

Introduction to Prolog (2)

Using some slides from Peter Flach's materials

Recap from the last lecture... Herbrand...

- Herbrand universe, Herbrand base, Herbrand interpretation, Herbrand model

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person(X) :- teaches(X,_).
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```

- **Herbrand universe:**

- {alice, prolog}

- **Herbrand base:**

- {teaches(alice, prolog), teaches(prolog,alice), teaches(prolog,prolog), teaches(alice,alice), course(alice), course(prolog), person(alice), person(prolog)}

Recap from the last lecture... Hebrand...

- **Herbrand universe, Herbrand base, Herbrand interpretation, Herbrand model**

- **For instance:**

```
teaches(alice, prolog).  
course(X) :- teaches(_,X).  
person(X) :- teaches(X,_).
```

- **Herbrand universe:**

- {alice, prolog}

- **Herbrand base:**

- {teaches(alice, prolog), teaches(prolog,alice), teaches(prolog,prolog), teaches(alice,alice), course(alice), course(prolog), person(alice), person(prolog)}

- **Herbrand interpretation:**

- {teaches(alice, prolog), course(prolog), person(alice)}
- ...which is also a Herbrand model of the program.

Recap from the last lecture... Functions

- **Function symbols** allow us to construct composite terms (besides that they also allow us to do **Skolemization**).

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- **Example:**

$\text{num}(0) .$

$\text{num}(s(X)) \text{ :- num}(X) .$

Recap from the last lecture... Functions

- **Function symbols** allow us to construct composite terms (besides that they also allow us to do **Skolemization**).

- **Example:**

$\text{num}(0) .$

$\text{num}(s(X)) \text{ :- num}(X) .$

- **Herbrand universe:** $\{0, s(0), s(s(0)), \dots\}$
- **Herbrand base:** $\{\text{num}(0), \text{num}(s(0)), \text{num}(s(s(0))), \dots\}$
- In this case, we do not have a finite Herbrand model.

A Digression: Skolemization

“Everybody knows somebody.”

A Digression: Skolemization

“Everybody knows somebody.”

Skolemization to avoid an existential quantifier

```
knows (X, person_known_by (X)) .
```

functor

term

complex term

```
knows (peter, person_known_by (peter)) .
```

```
knows (anna, person_known_by (anna)) .
```

```
knows (paul, person_known_by (paul)) .
```

...

Recap from the last lecture... Logical consequence

- A clause C is a **logical consequence** of a program (set of clauses) P iff every model of P is a model of C .

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Program **P**:

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Recap from the last lecture... Logical consequence

- A clause **C** is a **logical consequence** of a program (set of clauses) **P** iff every model of **P** is a model of **C**.

Program **P**:

```
num(0) .  
num(s(X)) :- num(X) .
```

Some logical consequences of **P**:

```
num(s(s(0))) .  
num(0)  
num(s(s(X))) :- num(X) .
```

An Example: Prague Metro



A Prolog DB (1)

```
connected(nemocnice_motol,petriny,green).
connected(petriny,nadrazi_veleslavin,green).
connected(nadrazi_veleslavin,borislavka,green).
connected(borislavka,dejvicka,green).
connected(dejvicka,hradcanska,green).
connected(hradcanska,malostranska,green).
connected(malostranska,staromestska,green).
connected(staromestska,mustek,green).
connected(mustek,muzeum,green).
connected(muzeum,namesti_miru,green).
connected(namesti_miru,jiriho_z_podebrad,green).
connected(jiriho_z_podebrad,flora,green).
connected(flora,zelivskeho,green).
connected(zelivskeho,strasnicka,green).
connected(strasnicka,skalka,green).
connected(skalka,depo_hostivar,green).
```


A Prolog DB (2)

```
connected(letnany,prosek,red) .
connected(prosek,trizkov,red) .
connected(trizkov,ladvi,red) .
connected(ladvi,kobylisy,red) .
connected(kobylisy,nadrazi_holesovice,red) .
connected(nadrazi_holesovice,vltavska,red) .
connected(vltavska,florenc,red) .
connected(florenc,hlavni_nadrazi,red) .
connected(hlavni_nadrazi,muzeum,red) .
connected(muzeum,i_p_pavlova,red) .
connected(i_p_pavlova,vysehrad,red) .
connected(vysehrad,prazskeho_povstani,red) .
connected(prazskeho_povstani,pankrac,red) .
connected(pankrac,budejovicka,red) .
connected(budejovicka,kacerov,red) .
connected(kacerov,roztyly,red) .
connected(roztyly,chodov,red) .
connected(chodov,opatov,red) .
connected(opatov,haje,red) .
```

A Prolog DB (3)

```
connected(zlicin, stodulky, yellow) .
connected(stodulky, luka, yellow) .
connected(luka, luziny, yellow) .
connected(luziny, hurka, yellow) .
connected(hurka, nove_butovice, yellow) .
connected(nove_butovice, jinonice, yellow) .
connected(jinonice, radlicka, yellow) .
connected(radlicka, smichov, yellow) .
connected(smichov, andel, yellow) .
connected(andel, karlovo_namesti, yellow) .
connected(karlovo_namesti, narodni_trida, yellow) .
connected(narodni_trida, mustek, yellow) .
connected(mustek, namesti_republiky, yellow) .
connected(namesti_republiky, florenc, yellow) .
connected(florenc, krizikova, yellow) .
connected(krizikova, invalidovna, yellow) .
connected(invalidovna, palmovka, yellow) .
connected(palmovka, ceskomoravska, yellow) .
connected(ceskomoravska, vysocanska, yellow) .
connected(vysocanska, kolbenova, yellow) .
connected(kolbenova, hloubetin, yellow) .
connected(hloubetin, rajska_zahrada, yellow) .
connected(rajska_zahrada, cerny_most, yellow) .
```

“Nearby”

Two stations are nearby if they are on the same line with at most one other station in between:

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Two stations are nearby if they are on the same line with at most one other station in between:

```
nearby(zlicin, luka) .  
nearby(luka, zlicin) .  
nearby(zlicin, stodulky) .  
nearby(stodulky, zlicin) .  
nearby(luka, luziny) .  
nearby(luziny, luka) .  
nearby(luka, hurka) .
```

“Nearby”

Two stations are nearby if they are on the same line with at most one other station in between:

```
nearby(zlicin,luka) .
nearby(luka,zlicin) .
nearby(zlicin,stodulky) .
nearby(stodulky,zlicin) .
nearby(luka,luziny) .
nearby(luziny,luka) .
nearby(luka,hurka) .
...
```

or better

```
nearby(X,Y):-connectedS(X,Y,L) .
nearby(X,Y):-connectedS(X,Z,L),connectedS(Z,Y,L) .
connectedS(X,Y,W):-connected(X,Y,W) .
connectedS(X,Y,W):-connected(Y,X,W) .
```

“Not too far”

Compare

```
nearby (X, Y) :- connectedS (X, Y, L) .
```

```
nearby (X, Y) :- connectedS (X, Z, L) , connectedS (Z, Y, L) .
```

with

```
not_too_far (X, Y) :- connectedS (X, Y, L) .
```

```
not_too_far (X, Y) :- connectedS (X, Z, L1) , connectedS (Z, Y, L2) .
```

“Not too far”

Compare

```
nearby (X, Y) :- connectedS (X, Y, L) .
```

```
nearby (X, Y) :- connectedS (X, Z, L) , connectedS (Z, Y, L) .
```

with

```
not_too_far (X, Y) :- connectedS (X, Y, L) .
```

```
not_too_far (X, Y) :- connectedS (X, Z, L1) , connectedS (Z, Y, L2) .
```

This can be rewritten with don't cares:

```
not_too_far (X, Y) :- connectedS (X, Y, _) .
```

```
not_too_far (X, Y) :- connectedS (X, Z, _) , connectedS (Z, Y, _) .
```

?-nearby(mustek, W)

?-nearby(mustek,W)

nearby(X1,Y1):-connectedS(X1,Y1,L1)

?-nearby(mustek,W)



nearby(X1,Y1):-connectedS(X1,Y1,L1)

{X1->mustek, Y1->W}

?-connectedS(mustek,W,L1)

?-nearby(mustek,W)



nearby(X1,Y1):-connectedS(X1,Y1,L1)

{X1->mustek, Y1->W}

?-connectedS(mustek,W,L1)

connectedS(X,Y,L) :- connected(X,Y,L)

?-nearby(mustek,W)



nearby(X1,Y1):-connectedS(X1,Y1,L1)

{X1->mustek, Y1->W}

?-connectedS(mustek,W,L1)



connectedS(X,Y,L) :- connected(X,Y,L)

{X->mustek, Y->W, L->L1}

?-connected(mustek,W,L1)

?-nearby(mustek,W)



nearby(X1,Y1):-connectedS(X1,Y1,L1)

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?-connectedS(mustek,W,L1)



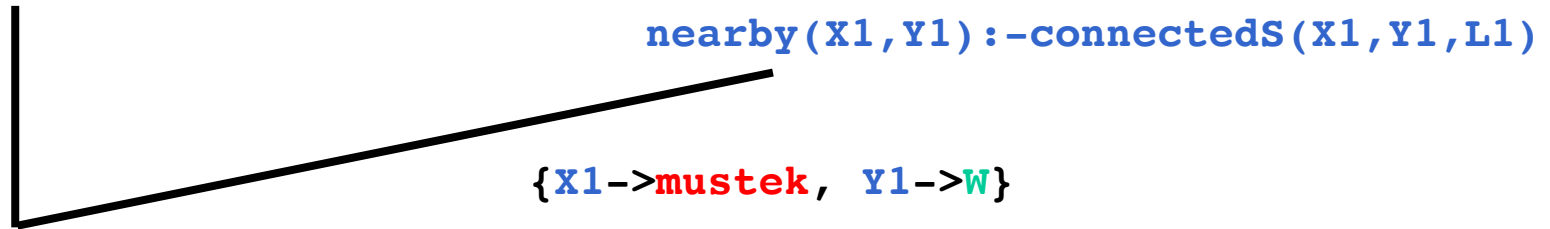
connectedS(X,Y,L) :- connected(X,Y,L)

{X->mustek, Y->W, L->L1}

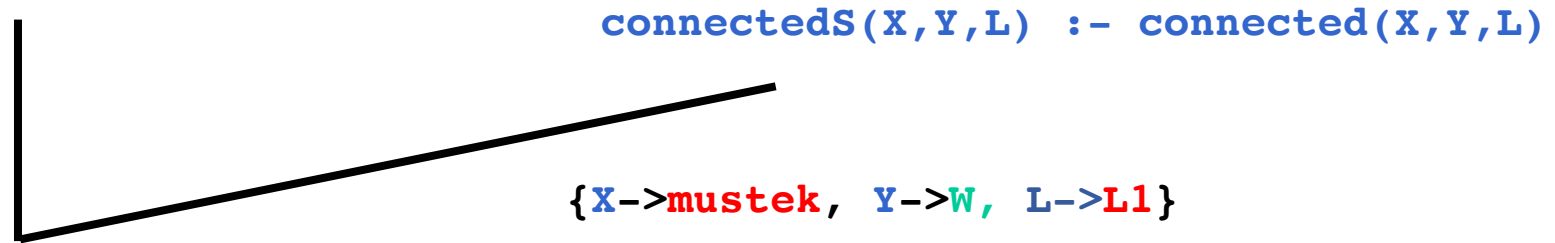
?-connected(mustek,W,L1)

connected(mustek,muzeum,green)

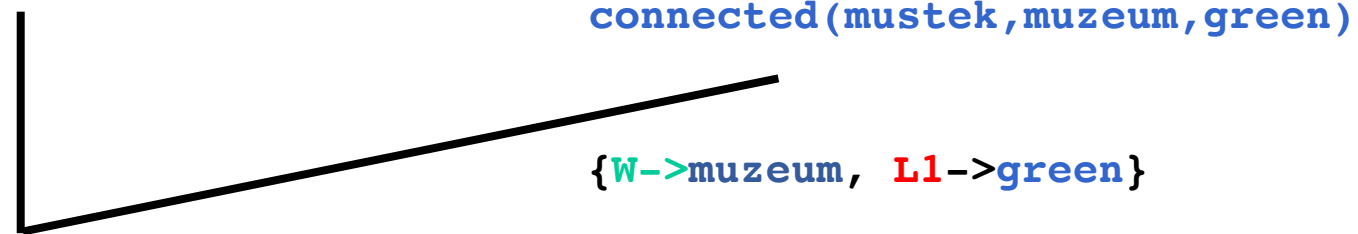
?-nearby(mustek,W)



?-connectedS(mustek,W,L1)



?-connected(mustek,W,L1)



[]

“Reachable”

A station is reachable from another if they are on the same line, or with one, two, ... changes:

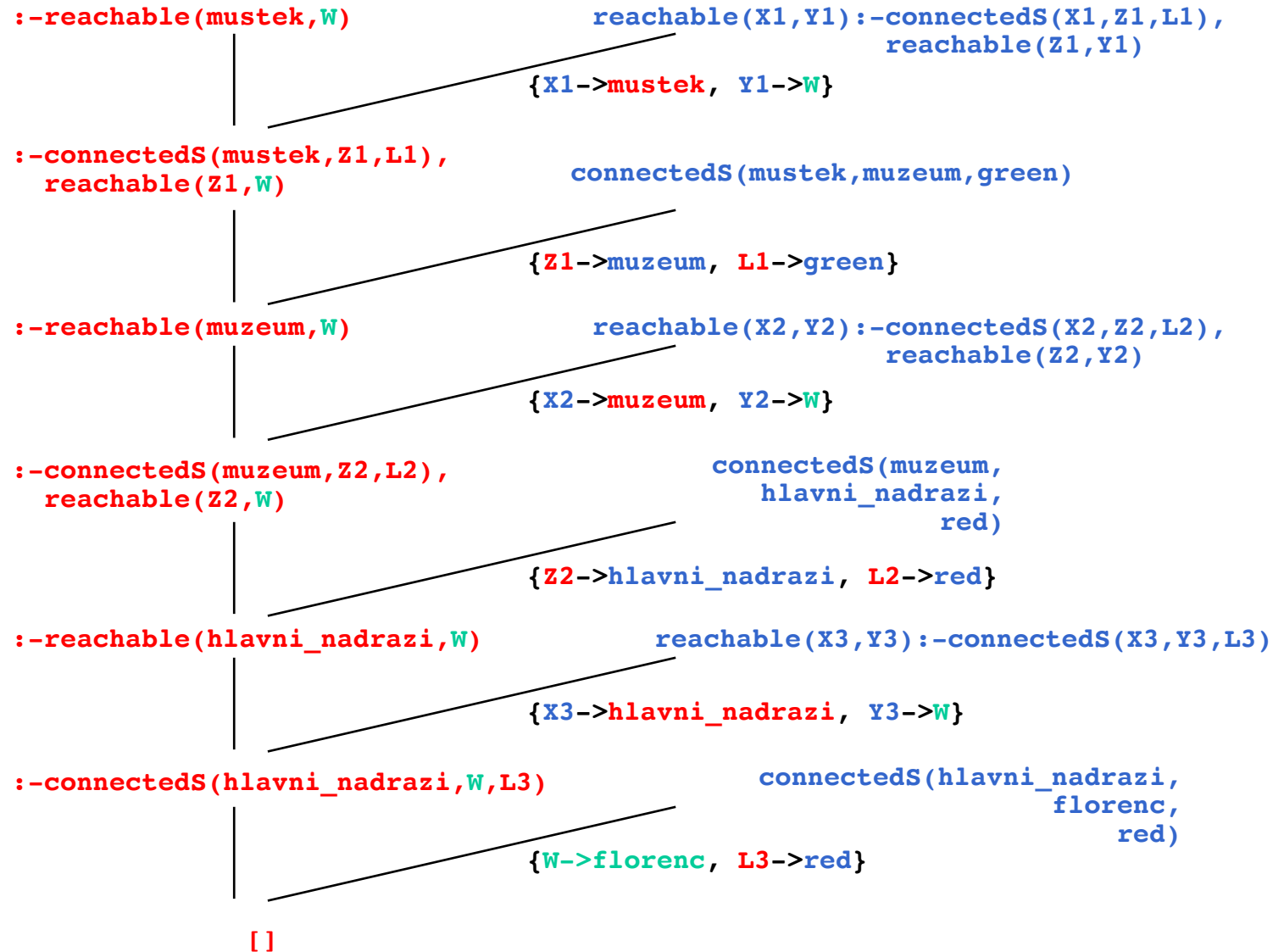
```
reachable (X, Y) :-connectedS (X, Y, L) .  
reachable (X, Y) :-connectedS (X, Z, L1) , connectedS (Z, Y, L2) .  
reachable (X, Y) :-connectedS (X, Z1, L1) , connectedS (Z1, Z2, L2) ,  
                    connectedS (Z2, Y, L3) .
```

...

or better

```
reachable (X, Y) :-connectedS (X, Y, L) .  
reachable (X, Y) :-connectedS (X, Z, L) , reachable (Z, Y) .
```

Note: To make the example more compact, we assume that we have the predicate `connectedS` given as ground facts.



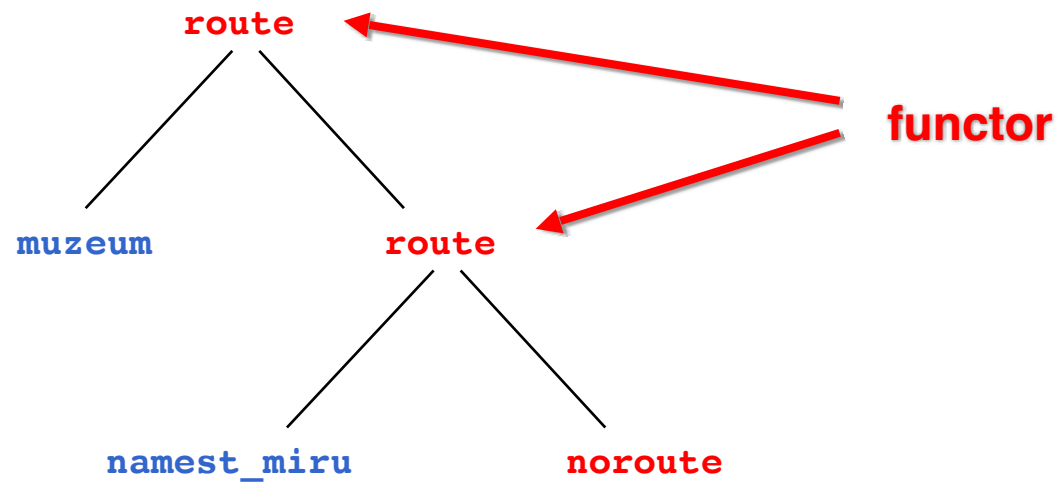
There is a catch!

- The answers that we get depend on the exact way Prolog works inside. We will talk about that next time.

“Recording the Path”

```
reachable(X,Y,noroute):-connected(X,Y,L).  
reachable(X,Y,route(Z,R)):-connected(X,Z,L),  
                             reachable(Z,Y,R).
```

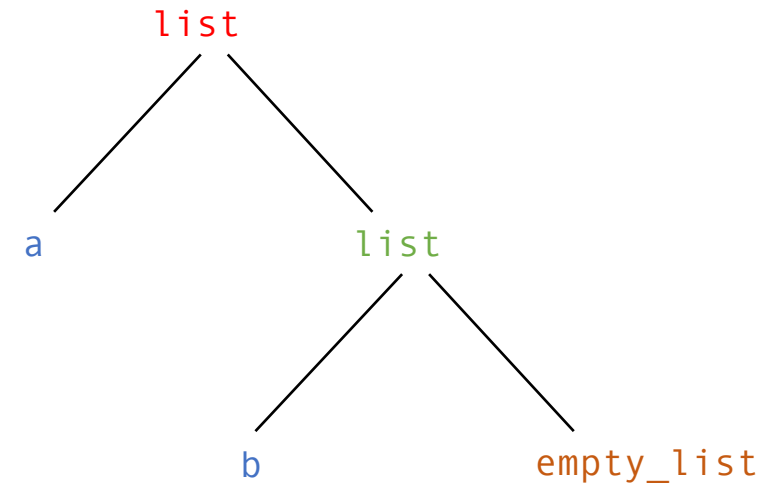
```
?-reachable(mustek,jiriho_z_podebrad,R).  
R = route(muzeum,route(namesti_miru,noroute));  
...
```



More examples (Still with our “own” lists)

- Here we will represent lists more or less like we represented routes (Prolog has built in support for lists but this is still to illustrate them as pure logic concepts), e.g.: `empty_list`, `list(X, empty_list)`, ...

```
list(a, list(b, empty_list))
```



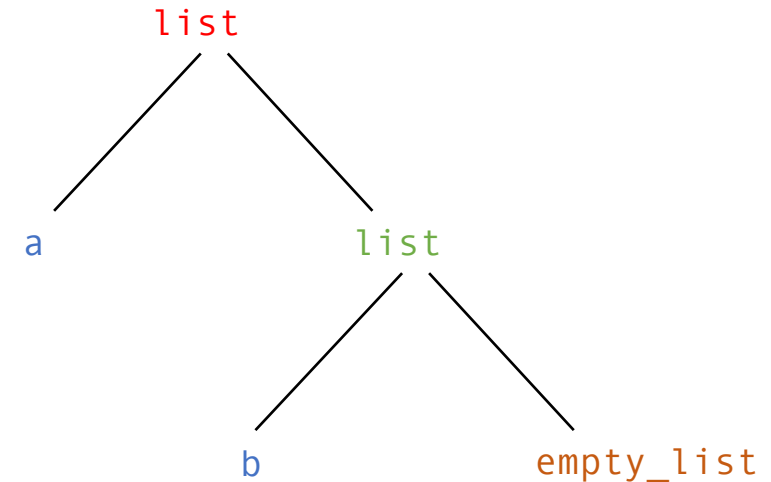
More examples (Still with our “own” lists)

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Example:

```
add_to_start(empty_list, X, list(X,empty_list)).  
add_to_start(list(H,T), X, list(X,list(H,T))).
```

```
list(a,list(b,empty_list))
```





Program x +

```
1 add_to_start(empty_list, X, list(X,empty_list)).  
2 add_to_start(list(H,T), X, list(X,list(H,T))).  
3  
4
```

```
add_to_start(empty_list, a, L1), add_to_start(L1, b, L2),  
add_to_start(L2, c, L3)
```

L1 = list(a, empty_list),

L2 = list(b, list(a, empty_list)),

L3 = list(c, list(b, list(a, empty_list)))

?-

```
add_to_start(empty_list, a, L1), add_to_start(L1, b,  
L2), add_to_start(L2, c, L3)
```

Examples▲

History▲

Solutions▲

 table results

Run!



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Program +

```
1 add_to_start(empty_list, X, list(X,empty_list)).  
2 add_to_start(list(H,T), X, list(X,list(H,T))).  
3  
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```
add_to_start(empty_list, a, L1), add_to_start(L1, b, L2),  
add_to_start(L2, c, L3)
```

```
L1 = list(a, empty_list),
```

```
L2 = list(b, list(a, empty_list)),
```

```
L3 = list(c, list(b, list(a, empty_list)))
```

?-

```
add_to_start(empty_list, a, L1), add_to_start(L1, b,  
L2), add_to_start(L2, c, L3)
```

Examples History Solutions

 table results

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Program x +

```
1 add_to_start(empty_list, X, list(X,empty_list)).  
2 add_to_start(list(H,T), X, list(X,list(H,T))).  
3  
4
```

```
add_to_start(empty_list, a, L1), add_to_start(L1, b, L2),  
add_to_start(L2, c, L3)
```

```
L1 = list(a, empty_list),
```

```
L2 = list(b, list(a, empty_list)),
```

```
L3 = list(c, list(b, list(a, empty_list)))
```

?-

```
add_to_start(empty_list, a, L1), add_to_start(L1, b,  
L2), add_to_start(L2, c, L3)
```

Examples▲

History▲

Solutions▲

 table results

Run!

Yet another example

```
add_to_end(empty_list, X, list(X,empty_list)).
```

```
add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,X,T2).
```




Program

```
1 add_to_end(empty_list, X, list(X,empty_list)).  
2 add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,
```

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```

L1 = list(a, empty_list),

L2 = list(a, list(b, empty_list)),

L3 = list(a, list(b, list(c, empty_list)))

?-

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```

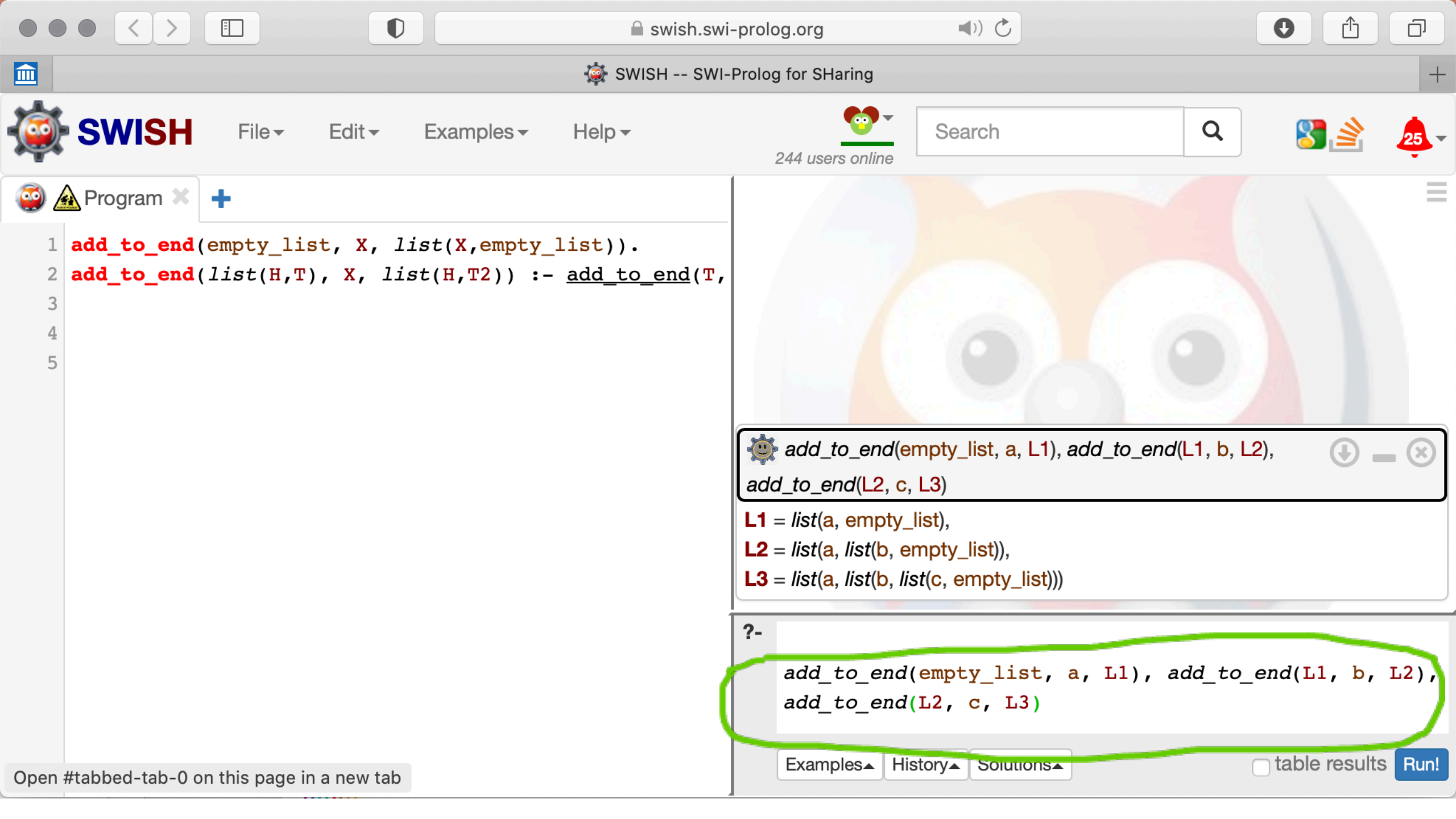
Examples

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Run!



Program

```
1 add_to_end(empty_list, X, list(X,empty_list)).  
2 add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,  
3  
4  
5
```

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```

```
L1 = list(a, empty_list),  
L2 = list(a, list(b, empty_list)),  
L3 = list(a, list(b, list(c, empty_list)))
```

?-

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```



Program

```
1 add_to_end(empty_list, X, list(X,empty_list)).  
2 add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,
```

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```

```
L1 = list(a, empty_list),
```

```
L2 = list(a, list(b, empty_list)),
```

```
L3 = list(a, list(b, list(c, empty_list)))
```

?-

```
add_to_end(empty_list, a, L1), add_to_end(L1, b, L2),  
add_to_end(L2, c, L3)
```

Examples

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Run!

Removing last element

We had:

```
add_to_end(empty_list, X, list(X,empty_list)).  
add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,X,T2).
```

We can use:

```
remove_last(List,NewList) :- add_to_end(NewList,_,List).
```

Removing last element

We had:

```
add_to_end(empty_list, X, list(X,empty_list)).  
add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,X,T2).
```

in *out*


We can use:

```
remove_last(List,NewList) :- add_to_end(NewList,_,List).
```

out *in*

Program  

```
1 add_to_end(empty_list, X, list(X,empty_list)).  
2 add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,X,T2).  
3  
4 remove_last(List,NewList) :- add_to_end(NewList,_,List).  
5  
6  
7
```

```
 add_to_end(empty_list, a, L1),  
add_to_end(L1, b, L2), add_to_end(L2, c, L3),  
remove_last(L3,L4)
```

L1 = list(a, empty_list),

L2 = L4, **L4** = list(a, list(b, empty_list)),

L3 = list(a, list(b, list(c, empty_list)))

Next 10 100 1,000 Stop

?-

```
add_to_end(empty_list, a, L1),  
add_to_end(L1, b, L2), add_to_end(L2, c,  
L3), remove_last(L3,L4)|
```

Examples History Solutions

table results **Run!**



Program x +

```

1 add_to_end(empty_list, X, list(X,empty_list)).
2 add_to_end(list(H,T), X, list(H,T2)) :- add_to_end(T,X,T2).
3
4 remove_last(List,NewList) :- add_to_end(NewList,_,List).
5
6
7

```

```

add_to_end(empty_list, a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c, L3),
remove_last(L3,L4)

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L1 = list(a, empty_list),

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Next 10 100 1,000 Stop

?-

```

add_to_end(empty_list, a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c,
L3), remove_last(L3,L4)|

```

Examples History Solutions

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Run!



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1 add_to_end(empty_list, X, list(X,empty_list)).
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4 remove_last(List,NewList) :- add_to_end(NewList,_,List).
```

```
add_to_end(empty_list, a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c, L3),
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Next 10 100 1,000 Stop

?-

```
add_to_end(empty_list, a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c,
L3), remove_last(L3,L4)|
```

Examples History Solutions

table results Run!

Built-in Lists in Prolog

Examples of lists:

`[], [a, b, c], [a, X, c, d], ...`

Instead of writing `list(H, T)`, we write `[H|T]` in Prolog.

```
add_to_end(empty_list, X, list(X, empty_list)).
```

```
add_to_end(list(H, T), X, list(H, T2)) :- add_to_end(T, X, T2).
```

```
add_to_end([], X, [X]).
```

```
add_to_end([H|T], X, [H|T2]) :- add_to_end(T, X, T2).
```

```
1 add_to_end([], X, [X]).
2 add_to_end([H|T], X, [H|T2]) :- add_to_end(T,X,T2).
3
4
5 remove_last(List,NewList) :- add_to_end(NewList,_,List).
6
7
8
```

```
add_to_end([], a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c, L3),
remove_last(L3,L4)
L1 = list(a, empty_list),
L2 = L4, L4 = list(a, list(b, empty_list)),
L3 = list(a, list(b, list(c, empty_list)))
```

```
add_to_end([], a, L1), add_to_end(L1, b, L2),
add_to_end(L2, c, L3), remove_last(L3,L4)
L1 = [a],
L2 = L4, L4 = [a, b],
L3 = [a, b, c]
```

```
?-
add_to_end([], a, L1), add_to_end(L1, b,
L2), add_to_end(L2, c, L3),
remove_last(L3,L4)
```

```
1 add_to_end([], X, [X]).
2 add_to_end([H|T], X, [H|T2]) :- add_to_end(T,X,T2).
3
4
5 remove_last(List,NewList) :- add_to_end(NewList,_,List).
6
7
8
```

```
add_to_end([], a, L1),
add_to_end(L1, b, L2), add_to_end(L2, c, L3),
remove_last(L3,L4)
L1 = list(a, empty_list),
L2 = L4, L4 = list(a, list(b, empty_list)),
L3 = list(a, list(b, list(c, empty_list)))
```

```
add_to_end([], a, L1), add_to_end(L1, b, L2),
add_to_end(L2, c, L3), remove_last(L3,L4)
L1 = [a],
L2 = [a], L4 = [a, b],
L3 = [a, b, c]
```

```
?-
add_to_end([], a, L1), add_to_end(L1, b,
L2), add_to_end(L2, c, L3),
remove_last(L3,L4)
```

SLD Resolution in Detail

Unification

Recap from the last lecture... Substitutions

- **Example:**

```
likes(X,Y) :- dogPerson(X), dog(Y).
```

Recap from the last lecture... Substitutions

- **Example:**

`likes(X,Y) :- dogPerson(X), dog(Y).`

`S = {X/alice, Y/barney}`

Recap from the last lecture... Substitutions

- **Example:**

```
likes(X,Y) :- dogPerson(X), dog(Y).
```

```
S = {X/alice, Y/barney}
```

```
likes(alice, barney) :- dogPerson(alice), dog(barney)
```


Unification

- (We need unification to perform resolution.)
- **Unifier:** Given two atoms or terms A and B with disjoint (!) sets of variables, their **unifier** is a substitution θ such that $A\theta = B\theta$.

Unification

- (We need unification to perform resolution.)
- **Unifier:** Given two atoms or terms A and B with disjoint (!) sets of variables, their **unifier** is a substitution Θ such that $A\Theta = B\Theta$.
- A unifier Θ is more general (*technically speaking, we should be saying more or equally general*) than a unifier Θ' iff there exists a substitution σ such that $\Theta' = \Theta \sigma$.

Unification - Example

- ...A unifier Θ is more general (*technically speaking, we should be saying more or equally general*) than a unifier Θ' iff there exists a substitution σ such that $\Theta' = \Theta \sigma$. If Θ is more general than Θ' and not vice versa, then Θ is strictly more general than Θ' .

Example:

A = studies(X, Y)

B = studies(alice, Z)

$\Theta = \{X/\text{alice}, Y/Z\}$

$\Theta' = \{X/\text{alice}, Y/\text{bob}, Z/\text{bob}\}$

$\Theta' = \Theta \sigma$, where $\sigma = \{Z/\text{bob}\}$

Unification – Most General Unifier

- **Most General Unifier:** Given two atoms or terms A and B with disjoint sets of variables, their **most general unifier (MGU)** is a unifier θ such that there is no strictly more general unifier of A and B .

- **Example:**

$A = \text{studies}(X, Y)$

$B = \text{studies}(\text{alice}, Z)$

$\theta = \{X/\text{alice}, Y/Z\}$

Here, θ is an MGU of A and B .

The Hebrand Unification Algorithm

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

- **Initialization:** $\Theta = \{\}$, `failure = false`, **push** $A = B$ on the stack

- **Loop:**

- 1: **Pop** $x = y$ from the stack

- 2: **If** x and y are constants and $x == y$ **then continue.**

- 3: **Else if** x is a variable and x does not appear in y **then** /* “Occurs check” */

- 4: Replace x with y in the stack and in Θ . Add the substitution x/y to Θ .

- 5: **Else if** x is a variable and $x == y$ **then continue.**

- 6: **Else if** y is a variable and x is not a variable **then push** $y = x$ on the stack.

- 7: **Else if** x and y are compound and x is $f(t_1, \dots, t_k)$ and y is $f(u_1, \dots, u_k)$ **then**

- 8: **Push** on the stack: $t_1 = u_1$, $t_2 = u_2$, ..., $t_k = u_k$.

- 9: **Else set** `failure = true`, $\Theta = \{\}$ and **return.**

Until stack is empty.

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$ */* In Prolog, "=" means "unify" */*

$\Theta = \{\}$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$$g(X, h(X, b)) = g(a, Z)$$

$$Z = Y$$

$$\Theta = \{\}$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$$X = a$$

$$h(X, b) = Z$$

$$Z = Y$$

$$\Theta = \{\}$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$$h(a, b) = Z$$

$$Z = Y$$

$$\Theta = \{X/a\}$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$$Z = h(a, b)$$

$$Z = Y$$

$$\Theta = \{X/a\}$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

$$h(a, b) = Y$$

$$\Theta = \{X/a, Z/h(a, b)\}$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

(Adapted from Pierre M. Nugues: Language Processing with Perl and Prolog)

Stack:

EMPTY

$$\Theta = \{X/a, Z/h(a, b), Y/h(a, b)\}$$

Result of unification - MGU:

$$f(g(a, h(a, b)), h(a, b))$$

Example

$$f(g(X, h(X, b)), Z) = f(g(a, Z), Y)$$

The screenshot shows the SWISH web interface. The browser address bar displays `swish.swi-prolog.org`. The page title is "SWISH -- SWI-Prolog for SHaring". The interface includes a menu bar with "File", "Edit", "Examples", and "Help", a search bar, and a notification bell showing "25". A "Program" tab is active, showing a list of lines 1 through 4. The main content area displays a Prolog program and its solution, which is circled in red:

```
f(g(X, h(X, b)), Z) = f(g(a, Z), Y)
```

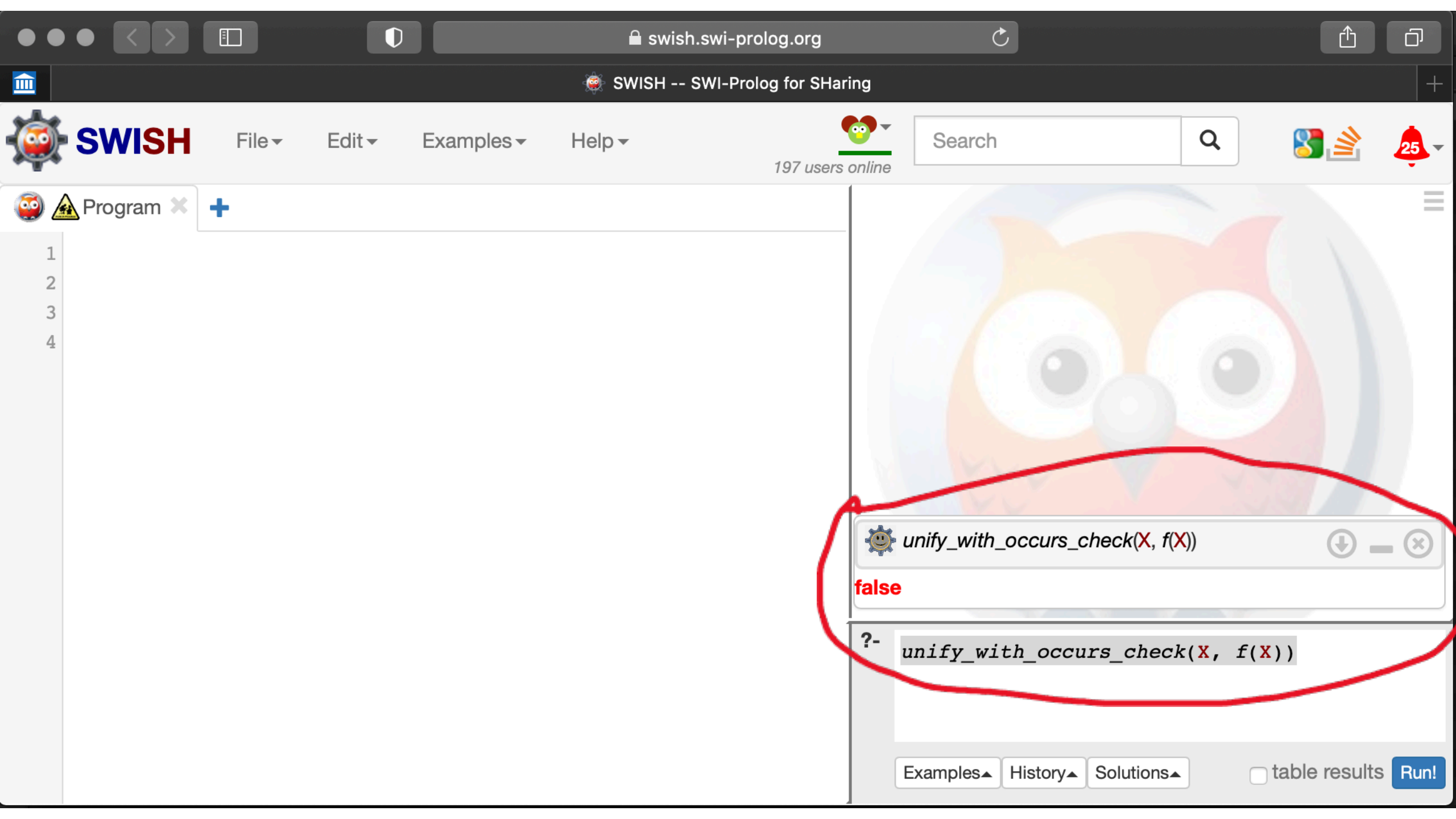
X = a,
Y = Z, Z = h(a, b)

?- `f(g(X, h(X, b)), Z) = f(g(a, Z), Y)`

At the bottom, there are buttons for "Examples", "History", "Solutions", a checkbox for "table results", and a "Run!" button.

The “Occurs Check”

- Line 3: **Else if** x is a variable and x does not appear in y **then** ...
- This is not implemented in many Prologs for efficiency reasons and programmers have to take care of it themselves. Without occurs check, Prolog inference is not really sound (this is a tradeoff for efficiency).
- If you want sound unification, use `unify_with_occurs_check/2`



1
2
3
4

`unify_with_occurs_check(X, f(X))`
false

?- `unify_with_occurs_check(X, f(X))`



Program

1
2
3
4

unify_with_occurs_check(X, f(X))

false

X = f(X)

X = f(X)

?- X = f(X)

Examples

History

Solutions

table results

Run!

Next week... SLD trees