseconds

Avoiding Local Minima in Artificial Potential Field

- Consider harmonic functions that have only one extremum

$$\nabla^2 f(q) = 0$$

- Finite element method

Dirichlet and Neumann boundary conditions



J. Mačák, Master thesis, CTU, 2009