

## *Logical reasoning and programming, lab session 2*

(October 2, 2023)

- 2.1** Use <http://fmv.jku.at/limboole/> on  $\varphi = (a \rightarrow (c \wedge d)) \vee (b \rightarrow (c \wedge e))$ .
- 2.2** Derive the empty clause from  $\{\bar{a}, b\}, \{\bar{b}, c\}, \{a, \bar{c}\}, \{a, b, c\}, \{\bar{a}, \bar{b}, \bar{c}\}$  using resolution.
- 2.3** A clause  $c_1 = \{l, l_1, \dots, l_n\}$  is blocked in  $\varphi$  by  $l$  if for every clause  $c_2 \in \varphi$  such that  $\bar{l} \in c_2$  the resolvent of  $c_1$  and  $c_2$  is a tautology. Prove that if a clause  $d$  is blocked (by a literal) in  $\psi$ , then  $\psi \in \text{SAT}$  iff  $\psi \setminus d \in \text{SAT}$ .
- 2.4** Formulate graph coloring (a vertex coloring) as a SAT problem. Namely, given a graph  $G$ , does  $G$  admit a proper vertex coloring with  $k$  colors?  
Discuss various possibilities how to formulate the problem. Moreover, are really all the constraints necessary?
- 2.5** Check this video, where zChaff colors the McGregor graph.
- 2.6** Let  $\varphi$  be a formula in CNF such that it contains only Horn clauses; they contain at most one positive literal. Show that SAT for  $\varphi$  is decidable in polynomial time.  
*Hint:* Perform all the unit propagations first.
- 2.7** In fact, it is possible to improve the previous algorithm in such a way that it works in linear time.
- 2.8** Express the pigeonhole principle in propositional logic. Namely, define a propositional formula  $\text{PHP}_n^{n+1}$ , which says that you have  $n + 1$  pigeons,  $n$  holes, every pigeon has a hole, and no two pigeons sit in the same hole.
- 2.9** Try PicoSAT/pycosat on  $\text{PHP}_n^{n+1}$  for small values of  $n$ . What is the maximal value of  $n$  for which you can solve  $\text{PHP}_n^{n+1}$  in one minute?
- 2.10** Install PySAT.
- 2.11** If you want to play with SAT solving a bit, then a standard exercise is to formalize Sudoku as a SAT problem and hence produce a Sudoku solver. Write a program that generates a problem specification in the DIMACS format in such a way that it is possible to specify an input (a partially completed grid) by appending<sup>1</sup> clauses saying which variables are true. You can use MiniSat or pycosat and some input is available from here, where each line has format  $XYZ$  meaning there is  $Z$  in cell  $(X, Y)$ .  
By the way, is it possible to obtain also a generator of Sudoku puzzles this way?
- 2.12** Check the details of two watched literals, for example, these and these slides illustrate the data structure.

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<sup>1</sup>Note that this changes the number of clauses, a parameter specified in the DIMACS format.