## Solving Extensive-form games

November 19, 2014

## Finding Nash equilibrium in imperfect information EFG

- Existence of pure strategy Nash not guaranteed
- Convert to normal-form game and find NE there

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- Backward induction does not work
- Approximative regret minimizing methods
- LP based on sequence form
- Iterative building with sequence-form LP

- Based on sequences
- Realization plans
- Even more compressed than behavioral strategies
- Extended utility function

## ${\sf Sequence-form} \ {\sf LP}$

- $I(\sigma)$  returns the information set where the last action of the sequence  $\sigma$  took place
- seq(l<sub>i</sub>) returns sequence of player i which leads to l<sub>i</sub> (there can be only one)

$$\max_{r_{1},v} v(l([]))$$
(1)  
s.t.  $v(l(\sigma_{2})) \leq \sum_{l_{2} \in seq_{1}(l_{2}) = \sigma_{2}} v(l_{2}) + \sum_{\sigma_{1} \in \Sigma_{1}} g_{1}(\sigma_{1},\sigma_{2})r_{1}(\sigma_{1}), \quad \forall \sigma_{2} \in \Sigma_{2}$   
(2)  
 $r_{1}([]) = 1$ (3)  
 $r_{1}(\sigma_{1}) = \sum_{a \in A(l_{1})} r_{1}(\sigma_{1} \cdot a), \quad \forall l_{1} \in \mathcal{I}_{1}, \sigma_{1} = seq(l_{1})$ (4)  
 $r_{1}(\sigma_{1}) \geq 0, \quad \forall \sigma_{1} \in \Sigma_{1}$ (5)

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