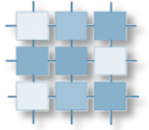
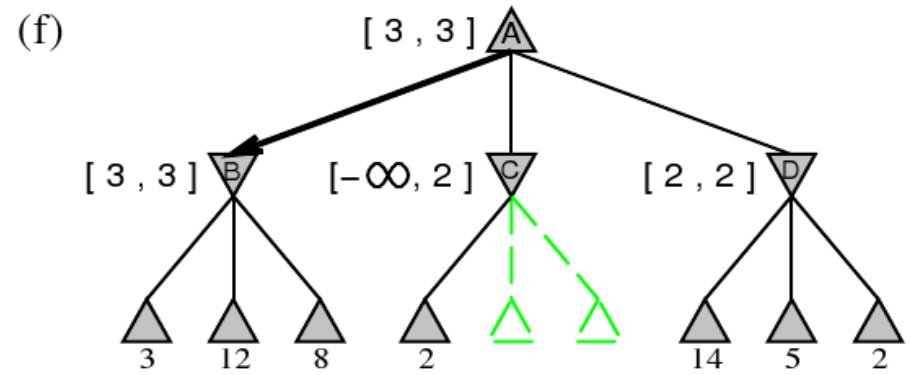
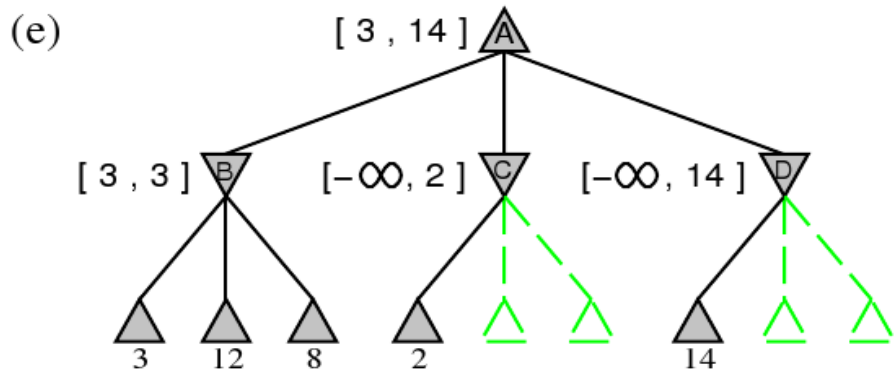
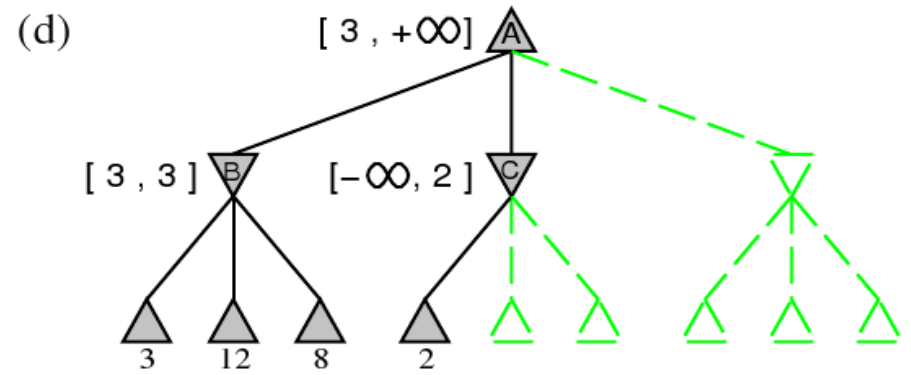
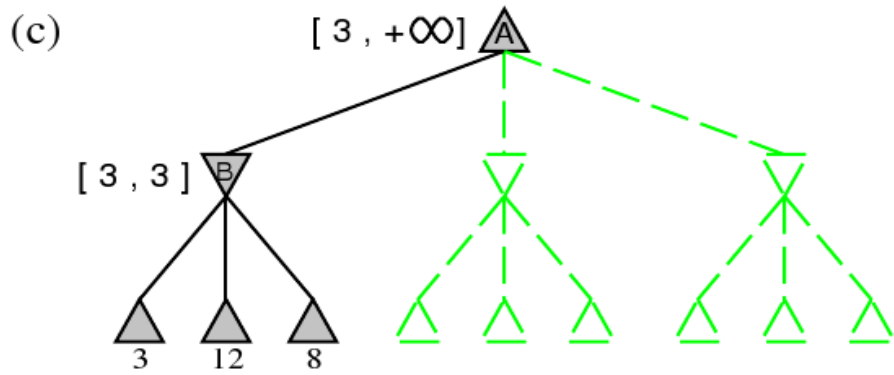
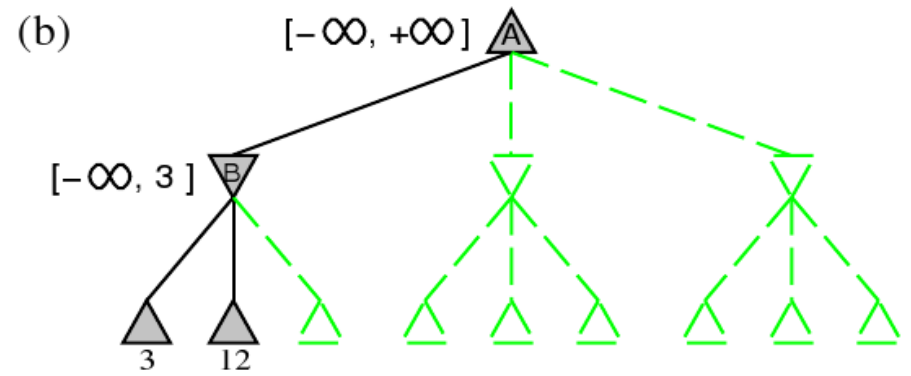
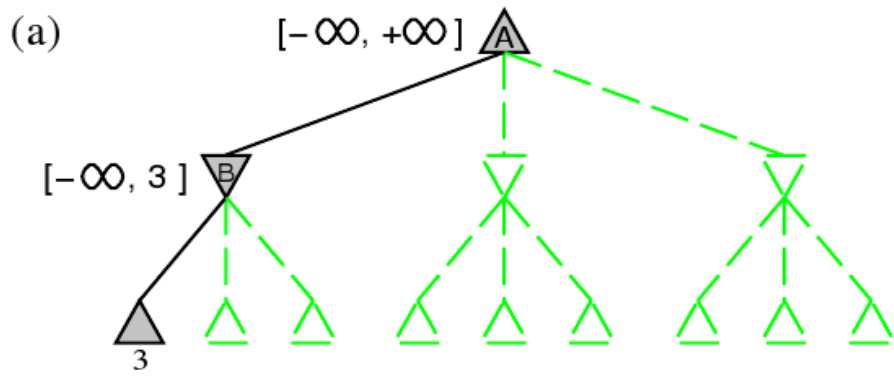


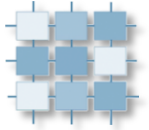
# Two-players Games



# Reminder and

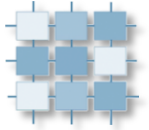
- two-players games
- zero-sum games
- no chance nodes
- perfect information
  
- chess, checkers, ...
  
- note – the algorithms return a value, but we need a solution
- backup function





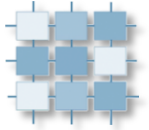
## Reminder (2)

- **function** alphabeta(node, depth,  $\alpha$ ,  $\beta$ , Player)
- **if** (depth = 0 or node is a terminal node) **return** the heuristic value of node
- **if** (Player = MaxPlayer)
- **for each** child of node
- $\alpha := \max(\alpha, \text{alphabeta}(\text{child}, \text{depth}-1, \alpha, \beta, \text{not}(\text{Player}))$ )
- **if** ( $\beta \leq \alpha$ ) **break**
- **return**  $\alpha$
- **else**
- **for each** child of node
- $\beta := \min(\beta, \text{alphabeta}(\text{child}, \text{depth}-1, \alpha, \beta, \text{not}(\text{Player}))$ )
- **if** ( $\beta \leq \alpha$ ) **break**
- **return**  $\beta$



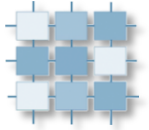
# Step 1 - Negamax

- **function** **negamax**(node, depth,  $\alpha$ ,  $\beta$ , color)
- **if** (depth = 0 or node is a terminal node) **return** the heuristic value of node
- **if** (Player = MaxPlayer)
- **for each** child of node
- $\alpha := \max(\alpha, \text{negamax}(\text{child}, \text{depth}-1, -\beta, -\alpha, \text{color}))$
- **if** ( $\beta \leq \alpha$ ) **break**
- **return**  $\alpha$
- **else**
- **for each** child of node
- $\beta := \min(\beta, \text{alphabeta}(\text{child}, \text{depth}-1, \alpha, \beta, \text{not}(\text{Player})))$
- **if** ( $\beta \leq \alpha$ ) **break**
- **return**  $\beta$



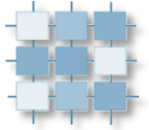
## Step 2 – Aspiration Search

- $[\alpha, \beta]$  interval – window
- alphabeta initialization  $[-\infty, +\infty]$
- what if we use  $[\alpha_0, \beta_0]$ 
  - $x = \text{alphabeta}(\text{node}, \text{depth}, \alpha_0, \beta_0, \text{player})$
  - $\alpha_0 \leq x \leq \beta_0$  - we found a solution
  - $x \leq \alpha_0$  - failing low (run again with  $[-\infty, x]$ )
  - $x \geq \beta_0$  - failing high (run again with  $[x, +\infty]$ )



## Step 3 – Scout – Idea

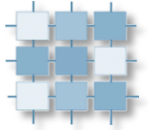
- assume we are in a MAX node
- we are about to search a child 'c'
- we already have obtained a lower bound ' $\alpha$ '
  
- Is it worth searching the branch 'c'?
  
- we need to have some test ...



## Step 3 – Scout – A Test

- what we really need at that moment is a bound (not the precise value)
- Remember Aspiration Search?
  - $x \leq \alpha_0$  - failing low (we know, that solution is  $\leq x$ )
  - $x \geq \beta_0$  - failing high (we know, that solution is  $\geq x$ )
- What if we use a null-window  $[\alpha, \alpha+1]$  (or  $[\alpha, \alpha]$ )?
  - we obtain a bound ...

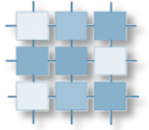




# Step 3 – NegaScout

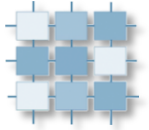
**function** negascout(node, depth,  $\alpha$ ,  $\beta$ , color)

- **if** ((depth = 0) or (node is a terminal node)) **return** eval(node)
- $b := \beta$
- **for each** child of node
- $v :=$  -negascout(child, depth-1, -b, - $\alpha$ , -color )
- **if** ((  $\alpha < v < \beta$  ) and (child is not the first child))
- $v :=$  -negascout(child, depth-1, - $\beta$ , - $\alpha$ , -color )
- $\alpha :=$  max( $\alpha$ , v)
- **if** ( $\beta \leq \alpha$ ) **break**
- $b := \alpha + 1$
- **return**  $\alpha$



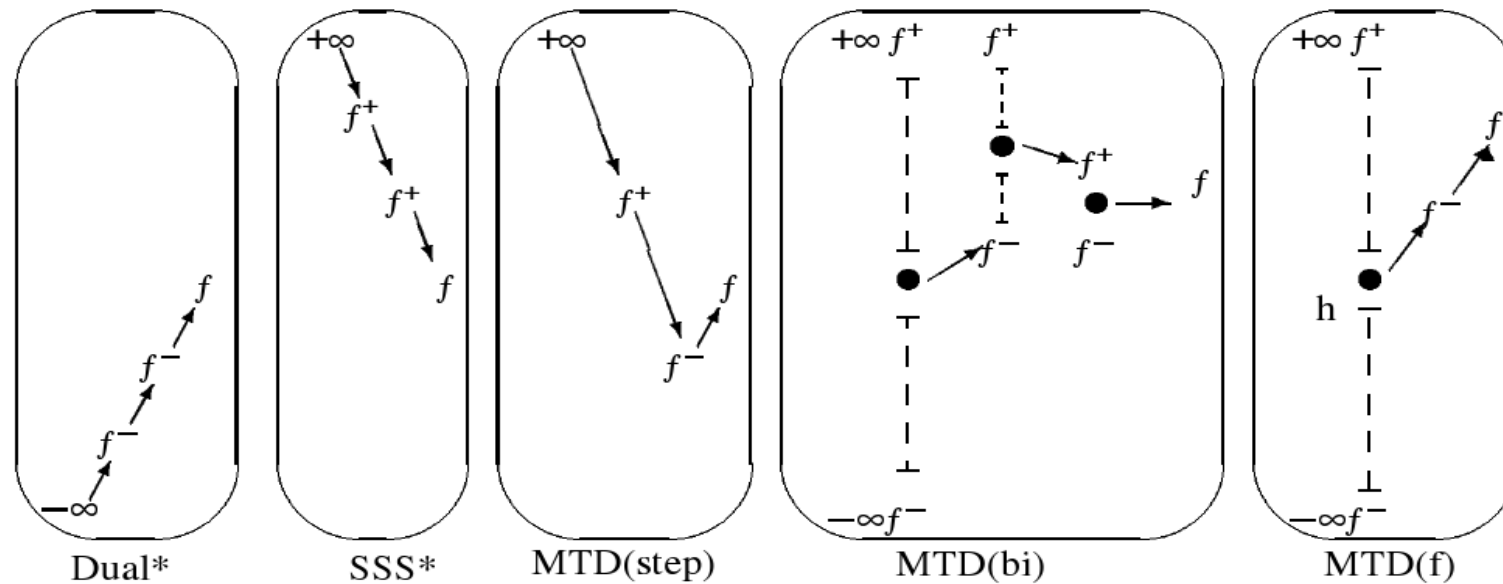
## Step 3 – NegaScout

- also termed Principal Variation Search (PVS)
- dominates alphabeta (never evaluates more nodes than alphabeta)
- depends on the move ordering
- can benefit from transposition tables
- generally 10-20% faster compared to alpha-beta

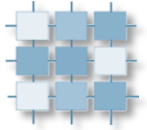


# Step 4 – MTD

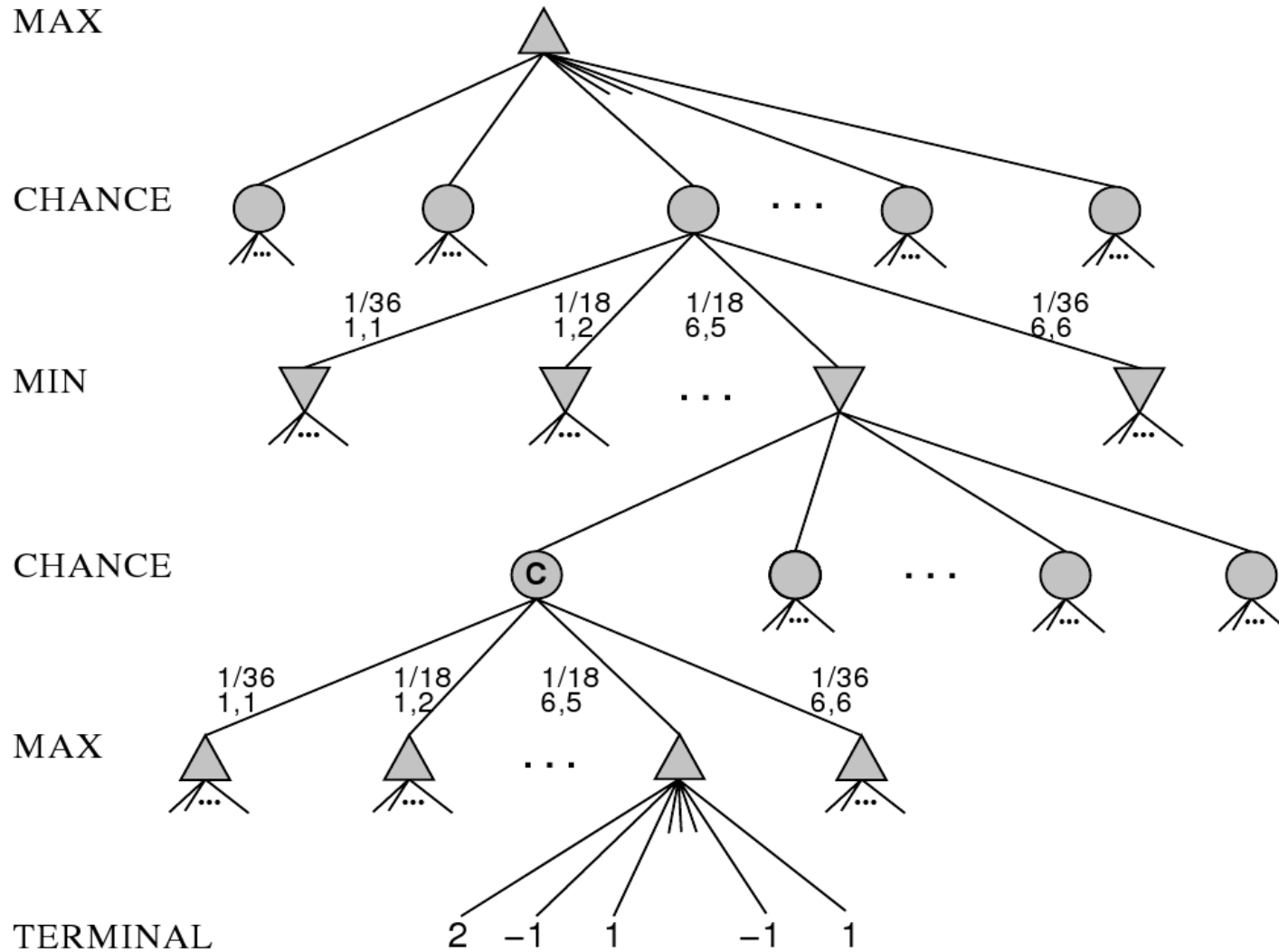
- Memory-enhanced Test Driver

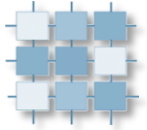


- Best-first fixed-depth minimax algorithms. Plaat et. al. , In *Artificial Intelligence*, Volume 87, Issues 1-2, November 1996, Pages 255-293

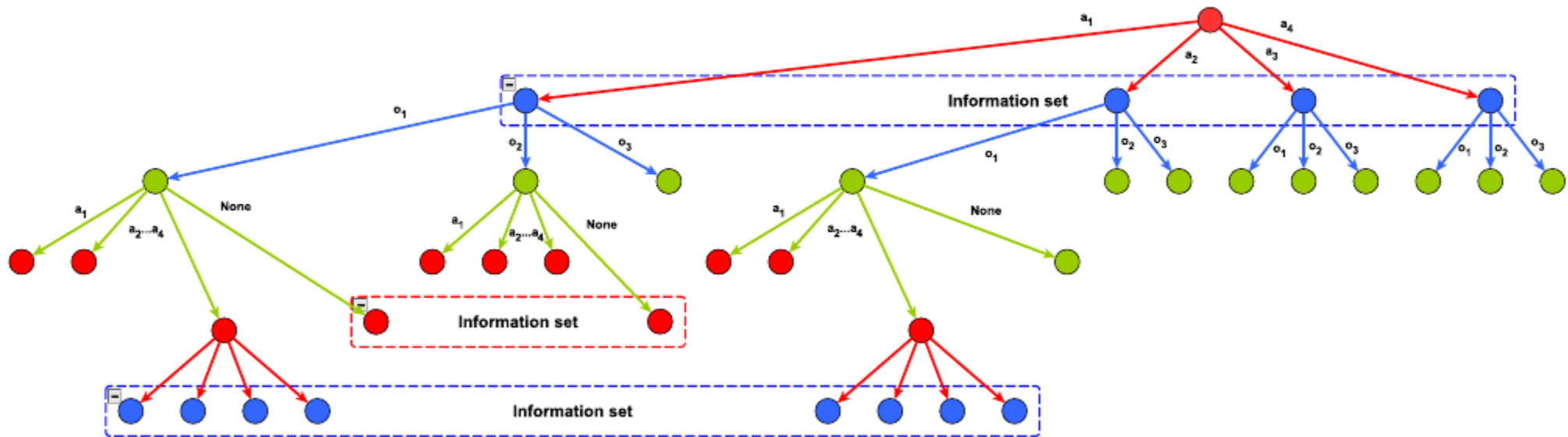


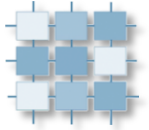
# Other Games - Chance nodes





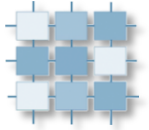
# Other Games – Imperfect Information





# Games and AI

- checkers – 1994 Chinook ... now a solved game (the program cannot lose)
  - chess – 1997 Deep Blue, ..., computers are now too strong
  - go – best human players are still undefeated, but ... (see <http://www.computer-go.info/h-c/index.html>)
  - poker – 2008 best program (Polaris) can beat a human master
  - ... and many, many others (Hex, Havannah, ...)
- 
- University of Alberta
  - Computer Olympiad



# Challenges?

- simultaneous moves, imperfect information
- durative moves (asynchronous chess, Google AI Challenge, ...)
- General Game Playing
  - an algorithm receives rules of the game and has to play
- ARIMAA (created in 2002)
  - $BF \approx 17,000$ ; no opening books; very few patterns
  - easy for people, very difficult for an algorithm
- using a 'real-AI-algorithms' in computer video-games
  - very few examples: F.E.A.R., World In Conflict, ...